

INFORMATION TO USERS

This was produced from a copy of a document sent to us for microfilming. While the most advanced technological means to photograph and reproduce this document have been used, the quality is heavily dependent upon the quality of the material submitted.

The following explanation of techniques is provided to help you understand markings or notations which may appear on this reproduction.

1. The sign or "target" for pages apparently lacking from the document photographed is "Missing Page(s)". If it was possible to obtain the missing page(s) or section, they are spliced into the film along with adjacent pages. This may have necessitated cutting through an image and duplicating adjacent pages to assure you of complete continuity.
2. When an image on the film is obliterated with a round black mark it is an indication that the film inspector noticed either blurred copy because of movement during exposure, or duplicate copy. Unless we meant to delete copyrighted materials that should not have been filmed, you will find a good image of the page in the adjacent frame.
3. When a map, drawing or chart, etc., is part of the material being photographed the photographer has followed a definite method in "sectioning" the material. It is customary to begin filming at the upper left hand corner of a large sheet and to continue from left to right in equal sections with small overlaps. If necessary, sectioning is continued again—beginning below the first row and continuing on until complete.
4. For any illustrations that cannot be reproduced satisfactorily by xerography, photographic prints can be purchased at additional cost and tipped into your xerographic copy. Requests can be made to our Dissertations Customer Services Department.
5. Some pages in any document may have indistinct print. In all cases we have filmed the best available copy.

University
Microfilms
International

300 N. ZEEB ROAD, ANN ARBOR, MI 48106
18 BEDFORD ROW, LONDON WC1R 4EJ, ENGLAND

7926020

YAZDANI-BAYGI, EBRAHIM
A HEURISTIC APPROACH FOR CLASS-SCHEDULING.

THE UNIVERSITY OF OKLAHOMA, PH.D., 1979

University
Microfilms
International

300 N. ZEEB ROAD, ANN ARBOR, MI 48106

THE UNIVERSITY OF OKLAHOMA
GRADUATE COLLEGE

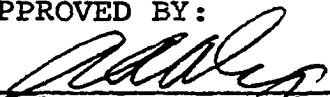
A HEURISTIC APPROACH FOR CLASS-SCHEDULING

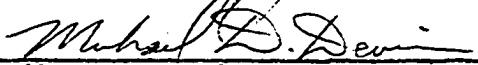
A DISSERTATION
SUBMITTED TO THE DEPARTMENT OF INDUSTRIAL ENGINEERING
in partial fulfillment of the requirements for the
degree of
DOCTOR OF PHILOSOPHY

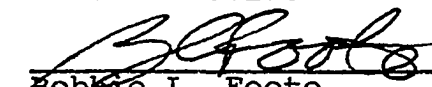
By
EBRAHIM YAZDANI-BAYGI
Norman, Oklahoma
May 1979

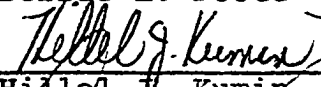
A HEURISTIC APPROACH FOR CLASS-SCHEDULING

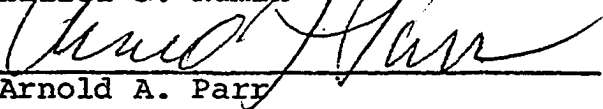
APPROVED BY:


Adel A. Aly, Chairman


Mike D. Devine


Bobbie L. Foote


Hillel J. Kumin


Arnold A. Parr

DISSERTATION COMMITTEE

ACKNOWLEDGEMENTS

I wish to express my appreciation to my committee chairman, Dr. Adel A. Aly, for his advice and encouragement during the course of my dissertation.

I am grateful to Dr. Mike Devine for his invaluable suggestions and other members of my committee, Dr. Bob Foote, Dr. Hillel Kumin and Dr. Arnold Parr for their encouragement throughout my doctoral program.

I also would like to express my very special thanks to my wife, Mandana, who, through her willingness to work and her moral support, has made my graduate work possible.

ABSTRACT

Scheduling problems are among the most challenging problems in applied mathematics. Solution procedures for most of these problems, however, are not among the most elegant.

Theoretical models have been constructed in some cases but the huge size of most real-life scheduling problems do not allow the solution procedures to these models to be computationally feasible or economical. In many cases, therefore, approximate solutions called heuristic solutions which generate "good" rather than optimum solutions, must be used. One such case is the class schedule construction.

In this research a theoretical model, based on network flow theory, is developed which lends itself to solution procedures that are computationally feasible. The model is for the case where a class schedule exists (such as a school's previous year class schedule) and the problem is to schedule students into their requested courses such that total number of students encountering conflicts is kept at a minimum.

A heuristic scheduling algorithm, composed of several algorithms for different types of courses, is developed which is suitable for traditional and modified school structures. The traditional school structure is considered as that in which each class meets once for one period of a scheduling cycle (normally a day). A modified school structure is similar to the traditional structure but allows for class meetings of more than one successive period.

The algorithms are implemented and tested on an actual data set for a school of complex, modified structure. The overall algorithm, after ranking the courses according to an index of difficulty processes one course or a group of courses at a time, using a special algorithm, in each case, based on the type of courses being processed.

The computational experience presented shows that:

- (a) The algorithm is quite efficient and therefore economical.
- (b) The algorithm is effective in class scheduling for traditional and modified school structures. In the case of a complex modified school structure the algorithm reduced the number of students encountering conflicts to less than one-half compared to a manually constructed schedule.

TABLE OF CONTENTS

	Page
ACKNOWLEDGEMENTS	iii
ABSTRACT	iv
 Chapter	
1. INTRODUCTION	1
2. LITERATURE REVIEW	5
2.1. Exact Mathematical Methods	5
2.2. Heuristic Approaches	20
3. A NETWORK FLOW MODEL FOR SCHEDULING STUDENTS INTO AN EXISTING CLASS SCHEDULE	30
3.1. Network Flow Model Formulation	31
3.2. Informal Description of the Network Model Formulation	32
3.3. Network Flow Problem Construction	42
3.4. Model Flexibility in More General Cases	45
3.5. Model Application	46
3.6. Computational Methods	47
3.7. Summary	48
4. A HEURISTIC APPROACH FOR CLASS-SCHEDULING IN TRADITIONAL SCHOOLS	49
4.1. General Problem Description	50
4.1.1. Basic Assumptions	51
4.1.2. Heuristics	51
4.2. The Heuristic Solution Approach	53
4.2.1. Treatment of Locked-In Courses	56
4.2.2. Scheduling Several-Hour Courses	56
4.2.3. Scheduling Single-Section Teacher or Room Restricted Courses	58

Chapter	Page
4.2.4. Scheduling Single-Section Unrestricted Courses (No Teacher or Room Restrictions) .	62
4.2.5. Scheduling Multi-Section Courses (2 to 5 Sections) . . .	64
3.2.6. Scheduling Multi-Section Courses (6 Sections or More) .	76
4.3. The Heuristic Algorithm	81
4.4. Implementation	98
4.4.1. Component 1: Program MASTERFILE	99
4.4.2. Component 2: Program SCHEDULE-1	102
4.4.3. Component 3: Program SCHEDULE-2	103
4.4.4. Component 4: Program SCHEDULE-3	104
4.4.5. Setting Up the Program and Data Decks	104
4.5. Computational Experience	105
4.6. Summary	110
5. CONCLUSIONS AND SUGGESTIONS FOR FURTHER RESEARCH	112
5.1. Conclusions	112
5.2. Suggestions for Further Research . . .	114
BIBLIOGRAPHY	117
APPENDICES	
I. TERMINOLOGY	123
II. MANUAL SCHEDULE, CENTRAL MID-HIGH SCHOOL . .	126
III. PROGRAM LISTINGS AND SAMPLE OUTPUTS	
Part A. Input Listings	
Part B. Output Listings	

SCHOOL SCHEDULING

CHAPTER I

INTRODUCTION

School scheduling is the process of building a time schedule of courses such that the number of student schedule conflicts are minimized, and the utilization of space, time, and teachers is optimized. It is assumed that students request the courses that they would like to take and based on these requests, the schedule, called a master schedule, is constructed. The master schedule is a time schedule of courses. When constructing a master schedule many factors must be considered, some of which include the following: There is a limited number of teachers, each having his own time limitations and class preferences. There is a limited number of classrooms, some special purpose classrooms such as laboratories and other general purpose rooms with different capacities. There is a limited number of time periods at which a large number of courses must be scheduled. Some courses may require more than one period per meeting. Some courses should be scheduled such that there is a minimum time lapse between two successive meetings of the course. A group

of courses should be scheduled at different time periods because of teachers or room constraints. There are courses which offer more than one section, each section having a limited number of seats available, etc. The above list of factors involved in scheduling is by no means complete. Other considerations may be important depending on the particular characteristics of the school.

There have been two general solution approaches in attacking a scheduling problem, mathematical or exact methods and heuristic or methods based on intuition and past experience.

The school scheduling problem, even in its general form, could be formulated mathematically. The solution to this mathematical formulation would give the best possible schedule for a given set of restrictions. It is not feasible, however, to solve such a problem by exact mathematical techniques with existing solution procedures and computers. Computer capabilities, however, have been exploited to develop general purpose scheduling programs. These programs are usually based on past experience, intuitive techniques, or mathematical techniques. Such scheduling programs generally are designed to perform scheduling for a wide range of schools. Since each school has its own type of constraints and limitations general purpose programs must consider every conceivable constraint, and limitation in order to be very general. However, such programs often can become highly

inefficient, uneconomical, and difficult to implement because of their complexity.

A more realistic approach would be to develop a relatively small program tailor made to consider some of the most important factors common in many local schools. Such a program could be designed small enough to be economical, flexible enough to be able to schedule a wide range of local schools, and incorporate several heuristics combined with mathematical optimization techniques to develop good schedules with few student conflicts.

As far as problem definition is concerned, three broad categories of practical scheduling or time-tabling problems could be identified:

- (a) First is the case where originally a master schedule exists and the problem concerns scheduling of students to time blocks.
- (b) Second is the case where a specified number of periods are given and classes and students must be scheduled to minimize the number of conflicts.
- (c) Third is the case where the assumption is that conflicts must not occur and it is tried to find a minimum number of periods.

In subsequent chapters possibilities are investigated in developing a model, applicable to category (a), for scheduling students into a given class-schedule in an optimal fashion.

It is also intended to develop a heuristic algorithm, applicable to category (b), for constructing a class-schedule, i.e., assigning a set of courses, which are requested by a group of students, to a set of time periods while scheduling students into the courses that they have requested. It is felt that heuristic and mathematical optimization techniques could be incorporated into the algorithm in order to develop class-schedules with relatively few student conflicts.

Specifically, optimization techniques would be used in an attempt to schedule a group of teacher or room restricted courses at the same time. An attempt would also be made to construct a model and an efficient technique for scheduling multi-section courses. Finally, it is desired to develop a computer program, based on the above-mentioned algorithm, specifically for schools with traditional and modified structures. This computer program will be implemented and tested using actual data and will be small as far as main memory is concerned such that it could also be implemented on small computer installations frequently available to many secondary schools.

CHAPTER II

LITERATURE REVIEW

2.1. Exact Mathematical Methods

A simplified but actual case of the problem of scheduling which is of the second category is described as follows:

There are N periods (non-overlapping) into which courses can be scheduled. A course can go into any of the N periods, and we can ignore teacher or classroom schedules. There are J courses, with some courses having only one section and others ranging up to five sections. There are no differences between several sections of a course. There are I students, and each student lists L courses he would like to take. The goal is to schedule courses so as to minimize the number of student "conflicts" that arise.

A more simple case of the above problem is when all courses have only one section which could be considered as a problem of scheduling activities among facilities in such a way to minimize the conflict cost. In this case the restrictions are:

1. Each activity (course) must be assigned to one and only one facility (time blocks).
2. Any number of activities could be assigned to a facility.
3. If two activities are assigned to a facility a cost would be incurred, such as number of conflicts when two courses are scheduled at the same time.
4. There is no conflict between facilities (time blocks do not overlap).

This problem could be formulated as a quadratic assignment problem as follows:

$$\begin{aligned} \text{Min } z &= \frac{1}{2} \sum_{j=1}^n \sum_{i=1}^m \sum_{k=1}^m a_{ik} x_{ij} x_{kj} \\ \text{s.t. } \sum_{j=1}^n x_{ij} &= 1 & i = 1, \dots, m \\ x_{ij} &= 0 \text{ or } 1 & \begin{matrix} i = 1, \dots, m \\ j = 1, \dots, n \end{matrix} \end{aligned}$$

where:

z = the total number of conflicts

m = number of courses

n = number of time periods (for nontrivial solution $m > n$).

a_{ik} = number of conflicts if course i and k are both assigned to time block j which equals the number of students requesting to take both courses i and k . a_{ii} is the cost of scheduling any course with itself, then $a_{ii} = 0$.

$$x_{ij} = \begin{cases} 1 & \text{if course } i \text{ is assigned to time block } j \\ 0 & \text{otherwise} \end{cases}$$

Since there are a finite number of feasible solutions total enumeration could be one way of finding an optimum solution. If N_i is the number of time periods to which course i can be assigned, the number of feasible solutions would be $\prod_{i=1}^n N_i$ where n is the number of courses. It is seen that for large n total enumeration is out of the question.

Except for relatively small sized problems, an exact solution to the quadratic assignment problem cannot be obtained at a reasonable computational cost. Exact procedures that have been developed so far are mostly branch and bound, or implicit enumeration. Francis and White [50] indicate that branch and bound procedures are normally unable to solve problems where number of variables are greater than 15. In the case of total enumeration they indicate that the time required to solve a problem with $n=12$ on an even new-generation computer would be infeasible. Heuristic solution procedures, therefore, are generally used to obtain "good" solutions. Work in both exact and heuristic areas has been done by Gilmore [1], Lawler [3], Broder [7], Hillier [46], Nugent, Vollmann and Rum1 [47], Pierce and Crowston [48], Hanan and Kurtzberg [49], Geoffrion and Graves [51].

Considering the formulation given above it was shown

by Carlson and Nemhauser [5] that (a) the integer constraints are not required; (b) because of the special structure of the model, an algorithm, considerably more efficient than existing quadratic programming algorithms, is possible; (c) since the cost function is not convex only local minima are obtainable.

Lawler [3] shows that the quadratic assignment problem is equivalent to an integer linear program as follows.

Suppose the quadratic assignment problem is

$$\begin{aligned}
 \text{Min} \quad & \sum_{ij} \sum_{pq} c_{ijpq} x_{ij} x_{pq} \\
 \text{S.T.} \quad & \sum_j x_{ij} = 1 & i = 1, \dots, n \\
 & \sum_i x_{ij} = 1 & j = 1, \dots, n \\
 & x_{ij} = 0, 1 & i, j = 1, \dots, n
 \end{aligned}$$

Then assuming

$$y_{ijpq} = x_{ij} x_{pq}$$

The problem becomes

$$\begin{aligned}
 \text{Min} \quad & \sum_{ijpq} c_{ijpq} y_{ijpq} \\
 \text{S.T.} \quad & \sum_j x_{ij} = 1 & i = 1, \dots, n \\
 & \sum_i x_{ij} = 1 & j = 1, \dots, n \\
 & \sum_{ijpq} y_{ijpq} = n^2 \\
 & x_{ij} + x_{pq} - 2y_{ijpq} \geq 0 & i, j, p, q = 1, 2, \dots, n \\
 & x_{ij} = 0 \text{ or } 1 & i, j = 1, \dots, n \\
 & y_{ijpq} = 0 \text{ or } 1 & i, j, p, q = 1, 2, \dots, n
 \end{aligned}$$

It is seen that the quadratic assignment problem has n^2 variables but the integer linear problem has $n^4 + n^2$ variables. This shows that transforming the problem from quadratic form to integer linear form is not computationally helpful, considering the present state of the integer programming algorithms. Lawler states that "the quadratic assignment problem is a finite problem. In principle, each of the $n!$ feasible solutions could be evaluated to find the minimum. Alternatively, existing integer programming techniques could be applied to the equivalent linear problem. However, neither of these alternatives is attractive."

Davis, Devine, and Lutz [7] take a problem such as the above and further assume that two facilities operations could be in conflict (overlapping time blocks). Their formulation of the problem is the same as the quadratic assignment problem above, in addition, they introduce a coefficient t_{pq} as a coefficient in the objective function as follows:

$$t_{pq} = \begin{cases} 1 & \text{if the operation of facility } p \text{ and facility } q \\ & \text{is in conflict, therefore } t_{pp}=1 \text{ for } p=1, \dots, m \\ 0 & \text{if the operation of the facility } p \text{ and facility } \\ & q \text{ is not in conflict, note that } t_{pq} = t_{qp}. \end{cases}$$

They develop necessary optimality conditions and then develop an algorithm (an extension of the algorithm developed in [5]) that converges to a local optimum or to a point that satisfies necessary conditions for optimality. Since the computation time of major steps of the algorithm is related linearly (but

not exponentially) to the number of variables, the algorithm is computationally feasible for large problems. Murikipudi [45] compares computational results using this algorithm called "LO-OPT" versus two other heuristic algorithms. One of the latter two algorithms is developed by Murikipudi himself and is called "PAM", short for "Preference Allotment Method." The other algorithm is called "ONE-OPT" and developed by Devine, Kumin and Lutz [44] for scheduling the one-section classes in the Oklahoma City public schools. The three algorithms are all designed to be used for single-section scheduling problems when conflicting time blocks are present. Murikipudi considers 15 problems where number of classes range between 8 and 40 and number of time blocks are between 4 and 15. His results indicate that LO-OPT generally gives better results. Once a solution of a problem reached a local optimum, however, LO-OPT cannot improve the solution further. The size of the computer program for obtaining solutions using LO-OPT is relatively large and therefore not very attractive for implementation on small computers. The problem of scheduling one-section courses to minimize number of conflicts is the same as the problem of scheduling final examinations. Efforts have been made to approach the problem of an examination schedule which requires a minimum number of periods and satisfies the constraint that no student be required to take two exams simultaneously. Example of work in this area is Hall and Acton [8]. Basically, they first

group together, in all possible ways, those courses which may have their examinations scheduled at the same period. Then a minimum number of these groups are selected such that each course is included at least once. Finally, they obtain a minimum schedule by removing multiple occurrences of courses and scheduling each group at a different period.

Another approach to solving one-section problem might be to find a way to decompose the problem into several smaller independent problems and solve these problems independently. It is noticed that the constraints could be decomposed by the course number. The question is whether the objective function could also be decomposed? Since students take courses basically in their own area of interest it seems logical that the conflict matrix may be changed to have the arrangement of blocks of course conflicts along the diagonal. Weil and Kettler [9] propose a rearrangement for the linear programming problem constraints matrix. They suggest that if a large linear programming problem has a suitable block-angular structure, one of a time-saving decomposition algorithms could be used. Another idea in this respect is the analogy between identification of subproblems in one-section problem and isolating the disconnected subgraphs of a directed graph. It seems reasonable that by using some ideas from theories of directed graphs the problem could be subdivided into several smaller problems and solved independently.

Another avenue of study in school scheduling problems has been teacher-class-time schedules. In this kind of

problem conflicts are not allowed and it is desired to find at least one possible schedule. Gotlieb [10] is among those who has done extensive work in this area. In his approach to the problem he assumes that the weekly requirements have been subdivided into five sets of daily requirements. This problem of dividing the weekly requirements into a set of daily requirements can be formulated in the same terms as the problem of finding the time-table to satisfy the daily requirements. Also, to meet these pattern requirements, to provide for hours when some teachers or classes may not be available, and to ensure from the beginning that certain critical facilities (e.g., lunch rooms, science laboratories and gymnasias) are fully utilized, some part of the time table is filled out in advance and the problem is one of completing a partially filled time-table. It is these initial assignments which make the problem difficult. The procedure used by Gotlieb is that the problem is first checked for feasibility in the light of the requirements and initial conditions. This feasibility test is carried out by systematically following through all the implications of the initial assignments. Once feasibility is assured the time-table can easily be constructed and, in fact, all solutions satisfying the initial conditions can be found.

Dempster [11] formulates the class-teacher-time problem as a quadratic assignment problem. His algorithms may be regarded as 3-dimensional generalizations of familiar

algorithms for the assignment problem. Matrix reduction techniques have been applied to time-table problems also. In this regard work of Lions [12] could be mentioned. In his article, he describes the application of Kuhn's Hungarian Method to the problem of matrix reduction as needed in Gotlieb's method for time-table generation. The method is suited to both hand and computer calculation. Devices to improve the efficiency of the basic algorithm are discussed. Many other researchers have worked in this area such as Almond [13], Berghuis [14], and Barraclough [15].

In the one-section quadratic assignment problem other constraints could be included to take care of particular situations. One example could be the case where a set of courses are supposed to meet in a certain class and then the restriction is that those courses should not be assigned to the same or conflicting time blocks. The constraint to be added to the previous set of constraints is

$$\sum_{i \in p} x_{ij} \leq 1 \quad j = 1, \dots, n$$

p is the set of all courses which
are supposed to use a particular class.

Other sorts of constraints such as the above (the same teacher for several courses) could be added to the problem.

A more general and realistic scheduling problem is the case where courses could have more than one section. These types of courses are called multi-section courses.

It is assumed that students request a multi-section course itself and not a particular section of the course, and therefore, there is the freedom to assign them arbitrarily to any of the sections of the course. Since, otherwise, if the students choose a section each section could be considered as a one-section course and as a result the problem becomes a one-section assignment problem which could be solved using one of the available algorithms for one-section scheduling. When there are courses which offer more than one section and there is the freedom of assigning the students to any of the sections, then the students individually should be considered because the decision must be made as to which student should be assigned to which section in order to avoid a conflict (or improve the solution). Also, if there is a limit on the number of students who can take a particular section of a course, the students must be scheduled into sections such that there is a balance between the sections at the end, and therefore, attention should be given to individual students. Accordingly, let us consider another formulation of the problem similar to a general formulation developed by Oakford et al. [34], in which the objective is to maximize the number of students' requests satisfied rather than minimizing the number of conflicts. Consider the following formulation for one-section problem when there are m courses, n students, and ℓ time periods.

$$\text{Max } z_0 = \sum_{i=1}^n \sum_{j \in E_i} \sum_{k=1}^l x(i,j,k)$$

$$\text{S.T. } \sum_k y(j,k) = 1 \quad j = 1, \dots, m$$

$$x(i,j,k) \leq y(j,k) \quad \begin{array}{l} i = 1, \dots, n \\ j \in E_i \\ k = 1, \dots, l \end{array}$$

$$\sum_{j \in E_i} x(i,j,k) \leq 1 \quad \begin{array}{l} i = 1, \dots, n \\ k = 1, \dots, l \end{array}$$

where:

$$x(i,j,k) = \begin{cases} 1 & \text{if student } i \text{ is taking course } j \text{ at time } k \\ 0 & \text{otherwise} \end{cases}$$

$$y(j,k) = \begin{cases} 1 & \text{if course } j \text{ is scheduled at time } k \\ 0 & \text{otherwise} \end{cases}$$

$$E_i = \text{the set of all courses requested by student } i$$

In the above formulation the objective function maximizes the number of students' requests satisfied.

The first constraint ensures that course j is scheduled at a period k and is scheduled only once. The second constraint ensures that student i could be scheduled into course j at time period k only if course j is scheduled at time k . The third constraint says that student i can take only one of the courses that he is requesting at a period k .

When there are courses which offer more than one section, then the problem could be formulated as follows:

$$\text{Max } z_0 = \sum_{i=1}^n \sum_{j \in E_i} \sum_{s \in S_j} \sum_{k=1}^l x(i,j,s,k)$$

$$\begin{aligned} \text{S.T. } \sum_k y(j,s,k) &= 1 & j &= 1, \dots, m \\ & & s &= 1, \dots, S_j \\ \sum_{s \in S_j} x(i,j,s,k) &\leq 1 & i &\in P_j \\ & & j &\in E_i \\ & & k &= 1, \dots, l \\ x(i,j,s,k) &\leq y(j,s,k) & i &= 1, \dots, m \\ & & j &\in E_i \\ & & s &= 1, \dots, S_j \\ & & k &= 1, \dots, l \\ \sum_{j \in E_i} \sum_{s \in S_j} x(i,j,s,k) &\leq 1 & i &= 1, \dots, n \\ & & k &= 1, \dots, l \end{aligned}$$

where:

$$\begin{aligned} x(i,j,s,k) &= \begin{cases} 1 & \text{if student } i \text{ is scheduled into section } s \text{ of course } j \text{ at period } k \\ 0 & \text{otherwise} \end{cases} \\ y(j,s,k) &= \begin{cases} 1 & \text{if section } s \text{ of course } j \text{ is scheduled at time period } k \\ 0 & \text{otherwise} \end{cases} \\ S_j &= \text{number of sections of course } j \\ P_j &= \text{the set of all the students requesting course } j \\ E_i &= \text{the set of all the courses that student } i \text{ is requesting} \end{aligned}$$

If there is a restriction on the number of students who can be scheduled into a section, a constraint such as the following could be added to the above sets of constraints:

$$\sum_{i \in P_j} x(i,j,s,k) = N_{js} \quad \begin{aligned} j &= 1, \dots, m \\ s &= 1, \dots, S_j \\ k &= 1, \dots, l \end{aligned}$$

where N_{js} is the limit on the number of students who could be scheduled into section s of course j .

If it is required that no more than one section of a course be scheduled at a time period, a constraint such as the following could be added.

$$\sum_{s \in S_j} y(j,s,k) = 1 \quad \begin{array}{l} j = 1, \dots, m \\ k = 1, \dots, \ell \end{array}$$

This model is very general and almost any other sort of restrictions could be incorporated in it (teacher restrictions, class restrictions, etc.).

Therefore, the problem, even in much more general cases than this could be formulated as a mathematical model. Now the question is that whether a solution could easily be obtained?

For a medium size school of 700 students and 100 courses assuming each student on the average requests 6 courses, each course on the average offers 4 sections, and 6 available time periods, then:

Number of $X(i,j,s,k)$ variables = $700 * 6 * 4 * 6 = 100,800$

Number of $Y(j,s,k)$ variables = $100 * 4 * 6 = 2,400$

Total number of variables = 103,200

Number of constraints of the first type = $100 * 4 = 400$

Number of constraints of the second type = $700 * 6 * 6 = 25,200$

Number of constraints of the third type = $700 * 6 * 4 * 6 = 100,800$

Number of constraints of the fourth type = $700 * 6 = 4,200$

Total number of constraints = 130,600

It is seen that for a medium size school total number of (0-1) variables and total number of constraints are both over 100,000. This problem belongs to a class of problems called NP-complete by Karp [58] and the difficulty is not in mathematically formulating the problem but in the computation times required by existing solution methods.

Existing literature about zero-one integer programming problems shows that only problems with small number of variables could be solved by exact algorithms. A good survey article in this area is the one by Geoffrion and Marsten [19]. They discuss computational results for a number of general purpose algorithms for integer linear programming problems. They indicate that using enumerative algorithms large size hospital scheduling problems up to 100 0-1 variables have been solved. Additional computerized experience is presented solving problems with 200 binary variables although strict optimum solutions were not obtained. Geoffrion and Marsten claim that computational experience with cutting plane algorithms has been confined to very small problems. For crew scheduling problems, however, they indicate that large size problems in order of 7,000 binary variables and 150 constraints can often be solved on the order of 40 minutes of CDC 3600 time. Garfinkel and Nemhauser [52] report that scheduling problems of the size of 40 constraints and 131 variables have been solved using enumerative algorithms. For cutting plane algorithms they state that computer programs

have been developed which can handle up to 60 constraints and 240 variables. On ILPC-based (Integer Linear Programs over Cones) they report that the largest problem solved had 104 constraints and 236 variables. For knapsack problems they state that problems with 5000 variables were solved. In the case of traveling salesman problem they report the solution of a 42-city problem which has 861 binary variables and 84 assignment constraints. Karney and Klingman [53] report that using an in-core, out-of-core primal network code with an imbedded branch-and-bound code they solved a plant location problem with 3300 nodes, 41,000 arcs in 40 minutes of CDC 6600.

Theoretical developments in recent years, reviewed by Steen [32], indicate that many of the $(0,1)$ integer linear programming problems have a history of unsuccessful searches for optimal solution methods where the computation times grow polynomially rather than exponentially with size of data and many of these problems can be put into a category which can be transformed into a mathematical logic problem called the satisfiability problem. This could strongly indicate the infeasibility of determining optimal solutions which may be computationally feasible. It is concluded by Steen that searches for sufficiently good solutions are now much more common than search for exact algorithms.

Other references in this area are Glover and Klingman [38], Geoffrion [54], Shih [55], Talbot and Patterson [56], and Zanakakis [57].

2.2. Heuristic Approaches

It is seen that exact techniques are theoretically applicable to school scheduling problems, however, the relatively large size of any typical school schedule makes the use of techniques impossible. The mathematical approaches would be used more and more in the near future to aid in solving the scheduling problem. Only the heuristic approach has so far had any real success. For example, Holz [2] has developed a program which has been successfully used for high schools and small colleges. This program is called GASP (Generalized Academic Simulation Programs) and is divided into two segments. In the first a time-table is constructed and includes the tasks of making assignments of time, classroom and instructors to each class. In this segment student requests are taken into account, but students are not definitely assigned to classes. In the second segment, students are assigned to classes designated in the time table. The segments may be run separately or together on one continuous run. The second segment of the program "registers" students into the classes of the time-table. Assignment of students to classes is done in such a way as to make possible a general consideration of student preferences, such as time distribution of classes or choice of instructor. To do this, the program can compute a few (10 to 25) workable schedules for each student and, on the basis of any special rules or measurements the school desires, pick the best of these schedules. The various schedules are computed so that they are likely to

be quite different from each other, and section sizes remain balanced with high probability. In most instances, the man-hours needed to produce a schedule using GASP is near 25 percent of the time required by traditional methods. The computer time necessary to schedule a school of about 1,000 students is between two and four hours of 7094 (IBM) use and perhaps 30 hours of 1401 (IBM). One of the first workable programs of school scheduling was developed at Purdue University in 1956. Shortly thereafter, International Business Machines developed a scheduling program called CLASS (Class Load and Student Scheduling) [21] based on some of the early work done at Purdue. The CLASS program assigns students to a given master schedule on the following basic priority.

1. The student must get what he asks for.
2. He should not be given a schedule with too many sequential study halls (determined by the school).
3. He will be scheduled into sections in a fashion that tends to maintain a balance in class size.

The way the program CLASS assigns students to courses is that the student is tried in those sections of courses he is requesting that have the most seats available at that given instant. The computer begins this operation for each individual student with those courses requested that have the fewest sections offered and progress to those with the most. In this way one-section courses are scheduled first--as they usually are in manual procedures. The CLASS program was used to schedule 165,000 students in the Cleveland, Ohio-Pittsburgh, Pennsylvania

area during one summer. The average number of computer passes made by these schools involving the 165,000 students was two and one-half to achieve scheduling percentage of 95 percent. Computer costs ran approximately 35 cents per student.

Another program taking a different approach from others is called SOCRATES (Scheduling of Classes Realized Automatically Through Effortless Systemization) developed by Wilkes [22]. SOCRATES takes the approach that all students should first be scheduled into those classes which can be taken at only one time during the day. By scheduling initially only "one-chance" courses for all students and passing over all courses that have two chances during the day, a basis of priority can be established. Later, when the time comes that there is no remaining student who has a priority course request, the first two chance course is scheduled. Scheduling the latter type of course could result in the creation of a new one-chance course; therefore, after scheduling a two-chance course for a student, a new attempt is made to schedule one-chance courses for him. Only when there are no single-chance courses left unscheduled is the next two-chance course scheduled. When there are neither single- nor double-chance courses remaining unscheduled, then the first three chance course is scheduled, again returning to the one-chance mode, then the two-chance mode, then finally returning to schedule three-chance courses, and so on.

An algorithm developed by Anderson [23] assumes that a master schedule exists and is already available to the

computer. A file of student course requests is translated into a file of actual class assignments very rapidly. This speed permits repeated use of the computer programs, with refinements to the master schedule based on experience with the previous trial. In this context, the assignment algorithm can be used to develop the final master schedule. Anderson's approach tries many alternative possibilities to fit each student into a predetermined master schedule. If after a reasonable number of trials the student cannot be assigned to sections, he is referred to the scheduling staff as an exception. The computer provides a diagnosis of the conflict which frustrated its efforts, but human intervention and analysis is still required.

A system was developed at Stanford University by Oakford et al. [24] called Stanford School Scheduling System (SSSS). They found that the construction of the school schedule was a formidable obstacle. High school administrators are able to construct usable schedules only when virtually all classes conformed to the traditional meeting pattern of one period per day, five days per week. Oakford et al. claim that the scheduling problem in its most general form could be formulated as a linear integer programming problem. The system that they have developed is very general and is supposed to be adaptable to any kind of situation. In solving the scheduling problem they use a heuristic approach as follows:

Schedule (or deschedule and reschedule) sections

successively until a situation has been reached where all sections are scheduled and no increase in the objective function could result from descheduling and rescheduling any section. At each step of the scheduling process, the scheduler takes the action that appears to be reasonable in the light of the information available to him at that step. The action at a given step must not conflict with the action taken previously. At a given step in the scheduling process a section should be scheduled (or descheduled and rescheduled) only if that action promises to increase the value of the objective function. A general description of one heuristic approach is given where a set S_c of students that have requested enrollment in course x of which c is a section is partitioned into subsets relevant to the computation of the value of scheduling section c at period p . Also the cost of descheduling section c at period p^* is computed. If the value of scheduling section c at period p is greater than the cost of descheduling section c at period p^* then deschedule section c from period p^* and schedule it at period p . Schedules have been made for schools ranging from 120 to 2600 students. The effectiveness of the schedules, as measured by course requests satisfied, have ranged from acceptable (satisfy at least 90 percent of the requests) to excellent (satisfy at least 98 percent of the requests). The past experience indicates that schedules can be computed at a cost that is acceptable to the schools.

A revised version of the CLASS program called UPDATE has been developed at the Iowa Educational Information Center IEIC of the University of Iowa as reported by Van Dusseldorp and Richardson [25]. The IEIC used what they call a "daisy chain" approach in which students are first passed through the scheduling program in any order they happen to come and are assigned temporary schedules. The results of this first scheduling operation are stored, and then each student is completely rescheduled. He is taken out of the classes into which he was first scheduled and is scheduled again. This method results in the improvement of both class-size balance and individual student schedules.

Washington State University has a program to help scheduling as reported by Faulkner [26]. Some other articles in this area are such as Busam [27], Macon and Walker [28], Lutz, et al. [29], and Holzman [35].

Salazar [37] has developed an approximate interactive algorithm for constructing schedules for traditional high school structures. The traditional structure is characterized by special requirements such as: each class meets once for one period during a scheduling cycle. Salazar indicates that by authoritative estimates over 90 percent of secondary schools in the United States employ a traditional structure.

The algorithm developed by Salazar is based on a integer linear programming formulation, in which after ranking the student file, processes one student at a time

by formulating and solving a linear assignment problem.

If the school structure requires multi-period classes the model does not easily lend itself to a simple solution procedure.

In Salazar's algorithm students are first ranked based on the courses that they are requesting. The ranking is done by finding the sum of reciprocals of the number of sections of all the courses requested by a student. As an example, if a student is requesting 6 courses having 1, 1, 2, 3, 4, and 4 sections, then his or her rank is $(1/1 + 1/1 + 1/2 + 1/3 + 1/4 + 1/4 = 3.333)$. Scheduling is done by considering one student at a time according to the ranking scheme. This means that the original scheduling problem is partitioned into small portions simply by considering one student at a time and the interaction between these portions are considered by selecting the order in which the portions are solved, i.e., the ranking scheme. Courses which are restricted to a certain time period or require more than one time period cannot be handled by the algorithm and therefore must be scheduled manually. When a student is considered for scheduling, the courses requested by him or her are identified. For each course that the student is requesting a section(s) is considered to be scheduled at time period $\underline{P}$ ($P = \text{no. of time periods, } P = 1, \dots, \underline{P}$). The following cases could happen:

1. The section is already scheduled at this time period.

A cost would be associated with scheduling the student into this section which is proportional to enrollment in the section divided by average number of students for a section of this course (average number of students for a section = total number of students requesting the course/number of sections).

2. The section is scheduled at another time period, or it is not feasible to schedule the section at this time period because another section is scheduled at this time period which must not be scheduled at the same time with this present section. A very high cost is associated with scheduling the section at this time period, represented by u_m .
3. It is feasible to schedule the section at time period $\underline{p}$. In this case the cost is the sum of several costs as follows:
 - (a) Section s could be restricted by a set of other sections say s_j (for example, s and s_j 's are taught by the same teacher) and therefore these sections must not be scheduled at the same time period. Suppose prior to scheduling s , s_1 of s_j sections could be scheduled at $\underline{p}$ and after scheduling s at $\underline{p}$, s_2 of these sections could be scheduled at $\underline{p}$ then $s_1 - s_2$ represents the decrease in the number of feasible possibilities of other sections not yet scheduled if s is scheduled at $\underline{p}$. The cost is proportional to $s_1 - s_2$.

- (b) Some of the sections in the set s_j may become critical in the sense that only one time period would remain at which they could feasibly be scheduled. There is a cost associated with creating such sections because of scheduling s at $\underline{P}$.
- (c) Some of the sections in the set s_j could become impossible to schedule because no other time period could be found at which their section could feasibly be scheduled after s is scheduled at $\underline{P}$. There is a cost associated with creating such sections because of scheduling s at $\underline{P}$.
- (d) A cost proportional to student enrollment at time period $\underline{P}$.

The total cost is the sum of the costs in (a), (b), (c), and (d).

The cost is computed as discussed above for all sections of a course that the student is requesting and the section with minimum cost selected. The procedure is repeated for all time periods and all courses that the student is requesting. An assignment problem is then set up where sections are assigned to time periods at a minimum cost.

All the heuristic approaches discussed here are general purpose approaches and their corresponding computer programs are large in size except for the last approach which is designed for traditional schools. In the last approach, even

though the corresponding computer implementation is suitable for small traditional schools, the algorithm itself is inflexible in the sense that it is not designed to consider locked-in courses, several-hour courses, and overlapping time blocks in general.

CHAPTER III

A NETWORK FLOW MODEL FOR SCHEDULING STUDENTS INTO AN EXISTING MASTER SCHEDULE

One type of scheduling problem is the case where originally a class schedule exists, i.e., courses and sections of courses are already assigned to time periods, and the problem concerns scheduling of students into the courses that they have requested.

Such a problem could arise for example at the beginning of a school year, when a school official tries to perform his class scheduling by assigning courses to time periods and then scheduling students into the courses that they have requested. The school official could use some intuitive method, his previous year's schedule, or some other procedure for assigning courses to time periods. Once courses are assigned to time periods, the problem is to schedule students into the courses that they have requested such that the number of students encountering conflicts is minimized. In this chapter a model is presented, based on network flow theory, which could be used to schedule students into a given class schedule in an optimal fashion.

3.1. Network Flow Model Formulation

To find an optimum solution to the problem described above the problem is partitioned into smaller subproblems. A feasible solution for each small subproblem is found and the sum of all the solutions of all small subproblems constitutes the overall optimum solution. The partitioning considered is one student at a time, namely, the small problem of scheduling a student into the courses that he or she is requesting. For each student, based on the courses that the student is requesting and the time periods at which these courses are scheduled a small network flow problem is constructed. The solution to this network problem, if one exists, would be a feasible schedule for the corresponding student. Those students for which a solution to their corresponding network flow problem does not exist are the ones which could not have a feasible schedule in the given master schedule and, therefore, would encounter a conflict. The total number of conflicts finally would be equal to the total number of infeasible solutions obtained by solving network flow problems for all students. This would be the minimum number of conflicts achievable for the given master schedule. Any improvements in order to reduce the number of conflicts at this point could be achieved by sensitivity analysis.

3.2. Informal Description of the Network Model Formulation

To show how the original problem of scheduling students into a given class schedule is partitioned into smaller subproblems, in the form of network flow problems, an example would be presented first. Let us consider a master schedule consisting of courses 1, 2, ..., 9, and the courses are assigned to time blocks as follows:

Time Blocks				
	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>
	1	2	3	4
Courses	5-1	6-1	5-2	6-2
	7-1	8-1	8-2	7-2
	9-1		9-2	9-3

The first row labeled "time blocks" represents time blocks 1, 2, 3, and 4. The second row indicates that course 1 is scheduled at time period 1, course 2 is scheduled at time period 2 and so on. The third row indicates that course 5, section 1, is scheduled at time period 1; course 5, section 2 is scheduled at time period 3, and so forth.

Suppose a student is requesting courses 2, 5, 8, and 9. To set up a network flow problem for this student let's consider the above table (existing master schedule) in the following form.

Time Blocks

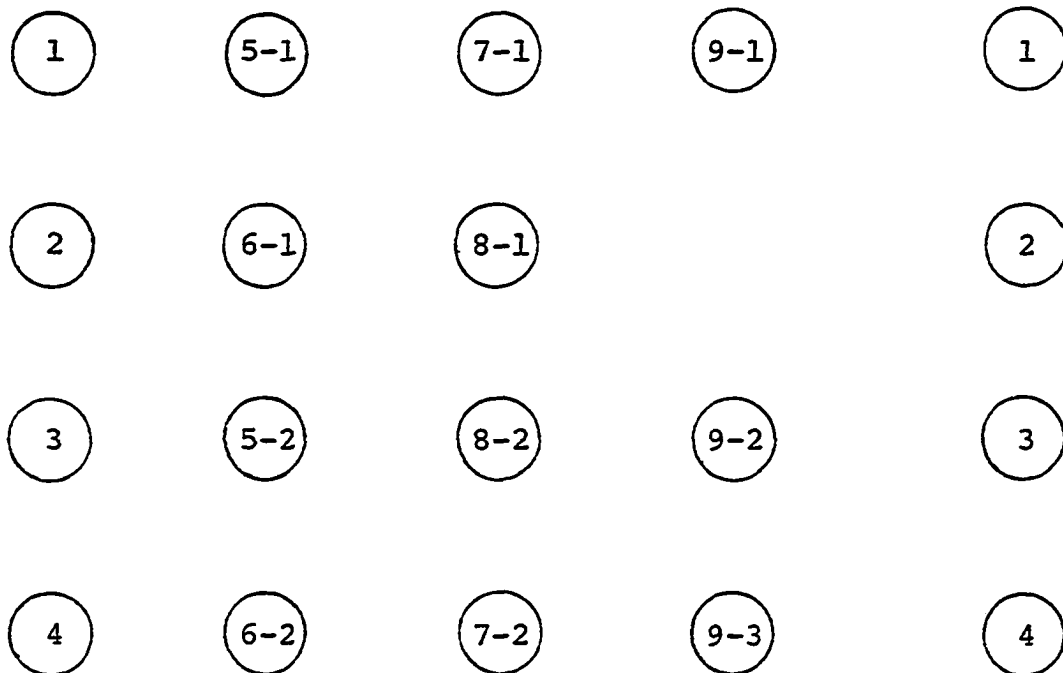


Figure 1.

In this representation a node represents a single section course or a section of a multi-section course and its location represents the time period to which it is assigned. A student would then be represented by a source node s and the arcs carry the flow to some of the nodes according to several restrictions. The restrictions ensure the existence of arcs between node s and the nodes corresponding to the courses that the student has requested and the arc has a lower and upper bound on flow as $(0-1)$. Each arc would then have a label with general form (L, U, C) where L and U

are lower and upper bounds on the flow, respectively, and C represents the cost of scheduling the student into a particular section.

To achieve a restriction for each multi-section course a super source would be considered which would receive only one unit of flow and therefore is able to send at most one unit of flow (Figure 2). This one unit can then flow into one of the nodes corresponding to different sections of a course and no flow into other nodes indicates that the student would not be scheduled into other sections of the same course.

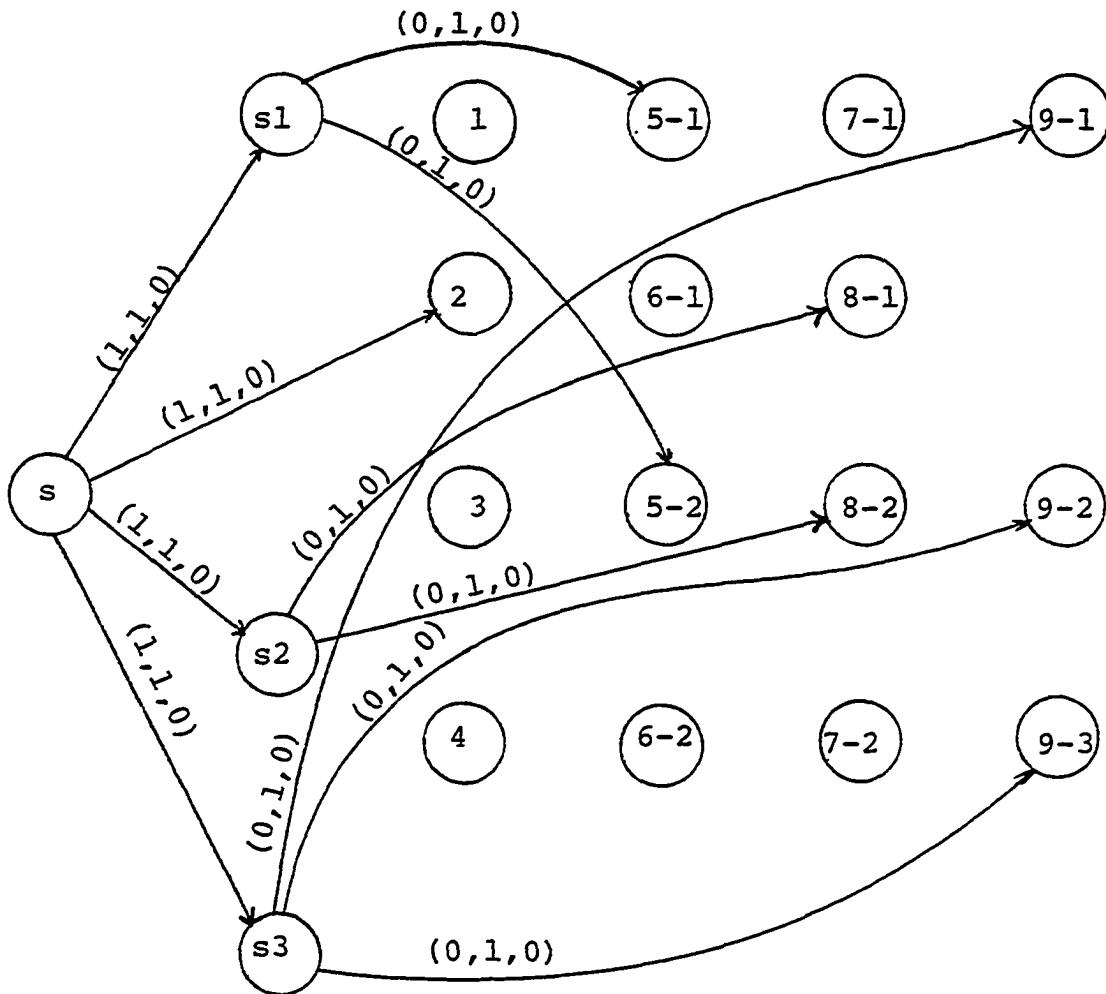


Figure 2.

In this figure node s is the supersource representing the student. Arcs (s,s_1) , $(s,2)$, (s,s_2) , and (s,s_3) all are of the form $(1,1,0)$. A $(1,1,0)$, for (s,s_1) means that one unit of flow must flow from s to s_1 . s_1 represents a supersource for course 5 which is a two-section course and therefore two arcs emanate from s_1 . Arcs $(s_1,5-1)$ and $(s_1,5-2)$ are of $(0,1,0)$ type which means that 0 or 1 unit of flow could flow from s_1 to either 5-1 or 5-2. Consider courses 5 and 9 and suppose one unit of flow moves from s to 5-1 and another unit of flow moves from s to 9-1 yielding a conflict. To prevent this from happening two nodes (supersinks) are added to the previous diagram (Figure 2) as shown in Figure 3 by nodes t_1 and t .

Flows of 0 or 1 unit could flow from 5-1 or 9-1 to t_1 but at most one unit could flow from t_1 to t .

The same line of reasoning would be applied to other courses and time periods producing the diagram in Figure 3.

The long arc from node t to node s with label $(4,4,0)$ is to make sure that 4 units of flow would move from t to s and puts this network flow problem into a form called a minimum cost circulation problem. The network is constructed in this form because the out-of-kilter algorithm developed by Fulkerson [39] deals with circulations and a circulation is an assignment of flow to arcs such that flow is conserved at each node. Totally there are 17 nodes and 24 arcs which form the network problem. The solution to this network flow

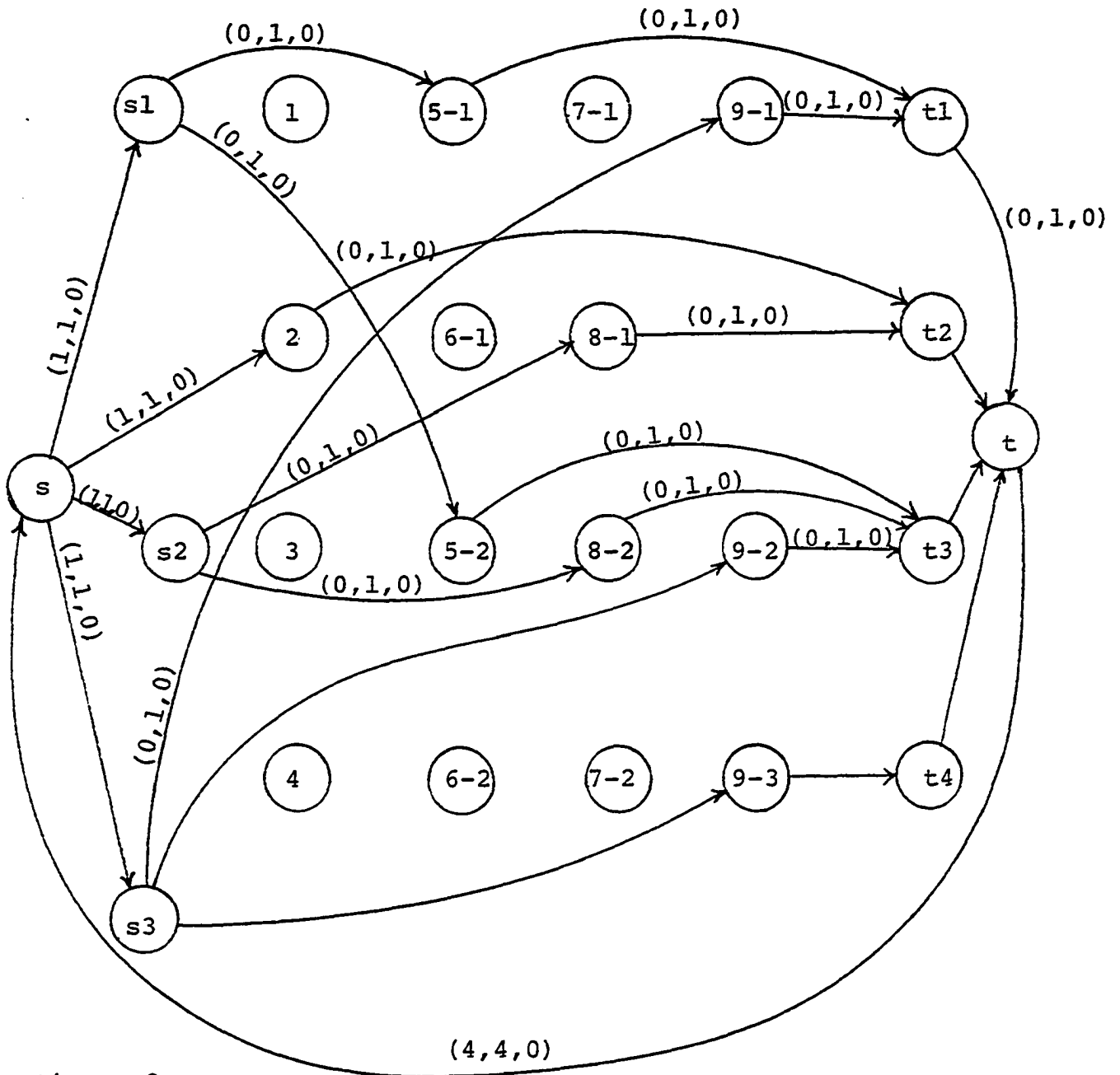


Figure 3.

problem, if one exists, would be a feasible schedule for the corresponding student. Two network flow problems such as above one with 12 nodes and 35 arcs and the other with 16 nodes and 35 arcs were solved using an out-of-kilter algorithm code. The two problems, network set up and solution, required .8 and 1.8 seconds of IBM

370/158 CPU time, respectively. A network flow problem, such as above, could be constructed for each student and then solved. All the feasible solutions represent all those students who could be successfully scheduled. All the infeasible solutions represent all those students for which a feasible schedule does not exist in the given class schedule. The overall solution, whether the objective is maximizing number of student requests satisfied or minimizing number of student conflicts is optimum for the existing class schedule. As to the size of each network flow problem consider the example used previously where a student is requesting six courses, a single-section, a 2-section, a 3-section, a 4-section, a 5-section, and a 6-section course. The network flow problem constructed for this student would have a maximum of 34 nodes and 54 arcs. It is therefore seen that, on the average, network flow problems constructed for students are relatively small both in terms of number of nodes and number of arcs.

At this point the general procedure for constructing a network flow problem for each student based on the assumption that a class schedule exists will be discussed. It is also assumed that time periods are non-overlapping and all the courses require one time period per meeting. Overlapping time periods and courses which require more than one time period per meeting are discussed in a later section.

Suppose a class schedule is given and a student's request list is selected to be considered for scheduling.

Consider each course and each section of a course that the student is requesting as a node in the network context. Assume there is a supersource node s (normally numbered 1) representing the student. This source node is capable of sending flow to other nodes and the amount of flow moving into and out of it must exactly equal the number of courses that the student is requesting. The network flow problem must be set up such that at most one unit of flow moves from the supersource to each course that the student is requesting. If the course is multi-section, then at most one unit of flow must move from the supersource s to all the nodes representing different sections of the course. This will ensure the restriction that the student is scheduled into at most one of the sections of a multi-section course. Units of flow move from a node to another through arcs connecting the two nodes and, therefore, the source node s is connected to all the nodes corresponding to the courses that the student is requesting through capacitated arcs.

The capacitated arcs all are directional and have labels of the form (L,U,C) , where L and U are lower and upper bounds on the amount of flow, respectively, and C represents the cost of utilizing the arc. For each single section course an arc is considered from s to a new node representing the course with a label $(0,1,0)$. This means that 0 or 1 unit of flow could move from s to the new node and the cost is 0. For each multi-section course an

intermediate supersource is considered. This intermediate supersource is connected to s through an arc with label $(1,1,0)$. This means that at least one ($L=1$) and at most ($U=1$) one unit of flow and therefore, exactly one unit of flow must flow from s to the intermediate supersource and the cost of utilizing this arc is 0. New nodes are then considered to represent different sections of the course, one node for each section. The intermediate supersource is connected to each of these new nodes by an arc labeled $(0,1,C_{j_m})$ where C_{j_m} is the number of students already scheduled in section m of course j . Consider the following example:

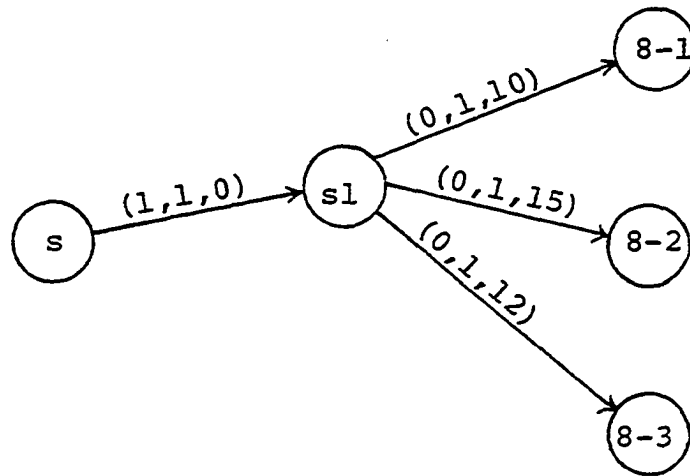


Figure 4.

Suppose course 8 is a three-section course requested by a student and each section has been assigned C_{j_m} students. If the existing class schedule and the courses that the student is requesting are such that one unit could flow from

s1 to either 8-1, 8-2, or 8-3, then that unit will flow into the arc with smallest cost and the student will be scheduled into the section with fewest number of students already scheduled in it. This is an attempt to achieve some balance on the number of students scheduled into different sections of a course.

New nodes are also considered for all the time periods at which a course or a section of a course requested by the student are scheduled. Arcs are then constructed from all the nodes corresponding to courses and sections of the courses to their corresponding time period nodes and they are all labeled $(0,1,0)$. For instance, in the previous example suppose course 8 section 1 is scheduled at time period 2, section 2 is scheduled at time period 4 and section 3 is scheduled at time period 5. A supersink node is considered and all the nodes corresponding to time periods are connected to this node in the following manner:

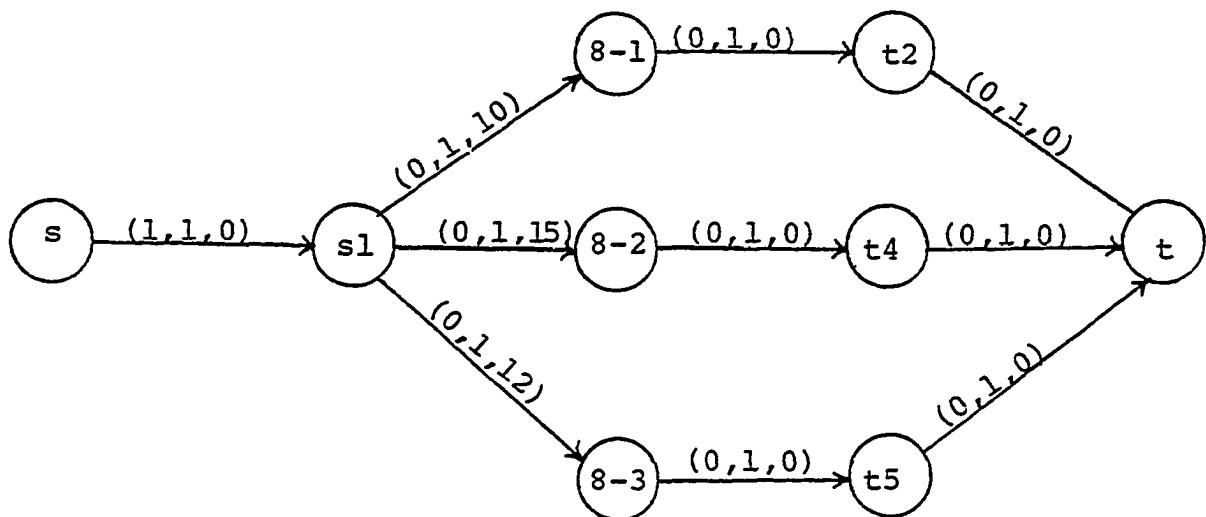


Figure 5.

Suppose this student, along with course 8, is also requesting a single section course, say course 7, scheduled at time period 4 and a double section course 9 scheduled such that section 1 is at time period 1 and section 2 is scheduled at time period 4. The network would be constructed as follows:

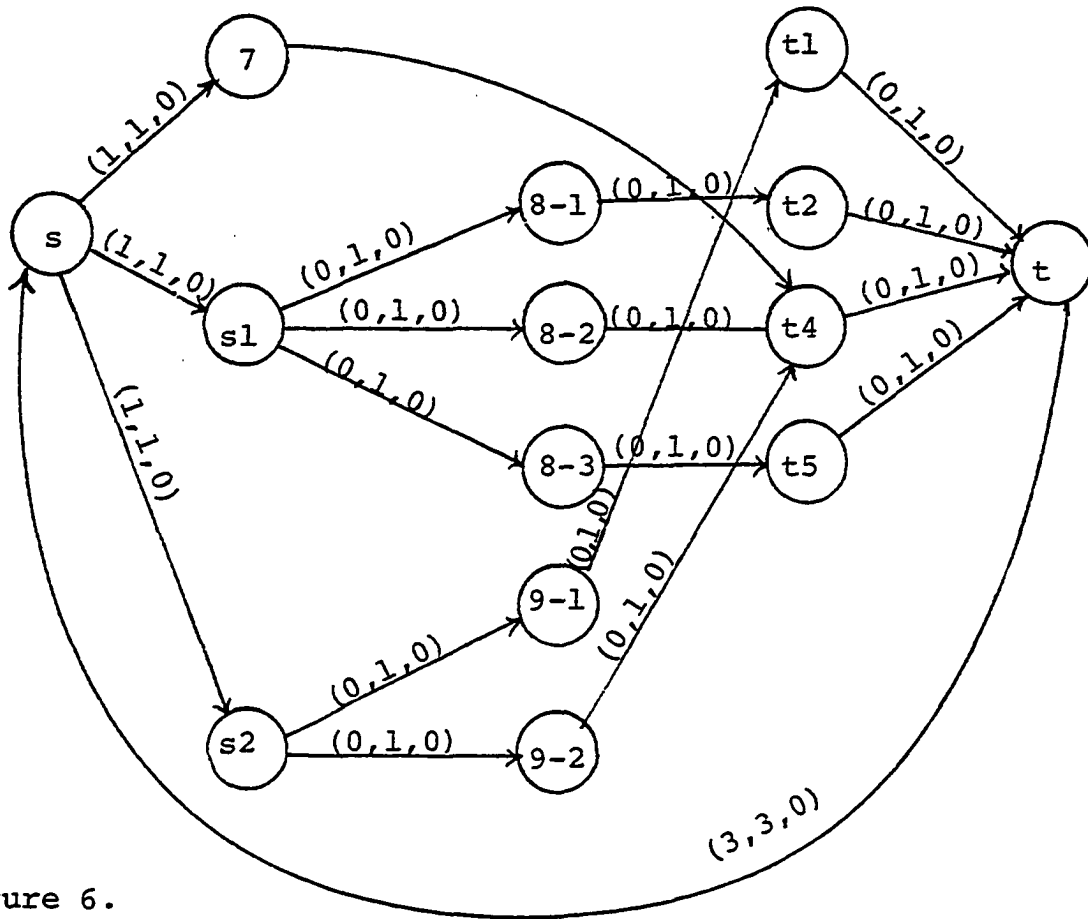


Figure 6.

A return arc labeled $(3, 3, 0)$ eventually ensures that exactly 3 units of flow (equal to the number of courses requested) moves from s to t if a feasible schedule exists.

3.3. Network Flow Problem Construction

NOS = Set of all students

N = Number of courses requested by a student

j = A course index, $j = 1, \dots, N$ set of courses
requested by a student

C_{j_m} = Number of students already scheduled into section
m of course j

S_j = Number of sections of course j

$NT = \sum_{j=1}^N S_j$ Total number of sections requested by a
student

$T(m, j)$ = Time period at which section m of course j is
scheduled

$TL(L)$ = A time period at which a section of a previously
considered course requested by student i is
scheduled. (Ex. when $j = 2$, the second course
requested by the student is considered, and
 $TL(1) = 3$, $TL(2) = 5$, ($TL(L) = 0$, $L = 3$, NT)
means that section 1 of the first course requested
by the student is scheduled at time period 3 and
section 2 of the first course is scheduled at
time period 5.)

K = A node index

K_s = A node index for multi-section courses

STEP 1: Select student $i \in NOS$. If NOS is empty STOP.

STEP 2: Compute number of nodes necessary to represent
supersources and all the courses and sections of

the courses requested by the student.

(a) Set $NN = N$ and $j = 0$

(b) Set $j = j + 1$

If $j > N$ Go to STEP 3

If $S_j \leq 1$ Go to STEP 2-(b)

Otherwise, set $NN = NN + S_j$, Go to STEP 2-(b)

STEP 3: Set $j = 0$ ($TL(L) = 0$, $L = 1$, NT), $t = 0$, $m = 1$

$K = 1$ (set up the first node; a supersource)

STEP 4: Set $j = j + 1$

If $j > N$ Go to STEP 7

If $S_j > 1$ Go to STEP 6

STEP 5: Set up node $K = K + 1$ (to represent the single section course j). Connect node 1 to node k to form arc $(1,k)$ and label the arc $(0,1,0)$.

(a) If there is an L ($L = 1, \dots, NT$) for which

$T(m,j) = TL(L)$ then a node $NN + L$ has been

already set up. Connect node k to node $P = NN + L$

to form arc (k,p) and label the arc $(0,1,0)$,

go to STEP 4.

(b) If $T(m,j) \neq TL(L)$ for $L = 1, \dots, NT$ set $t = t + 1$

and $TL(t) = T(m,j)$, then set up node $p = NN + t$,

connect node k to node p to form arc (k,p) and

label the arc $(0,1,0)$, go to STEP 4.

STEP 6: Set up node $K_s = K + 1$ (to represent a supersource for multi-section course j). Connect node 1 to node K_s to form the arc $(1,K_s)$ and label the arc

$(1,1,0)$. Set up nodes $K_s + 1, K_s + 2, \dots, K_s + m, \dots, K_s + S_j$ (to represent all sections of course j) and connect node K_s to each one of these nodes to form arcs $(K_s, K_s+1), (K_s, K_s+2), \dots, (K_s, K_s+S_j)$ and label all these arcs $(0,1,C_{j_m})$, set $m = 0$.

(a) $m = m + 1$

If $m > S_j$ go to STEP 4

(b) 1. If there is an L ($L = 1, \dots, NT$) for which $T(m,j) = TL(L)$, then a node $P = NN + L$ has already been set up. Connect node $K_s + m$ to node p to form arc (K_s+m,p) and label it $(0,1,C_{j_m})$, go to STEP 6-(a).

2. If $T(m,j) \neq TL(L)$ for $L = 1, \dots, NT$, set $t = t+1$ and $TL(t) = T(m,j)$, then set up node $p = NN + t$, connect node K_s+m to node p to form arc (K_s+m,p) and label the arc $(0,1,C_{j_m})$, go to STEP 6-(a).

STEP 7: Set up node $Q = NN + t + 1$ (to represent the last supersink), then connect nodes $NN + 1, NN + 2, \dots, p$ to node Q to form arcs $(NN + 1, Q), (NN + 2, Q), \dots, (P,Q)$ and label these arcs $(0,1,0)$.

STEP 8: Connect node Q to node 1 to form the reverse arc $(Q,1)$ and label the arc $(N,N,0)$. This completes the formulation of a minimum cost circulation network flow problem for student i . The solution to this problem, if one exists, would show a feasible schedule for the student. Go to STEP 1.

3.4. Model Flexibility in More General Cases

More complex situations could be modeled for more sophisticated class schedules. For example, if the school offers courses of several-hour type or courses with overlapping time periods a network flow problem could still be set up for each student. Limitations on the number of students allowed to be scheduled into a particular section could also be incorporated into the network problem formulation. When there are several-hour courses a node is considered for the course at each time period and a supersource considered for all these nodes. A flow equal to the number of hours that the course requires is forced into the supersource and therefore the student would be forced to be scheduled into all those time periods that the course requires. If there are overlapping time blocks for two courses and a student is requesting both of these courses, then there is no feasible solution for this student.

Suppose there is a limit on the number of students who could be scheduled into a particular section of a course. In this case, schedule students as described originally and if the number of students scheduled into this particular section is more than the limit indicated, reschedule all the students who are requesting this course. This time assign a higher per unit cost to the arc going into the node representing the particular section. If the number

of students scheduled into the section is still higher than the limit, the per unit cost could be raised again and students previously scheduled into this course rescheduled until the number of students scheduled into the section is less than or equal to the required limit. If a very high per unit cost for the above-mentioned arc is selected, then only those students would be scheduled into this section for which there is only one feasible schedule. If the number of students scheduled into this section is still higher than the limit, then the difference between number of students scheduled into the section and the indicated limit would show the number of conflicts (students with no feasible schedule) that actually would result if a limiting constraint is imposed.

3.5. Model Application

Suppose a heuristic approach is used in scheduling a school and a class schedule is generated, students are scheduled into classes, and a number of students are indicated as those encountering a conflict. As explained above a number of these conflicts could possibly be resolved by moving students from one section of a multi-section course to another. This means that the class schedule generated by a heuristic approach would not normally produce the best possible result. It is, however, possible to use the model discussed here and an algorithm similar to the one discussed in section 3.3 to find a feasible schedule for some conflicts labeled resolvable. A network flow problem could be set up for each student

encountering a conflict based on the courses that the student is requesting and the time periods at which these courses are scheduled. A feasible solution obtained for each network problem would indicate that the corresponding student's conflict could be resolved and a feasible schedule for this student exists. All the feasible solutions would show all the conflicts which could be resolved. All the infeasible solutions would show the total number of irresolvable conflicts. This would be the minimum number of conflicts for the class schedule.

3.6. Computational Methods

Recent computational advances in the development of computer codes to solve network flow problems show that many of these problems could be solved far more effectively than in the past. As Glover and Klingman [36] indicate new computer codes have succeeded in solving even large size problems (several thousand nodes and several thousand arcs) with an efficiency many times beyond codes previously available. The network flow problem considered here, for each student, is of a simple type normally, with less than 50 nodes and 100 arcs. These problems could be easily and efficiently solved using very little central computer memory. Glover and Klingman [38] indicate that new computer codes for solving network problems are extremely efficient and, as an example, a network problem with 1000 nodes (equations) and 7000 arcs (variables) can be solved in 8 seconds on an

IBM 360/65 using the highly efficient network code [40]. As indicated earlier two network flow problems, of the type discussed here, were solved using an out-of-kilter algorithm code. The problems one with 12 nodes and 35 arcs and the other with 16 nodes and 35 arcs required .8 and 1.8 seconds of the IBM 370/158 CPU time, respectively. Some of the references in this area are Barr et al. [41], Clasen [42], and Edmonds and Karp [43].

3.7. Summary

In this chapter a network flow model was developed for scheduling students into an existing class schedule in an optimal fashion. Existing solution procedures for the type of network flow problems developed and other applications of the model were discussed.

CHAPTER IV

A HEURISTIC APPROACH FOR CLASS-SCHEDULING IN TRADITIONAL SCHOOLS

As indicated earlier, general purpose scheduling programs have been developed which are very flexible and can be used to perform scheduling for a wide range of schools. To obtain more flexibility, scheduling programs have been developed which utilize what is called "modular scheduling." Modular scheduling is the case where the basic time unit considered is small and any class period could be composed of one or more of these basic units. This type of scheduling offers a great amount of flexibility which then allows the school program to incorporate all sorts of flexible schedules, programs, and new ideas. On the other hand, modular scheduling requires a great amount of information regarding students' course requests, course requirements, teacher requirements, preferences, classroom requirements, etc. This obviously indicates that very flexible scheduling requires a long planning period for data collection, a huge computer program to consider all types of situations and restrictions, technical assistance for data setup and computer runs, and

could become uneconomical in terms of personnel time, computer time, and computer memory requirements.

Since most schools follow a traditional or conventional scheduling system with little flexibility, a relatively small and simple to implement computer program for their class-scheduling could provide them with efficiency and economy. Conventional schools are characterized by a school structure where the cycle is one day, all the courses meet for one period during the day, and courses could offer more than one section. A modified school structure is very similar to a conventional school structure except that some of the courses offered would meet daily for more than one hour.

A heuristic approach and a computer program is developed intended to be used by schools with a modified structure or a traditional structure.

Heuristic and mathematical optimization techniques are incorporated in this program to generate class schedules with relatively few student conflicts, be fast as far as CPU time is concerned, and be small in terms of memory requirements.

4.1. General Problem Description

The approach is basically designed for small to medium size schools (300-800 students), although it could be used for large size schools as well.

Schools with a modified structure and as a special

case of this structure, schools with a traditional structure, would be considered in the approach and program development. This provides some flexibility in the sense that a large number of schools could use the program. In the above structures a school day is divided into six time periods of an hour each. If a class-schedule is constructed, it would remain the same every day of the week. Students, on the average, request six courses per semester.

There are two semesters in a school year. Course offerings are such that some courses are offered only in the first semester, some only in the second semester, and other courses which start in the first semester and continue through the second semester. These types of courses are called first-semester, second-semester, and all-year courses, respectively. In each category there are courses which offer only one section, and courses which offer more than one section. Most of these courses are only one hour long, but there are some courses, called several-hour courses, which require time periods longer than one hour. The latter type courses normally could require successive time periods of up to four hours long and are such that a group of students and one or more classrooms are assigned to a group of teachers for several hours. The teachers then plan and perform whatever program they are interested in. There are some courses which must be scheduled at specific time periods. These courses are called locked-in courses. Still, there are other types of courses where a group of them must be scheduled at

different time periods because of lack of enough teachers or rooms.

4.1.1. Basic Assumptions

The basic assumptions, based on which the heuristic approach will be developed, apply to the type of problem discussed above and are as follows:

A limited set of time periods (non-overlapping) are given. A number of courses are to be scheduled at these time periods such that a set of constraints are satisfied and the number of students encountering a conflict is minimum. Any number of courses could be scheduled at a given time period. A course or a section of a course must be scheduled at one and only one time period except for several-hour courses which must be assigned to several successive time periods. A student can take one and only one course (or a section of a course) at a time period. If two courses are scheduled at a time period some conflicts could encounter and the number of conflicts is equal to the number of students requesting both courses. A set of courses taught by a single teacher or a group of courses utilizing a certain room, must be scheduled at different time periods.

4.1.2. Heuristics

School officials who have been doing scheduling manually for several years know that, in general, some courses or some students are harder to schedule than other courses or other students. This knowledge, based on their past

experience, gives them the ability to recognize those courses which are hard to schedule and those students who are hard to schedule. Accordingly, they would schedule those courses or students at the beginning of their scheduling process as a rule. There are several heuristics that could advantageously be used in developing a scheduling program, some of which are as follows:

- (a) Generally speaking, single-section courses are harder to schedule than multi-section courses.
- (b) In order to make the choice as wide as possible, different sections of a course should be scheduled at different time periods.
- (c) The more single-section courses a student is requesting the harder it is to schedule him.
- (d) The more crowded courses should be scheduled before the less crowded courses.
- (e) Two courses offered in different semesters but both requested by many students should be scheduled at the same time.

4.2. The Heuristic Solution Approach

The solution procedure for the above scheduling problem involves dividing the original problem into several subproblems according to the type (locked-in, single-section, multi-section, etc.) of courses which are being considered for scheduling.

The subproblems are considered for scheduling in a

successive manner and the interaction between them is the order in which they are considered. The order in which a subproblem is considered for scheduling is based on the relative scheduling difficulty of courses comprising the subproblem. Although determining which types of courses are more difficult to schedule than other types is somewhat subjective, a good heuristic is that the more restricted courses are harder to schedule than less restricted courses.

This heuristic is used to identify subproblems representing course types and arranging them in the following manner: Locked-in, Several-hour, Teacher or room restricted single-section, Unrestricted single-section, and Multi-section courses. Locked-in courses are normally single section and restricted to a single time period and therefore, involve no decisions.

Each subproblem is further partitioned into smaller subproblems of one or more courses to be considered for scheduling at a time. The interaction between new subproblems is again determined by the order in which they are considered for scheduling. The order is determined by assigning a number, called an index of difficulty, to each course which represents its relative scheduling difficulty with respect to other courses. The index of difficulty is assigned based on one or more heuristics and is computed for different course types differently. The solution approach originally is designed to arrange the courses according to the five general types and according to descending order of index of difficulty within

each type. The scheduling process starts by scheduling locked-in courses one by one first. Then, other course types are scheduled successively using algorithms developed based on heuristics or mathematical optimization techniques or a combination of both.

The only course type where the developed algorithm considers more than one course for scheduling at a time is the teacher or room restricted single-section courses.

In addition to the use of heuristics and mathematical optimization techniques, the special structure of the problem has been used advantageously, whenever possible, to reduce computer time and memory requirements. Finally, the program developed for the implementation of the solution approach is developed in several components such that each component could be processed to schedule a portion of the overall problem independently. The advantages are that:

- (a) Instead of one huge program several small programs are processed successively so that memory requirements are kept down.
- (b) If after scheduling a portion of the courses the resulting schedule is not satisfactory changes could be made before all other courses are scheduled.
- (c) If an error is detected at a stage only the component corresponding to that stage would have to be processed again.

A conflict matrix is computed for all locked-in,

several-hour, and single-section courses. As an example consider:

Courses			
<u>Courses</u>	<u>529</u>	<u>550</u>	<u>571</u>
529	0	6	0
550		0	12
571			0

The first row indicates that if courses 529 and 550 are scheduled at the same time 6 students would encounter conflicts and if courses 529 and 571 are scheduled at the same time no conflicts would result.

The next sections discuss computation of index of difficulty, heuristic or mathematical techniques utilized, and the algorithm used for scheduling of each type of course.

4.2.1. Treatment of Locked-In Courses

These courses are supposed to be scheduled at a specific predetermined time period. Arrays are used to store relevant information to be used when other courses are scheduled later.

4.2.2. Scheduling Several-Hour Courses

These courses require several successive time periods and are originally ordered such that the lengthier courses would be considered for scheduling first. For example, courses which require four-hour time blocks are considered for scheduling before courses which require three-hour time

blocks and so on. Since these courses would be arranged based on the number of time periods that they require, they are arbitrarily assigned an index of difficulty of zero. These courses could be time restricted. For example, a four-hour course could be restricted to morning time periods of 1, 2, 3, and 4.

To explain the heuristic algorithm used for scheduling several-hour courses an example is given. Consider a three-hour course 571 which has no time restrictions and all the locked-in courses are already scheduled as follows:

Time Period	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>
Course	5	23	538	540	360	529

Course 571 would temporarily be scheduled at time periods 1, 2, 3 such that:

Time Period	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>
Course	5	23	538	540	360	529
	571	571	571			

And total number of conflicts between 571 and other courses is computed. The total number of conflicts between 571 and all other courses is computed by identifying (from stored information in arrays) all the courses which are already scheduled at time periods 1, 2, and 3 and summing up the number of conflicts between each one of these courses and 571. Course 571 would temporarily be scheduled again at time periods 2, 3, 4 and again total number of conflicts computed.

The above procedure would be repeated for course 571 for time periods 3, 4, 5 and 4, 5, 6. Finally, course 571

is scheduled permanently at those set of successive time periods for which the penalty is a minimum.

4.2.3. Scheduling Single-Section Teacher or Room Restricted Courses

These courses are arranged such that all the courses (single-section) restricted by a teacher or a room are in one group and the order of the courses within a group is not important. The groups of restricted single-section courses are arranged such that a group of six courses are before a group of five courses which are in turn before a group of four courses and so forth. This arrangement represents the relative scheduling difficulty of different groups and implies that it is more difficult to schedule a group of six restricted courses than it is to schedule a group of five restricted courses and so on.

If there are sets of restricted courses such that some of the courses are single-section and some others are multi-section, only single-section courses of each set would be considered as a group in this category. The multi-section restricted courses would be considered later when multi-section courses are considered for scheduling. To explain the method used for scheduling restricted single-section courses let us consider the following example:

Suppose all the locked-in and several-hour courses are already scheduled as shown in the following table:

Time Period	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>
Course	5	23	538	540	360	529
	571	571	571			
				573	573	573

Suppose courses 9, 10, 11, and 12 are a group of teacher-restricted single-section courses which must be scheduled at different time periods. Consider the conflict matrix C_{jk} computed below.

Time Periods		<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>
Courses	9	2	1	4	3	0	5
	10	3	1	6	2	1	0
	11	1	3	8	4	6	2
	12	5	1	3	7	2	3

Entries represent number of conflicts for scheduling a particular course at a particular time period considering all other courses which are already scheduled. The first row, for example, indicates that if course 9 is scheduled at time period 1, the number of students who would encounter a conflict is 2 (the two conflicts would result because of students requesting courses 9, 5 or 9, 571 or 9, 5, 571).

The entries in the first row also indicate that if course 9 is scheduled at time periods 2, 3, 4, 5, or 6 the number of potential conflicts would be 1, 4, 3, 0, or 5, respectively. Total number of potential conflicts for a course x if scheduled at a time period t is computed such

that C_{xt} is the sum of the number of conflicts between x and all other courses already scheduled at time period t which are all-year courses or are offered in the same semester as x is, minus sum of all the conflicts between x and all other courses scheduled at time period t which are offered in a different semester than x is.

A network flow problem is then constructed as follows:

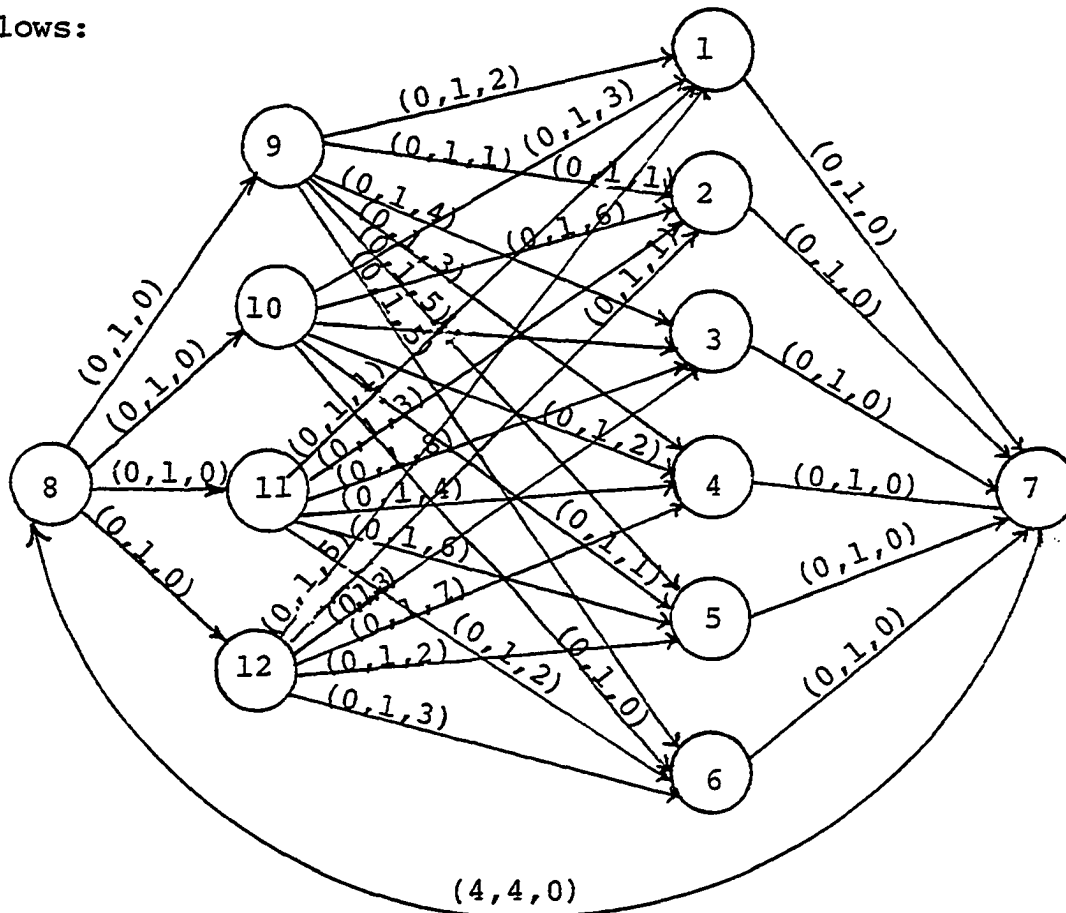


Figure 7.

Nodes 1, 2, 3, 4, 5, and 6 represent time blocks; nodes 9, 10, 11, and 12 represent courses. Arcs from a certain course node to a certain time node indicate that the certain course could be scheduled at the corresponding time

block and the label $(0,1,C_{jk})$ indicates that zero or one unit of flow could move from the course node to the time node and C_{jk} is the cost of per unit flow. For example, node 9 represents course 9 which can send 0 or 1 unit of flow to node 1 which represents time period 1 and the cost of sending a unit of flow C_{jk} would be equal to 2, the number of conflicts which would result if course 9 is scheduled at time period 1. Node 8 is a supersource and node 7 is a supersink. The arc (7-8) has a label $(4,4,0)$ which indicates that exactly four units of flow must flow from node 8 to all course nodes 9, 10, 11, and 12 (one unit of flow to each node) and then to four of the time nodes 1, 2, 3, 4, 5, and 6 and finally to 7. This is the general form of a network flow minimum cost circulation problem. The solution to this problem, using an efficient algorithm such as Out-of-Kilter algorithm by Fulkerson [39], would determine the flow of the four units from course nodes 9, 10, 11, and 12 to time nodes 1, 2, 3, 4, 5, and 6 that minimizes the total cost. The solution to the above network problem using a computer code for Out-of-Kilter algorithm indicates that the minimum cost corresponds to the case where one unit flows from node 9 to node 5, another unit flows from node 10 to node 6, a third unit flows from node 11 to node 1, and a fourth unit flows from node 12 to node 2 for a total cost of 2 conflicts. This means that courses 9, 10, 11, and 12 must be scheduled at time periods 5, 6, 1, 2 respectively to generate an optimum schedule for this course

group with respect to all the courses already scheduled, and with a minimum of 2 conflicts.

Finally, the above network flow problem is equivalent to the following assignment problem.

$$\begin{aligned}
 \text{Min} \quad & \sum_{j=9}^{12} \sum_{k=1}^6 C_{jk} x_{jk} \\
 \text{S.T.} \quad & \sum_{j=9}^{12} x_{jk} \leq 1 \quad k = 1, \dots, 6 \quad \text{at most one course could} \\
 & \hspace{15em} \text{be scheduled at a time} \\
 & \hspace{15em} \text{period} \\
 & \sum_{k=1}^6 x_{jk} = 1 \quad j = 9, 10, 11, 12 \quad \text{course } j \text{ must be} \\
 & \hspace{15em} \text{scheduled at a time} \\
 & \hspace{15em} \text{period } k \\
 & x_{jk} = 0 \text{ or } 1 \quad j = 9, 10, 11, 12 \quad k = 1, \dots, 6
 \end{aligned}$$

where C_{jk} is the potential cost of scheduling course j at time period k considering all the courses already scheduled at time period k .

4.2.4. Scheduling Single-Section Unrestricted Courses (No Teacher or Room Restrictions)

Single-section courses are considered for scheduling one by one according to descending order of their index of difficulty. The index of difficulty for this type of course is computed as follows:

Let S represent all the students requesting a single-section course J . Suppose a student $s \in S$ is requesting N courses and $C_s(K)$, $k = 1, \dots, N$ represent the set of courses that the student is requesting.

For student s an index of difficulty $IDFS(s)$ is computed as follows:

$IDFS(s) = 0$

for $K = 1, N$

If $C_s(K)$ is a locked-in, several-hour, or single-section restricted course, then $IDFS(s) = IDFS(s) + 3$

If $C_s(K)$ is another single-section course not restricted by teachers or rooms, then $IDFS(s) = IDFS(s) + 2$

If $C_s(K)$ is a two-section course, then $IDFS(s) = IDFS(s) + 1$

If $C_s(K)$ offers more than two sections, then $IDFS(s) = IDFS(s) + 0$

The course index of difficulty is now computed as:

$$IDF(J) = \sum_{s=1}^S IDFS(s) + a + b*TR$$

where a is an arbitrary large number which keeps the index of difficulty for single-section courses higher than all the multi-section courses.

b is a constant considered for time restrictions, and TR is the number of time period restrictions.

For example, a course which is restricted to morning time periods of 1, 2, 3, and 4 should not be scheduled at time periods 5 or 6 and therefore has a time restriction coefficient of $TR = 2$.

To schedule single-section courses, the course with the highest index of difficulty would be considered first.

The course considered for scheduling would be temporarily scheduled at time period 1, then 2, then 3, and so

on. At each time period, the number of potential conflicts between this course and all other courses already scheduled at this time period would be computed. The time period with minimum number of potential conflicts would be selected. The course presently being considered would be scheduled at this time period permanently. In the case of a tie that time period would be selected at which there are fewest number of courses already scheduled.

4.2.5. Scheduling Multi-Section Courses (2 to 5 Sections)

Multi-section courses are types of courses which offer more than one section. In this section only multi-section courses which offer less than six sections (number of periods) are considered. Courses which offer six sections or more will be considered in the next section. The index of difficulty for these types of courses is computed based on such factors as number of sections, course restrictions, time restrictions, alternative choices, and number of students requesting a course. Course restriction means that a multi-section course presently being considered for scheduling is restricted by another course already scheduled, and different sections of the present course must not be scheduled at the same time as that of the course already scheduled. For example, suppose courses 560, 561, and 562 are all physics courses such that course 560 is a single-section course, 561 and 562 are two-section courses and there is only one physics teacher available. Since course 560 is a single section

course, it would be scheduled when single-section courses are considered for scheduling prior to multi-section scheduling process. Then, course 561 is restricted by course 560 and course 562 is restricted by 560 and 561 (both sections).

In the above example if everything is assumed as it is except that course 562 is a three-section course then the three sections must be scheduled at those three time periods where courses 560 and 561 are not scheduled. A three section course such as 562 should then have a higher index of difficulty than those three-section courses which have more than one alternative scheduling choice. It should also be mentioned here that since course 562 is restricted by 561 it must be scheduled after course 561 is scheduled and therefore course 561 must have an index of difficulty higher than 562.

Considering the above discussion, the index of difficulty for a multi-section course J is computed as follows:

Suppose K = total number of available time periods

$ITST(J)$ = total number of students requesting course J

$L1$ = number of time periods at which course J could be scheduled

$MSEC(J)$ = number of sections offered by course J

$K11 = L1 - MSEC(J)$

$$TR = \text{Time Restrictions} = \begin{cases} 0, & \text{if there are no time restrictions} \\ K-L1 & \text{if there are time restrictions} \end{cases}$$

$$CR = \text{Course Restriction} = \begin{cases} \text{Sum of number of sections} \\ \text{of a set of courses } C(J); \\ \text{if course } J \text{ is restricted} \\ \text{by } C(J) \\ 0, \text{ otherwise} \end{cases}$$

Then, the index of difficulty for course J is:

$$IDF(J) = \begin{cases} c + ITST(J) + d \cdot TR + e \cdot CR - f \cdot MSEC(J) & \text{if } K11 > 0 \\ c + ITST(J) + d \cdot TR + e \cdot CR - f \cdot MSEC(J) + g & \text{if } K11 = 0 \end{cases}$$

where c, d, e, f, and g are constants. The constant c is a large number which is added to index of difficulty of all those courses which offer more than one section but less than six sections. The value of c is selected such that the index of difficulty of all courses offering two to five sections would be higher than the index of difficulty for all the courses offering more than five sections but less than the index of difficulty for single-section courses. In summary, c is selected such that

indices of difficulty for courses offering more than 5 sections	<	indices of difficulty for courses offering 2-5 sections	<	indices of difficulty for single-section courses
---	---	--	---	---

d, e, and f are all arbitrary constants representing scheduling difficulty of a multi-section course with respect to time restrictions, course restrictions, and the number of sections offered by the course, respectively. Finally, g is an arbitrary constant representing scheduling difficulty when there is only one alternative scheduling choice available.

To explain the algorithm for scheduling multi-section courses an example is presented. Suppose a class

schedule is being prepared for a school and all the locked-in, several-hour, single-section teacher or room restricted, and single-section unrestricted courses are already scheduled as follows:

Time Periods	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>
Courses						
Locked-in	5	23	538	540	360	529
Several-hour	571	571	571			
				573	573	573
Teacher or room restricted single-section	11	12			9	10
Unrestricted single-section	320		560	118		559

Course 320 is a second-semester course and all other courses are either first-semester or all-year courses.

Suppose now that the two-section course 61 is being considered for scheduling and this course is (a) restricted by course 560 (course restriction) and (b) must not be scheduled at time period 6 because of teacher unavailability (time restriction), and (c) it is a first-semester course.

Using restriction information course 560 could be located and noticed that course 61 cannot be scheduled at time period 3 (because of course 560) and time period 6. Since course 560 is already scheduled at time period 3, course 61 could be scheduled at time periods 1, 2, 4, and 5. Several alternative scheduling choices are available:

1. Schedule course 61 section 1 at time period 1 and section 2 at time period 2.
2. Schedule course 61 section 1 at time period 1 and section 2 at time period 4.
3. Schedule course 61 section 1 at time period 1 and section 2 at time period 5.
4. Schedule course 61 section 1 at time period 2 and section 2 at time period 4.
5. Schedule course 61 section 1 at time period 2 and section 2 at time period 5.
6. Schedule course 61 section 1 at time period 4 and section 2 at time period 5.

Total number of alternatives available = $C_2^4 = \frac{4!}{2!2!} = 6$

(For a three-section course and 6 time periods the number of alternative choices would be $C_3^6 = \frac{6!}{3!3!} = 20$.)

To select one of the alternative choices two heuristics are considered.

1. First heuristic is used to avoid generating conflicts and is stated as: schedule different sections of the course at that alternative choice which generates the minimum number of conflicts considering all the courses already scheduled.
2. The second heuristic is used to avoid generating conflicts in the future when other multi-section courses are scheduled and is stated as: schedule a first-semester

course at the same time as a second-semester course and vice versa, if a group of students have requested both courses.

As an example of the latter heuristic, suppose a student has requested courses 5, 23, 306, 324, 538, 540, 360, and 529. Suppose further that courses 5 and 23 are first-semester courses, 306 and 324 are second-semester courses and 538, 540, 360, and 529 are all-year courses. Considering that there are six time periods, the student could have a feasible schedule if the courses are scheduled in an arrangement such as follows:

Time Periods	<u>t₁</u>	<u>t₂</u>	<u>t₃</u>	<u>t₄</u>	<u>t₅</u>	<u>t₆</u>
Courses	5	23	538	540	360	529
	306	324				

The student would not encounter a conflict because although scheduled at the same time courses 5 and 306 are offered during two different semesters. The same argument is true for courses 23 and 324. Another arrangement in which the student could have a feasible schedule is:

<u>t₁</u>	<u>t₂</u>	<u>t₃</u>	<u>t₄</u>	<u>t₅</u>	<u>t₆</u>
5	23	538	540	360	529
324	306				

But, if courses 5, 23 and 306, 324 are scheduled at more than two time periods, if not now, the student will

encounter a conflict later when eventually all these courses are considered for scheduling.

The problem of finding the best alternative choice for a multi-section course J with respect to all the courses already scheduled could be formulated as an integer programming problem as follows:

$$\begin{aligned}
 &\text{Max} \quad \sum_{i,s,k} x(i,s,k) \\
 &\text{S.T.} \quad \sum_k y(s,k) = 1 \quad s = 1, \dots, S_j \\
 &\quad \quad x(i,s,k) \leq y(s,k) \quad \begin{array}{l} i = 1, \dots, I \\ s = 1, \dots, S_j \\ k = 1, \dots, N \end{array} \\
 &\quad \quad \sum_{s \in S_j} x(i,s,k) \leq 1 \quad \begin{array}{l} i = 1, \dots, I \\ k = 1, \dots, N \end{array} \\
 &\quad \quad \sum_{i \in P_k} \sum_{s \in S_j} x(i,s,k) = 0 \quad k = 1, \dots, N
 \end{aligned}$$

where:

$$\begin{aligned}
 x(i,s,k) &= \begin{cases} 1 & \text{if student } i \text{ is scheduled into section } s \text{ at time period } k \\ 0 & \text{otherwise} \end{cases} \\
 y(s,k) &= \begin{cases} 1 & \text{if section } s \text{ is scheduled at time period } k \\ 0 & \text{otherwise} \end{cases}
 \end{aligned}$$

I is the set of students requesting course J

S_j is the number of sections offered by course J

P_k is the set of students who cannot be scheduled

at time period k (without encountering a conflict).

In this formulation, the objective function maximizes

the number of students successfully scheduled into different sections of the course presently being scheduled. The first constraint insures that a section s is scheduled at one and only one time period. The second constraint states that student i could be scheduled into section s at time period k if section s is scheduled at time period k . The third constraint insures that student i is scheduled into at most one of the sections of the course presently being considered for scheduling. The fourth constraint forces $x(i,s,k)$'s to zero for those students who cannot have a feasible schedule at time period k considering all other courses already scheduled at k .

This model does not easily lend itself to a simple solution procedure and can still be a large size problem if a multi-section course is requested by many students. As an example if a 5-section course is requested by 200 students and the sections could be scheduled to any of the 6 time periods then the number of $x(i,s,k)$ variables would be $200 \times 5 \times 6 = 6000$ and the number of $y(s,k)$ variables would be $5 \times 6 = 30$. Total number of constraints would be 7211.

A different heuristic procedure is employed here for scheduling multi-section courses which is extremely efficient and produces the best solution (or one as good as the solution which could be obtained by utilizing the above model) for a multi-section course presently being scheduled with respect to all other courses previously scheduled. To show how this procedure works let us consider the previous

example where a set of courses are already scheduled and course 61, a first semester two-section course, could be scheduled such that different sections may be assigned to time periods 1, 2, 4, and 5. Suppose further that the identification numbers for 18 students requesting course 61 are given as:

1014 1030 1075 1091 1094 1112 1148 1156 1160
1169 1174 1206 1240 1257 1279 1282 1286 1288

All these students have requested other courses, some of which are already scheduled, and some other ones which will be scheduled after course 61 is scheduled. If any of these students has requested a course which is already scheduled at time period k , then the student cannot be successfully scheduled at period k and if a section of course 61 is scheduled at time period k the student cannot be successfully scheduled into this section. Based on this concept a list is determined of all the students who could successfully be scheduled at each time period 1, 2, 4, and 5 as follows:

Time Period	Student ID								
1	1014	1030	1075	1112	1160	1169	1240	1257	
	1279	1282	1288						
2	1112	1257							
4	1030	1075	1112	1148	1156	1169	1174	1206	
	1279	1282	1286	1288					
5	1030	1075	1112	1148	1156	1169	1174	1279	
	1282	1288							

The above list simply indicates that, as an example, students whose ID's are 1014, 1030, 1075, ... could successfully be scheduled at time period 1, but students whose ID's are 1094, 1148, 1156, ... cannot be successfully scheduled at this time period because these ID's are not included in the list for time period 1.

Since 18 students have requested course 61, a desired situation would be that approximately 9 students be scheduled into each of the two sections of the course, if possible. It is seen from the above lists that at time period 2 a maximum of 2 students could successfully be scheduled. This means that if one of the sections is scheduled at time period 2 and the other section at one of the time periods 1, 4, or 5, a balance between the two sections is almost impossible and it is highly likely that many students would encounter a conflict. Therefore, it could very well be advantageous to disregard time period 2 from further considerations and only investigate time periods 1, 4, and 5 for determining the best alternative choice. The selection of the best alternative choice at this point is simply determining, from the ID lists given above, the combination of time periods (1,4), (1,5), or (4,5) which would produce minimum number of students who would encounter a conflict.

Since course 61 is a first-semester course, if the sections are scheduled such that they are at the same time periods as a second-semester course (or sections of a second-

semester course), most of the students requesting both courses could be scheduled at the same time into both courses and possibly avoid many future conflicts. For this reason, the number of students requesting course 61 and all the second-semester courses already scheduled at each of the alternative choices (1,4), (1,5), and (4,5) are subtracted from the number of conflicts which is computed at the corresponding alternative choice and then the minimum selected. For example, if course 320 is the only second-semester course already scheduled at time period 1 and there are three students requesting both courses 320 and 61 these three students scheduled into a section of 61 at time period 1 could be prevented from encountering a conflict in the future.

Assume a situation for course 61 and alternative choices (1,4), (1,5), and (4,5) with respect to all the courses which are already scheduled such that for course 61 if:

1. Section 1 is scheduled at time 1 and section 2 at time 4 number of conflicts would be 5.
2. Section 1 is scheduled at time 1 and section 2 at time 5 number of conflicts would be 6.
3. Section 1 is scheduled at time 4 and section 2 at time 5 number of conflicts would be 3.

It is seen that the minimum number of conflicts would be for the case where section 1 is scheduled at time

period 4 and section 2 at time period 5 and, therefore, the third alternative choice seems to be the best at this point. But considering the fact that three students are requesting both courses 61 and 320 they would encounter a conflict, if a section of 61 is not scheduled at time period 1, now or later when some of their other courses are scheduled. By scheduling section 1 at time 1 (where the three students could successfully be scheduled into both courses) and section 2 at time 4 the number of conflicts could be computed as:

$$\begin{aligned} \text{Number of conflicts} &= \text{Number of conflicts which would encounter considering all other courses already scheduled} \\ &\quad - \text{Number of potential conflict reduction in the future} = \\ &\quad 5 - 3 = 2 \end{aligned}$$

And, therefore, alternative choice (1,4) is the best choice.

Now, assuming the selection is alternative choice (1,4) for permanently scheduling the two sections of course 61 the new problem is to schedule the students into the different sections such that potential future conflicts are prevented and, also, there is a balance between the two sections of the course. For our example suppose section 1 is permanently scheduled at time period 1 and section 2 at time period 4. The first student requesting course 61 is the student whose ID is 1014. If this student is requesting a second-semester course which is already scheduled at time period 1, the student would be scheduled into section 1 (at

time period 1). If this student is requesting a second-semester course which is already scheduled at time period 4, the student would be scheduled into section 2 (at time period 4). If this student is not requesting a second-semester course already scheduled at either time periods 1 or 4, then he would be scheduled into that section (1 or 2 of 61) which has the fewest number of students already scheduled in it (assuming he could successfully be scheduled into both time periods). This is an attempt to develop a balance between the two sections. If this student could be successfully scheduled into one of the time periods, then he would be scheduled into the section corresponding to that time period. Finally, if the student cannot be successfully scheduled into either of the two time periods, the student is encountering a conflict.

The above procedure is repeated for all the students requesting course 61, one by one, until all the students are either successfully scheduled into a section or identified as encountering a conflict.

The scheduling of all multi-section courses (2 to 5 sections) is the same as explained here for course 61 and is performed for all courses successively according to the descending order of their indices of difficulty.

4.2.6. Scheduling Multi-Section Courses (6 Sections or More)

These are types of courses which offer 6 sections or more. Since there are totally 6 time periods available, it

is best to schedule at least one section at each time period, if possible, to prevent conflicts. Then, schedule additional sections at those time periods where many students could successfully be scheduled such that eventually there is some balance between the sections and some balance between the time periods. Once it is decided how many sections would be scheduled at each time period, then the scheduling of students into these sections is the same as student scheduling in the previous section (4.2.5).

To show how scheduling is done for those courses which offer 6 sections and more, let us consider the following example. Suppose a class schedule is being prepared for a school and all the locked-in, several-hour, single-section and multi-section (less than 6 sections) are already scheduled as follows:

Time Periods	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>
Locked-in courses	5	23	538	540	360	529
Several-hour courses	571	571	571	573	573	573
Single-section courses	11 320	12	560	118	9	10 559
Multi-section courses	61-1 530-1		530-2	61-2	530-3	530-4

Suppose further that course 580 offers 10 sections, is requested by 200 students, and is considered for scheduling at this time.

It is desirable that each section of the course average approximately 20 students. This could formally be defined by $N0 = W \cdot K1 / M$, where W is a section balance indicator, $0 \leq W \leq 1$, and $K1$ and M are number of students requesting the course and number of sections offered by the course, respectively. Suppose the number of students who could successfully be scheduled at each time period K , $SN(K)$, could be computed and is given as:

Time Periods	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>
$SN(K)$	8	62	36	110	85	28

This means that out of 200 students requesting course 580, 8, 62, 36, 110, 85, and 28 could successfully be scheduled at each time periods 1, 2, 3, 4, 5, and 6, respectively.

It is noticed that only 8 students could successfully be scheduled at time period 1 and this is much lower than an average of 20 students per section. It could very well be decided then, that no sections be scheduled at time period 1.

Considering time periods 2, 3, 4, 5, and 6 indicates that at each one of these periods more than an average of 20 students could successfully be scheduled and, therefore, a section could be scheduled at each indicated time period. The remaining 5 sections should be scheduled at those time periods at which more than $2 \cdot 20 = 40$ students could successfully be scheduled. These time periods are 2, 4, and 5 where 62, 110, and 85 students could successfully be scheduled.

This indicates that 3 more sections, e.g., sections 6, 7, and 8 could be scheduled at time periods 2, 4, and 5, respectively. The remaining 2 sections should be scheduled at those time periods at which more than $3 \times 20 = 60$ students could successfully be scheduled. These time periods are again 2, 4, and 5. This indicates that the two remaining sections could be scheduled at two of the three time periods 2, 4, and 5. To select two of the three time periods, for the remaining two sections, it is noticed that only three courses are already scheduled at time period 2 and four courses are already scheduled at each one of the time periods 4 and 5. To have some balance between the periods one section could be scheduled at time period 2 and the final last section could arbitrarily be scheduled at either time periods 4 or 5, say period 4.

An attempt is made to keep conflicts at a minimum, have some degree of balance between the sections, and some degree of balance between the time periods. The course would then actually be scheduled as above and students scheduled into sections such that some balance eventually would exist between the sections and the procedure of scheduling students is the same as described in the previous section.

Suppose for the above same example a situation such as the following exists for number of students who could successfully be scheduled at different time periods.

Time Period	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>	<u>6</u>
SN(K)	8	42	36	68	42	28

On the average, it is desirable that approximately $NO = W \cdot 20 = 1 \cdot 20$ students be scheduled into each section. It is seen, from the above information, that at each one of the periods 2, 3, 4, 5, and 6 one section could be scheduled and approximately 20 students scheduled into it. This will take care of 5 sections 1, 2, 3, 4, and 5. For the remaining 5 sections it is seen that at each one of the time periods 2, 4, and 5 approximately 2 sections could be scheduled since at these time periods more than $2 \cdot 20 = 40$ students could successfully be scheduled. Therefore, sections 6, 7, and 8 could be scheduled at time periods 2, 4, and 5, respectively. Two more sections, i.e., sections 9 and 10 remain to be considered. One more section, say section 9 could be scheduled at time period 4, since $3 \cdot 20 = 60$ students could successfully be scheduled at this time period. The final section, i.e., section 10, then, cannot be scheduled at any of the time periods.

In a situation such as this the average number of students to be scheduled into each section would be redefined as follows: $NO = W \cdot 20 = .9 \cdot 20 = 18$, where $W = .9$ and the same procedure as above would be repeated to see if all 10 sections could successfully be scheduled, only this time on the average, it is desired that approximately 18 students be scheduled into each section. In this case the

procedure yielded a feasible schedule.

If the attempt to find successful schedules for all 10 sections would have failed for the second time, a new $NO = W \cdot NO = .9 \cdot NO = .9 \cdot 18 = 16$ would have been redefined and the above procedure repeated again.

The index of difficulty for these type of courses are determined simply based on their number of student requests. The more the number of students requesting a course the higher the index of difficulty of that course. For a course j the index of difficulty is equal to the number of students requesting the course.

4.3. The Heuristic Algorithm

List of Variables

NOS	Total number of students
IDS	Student identification number
NSCO	Number of courses requested by a student
JCRS(L)	ID number of courses requested by a student
L	An index $L = 1, \dots, NSCO$
KSEC(L)	That section of JCRS(L) at which the student is scheduled. If the student is not scheduled or the course is not scheduled $KSEC(L) = 0$.
NAME	Student name
IGRD	Student grade (9th or 10th, etc.)
ISEX	Student sex (M = Male, F = Female)
IDFS	Student index of difficulty
$C_s(L)$	Set of courses requested by student s , $L = 1, \dots, NSCO$

J	Course index number
ITST(J)	Total number of students requesting Jth course
IST1(N1)	ID of students requesting the Jth course where N1 = 1,...,ITST(J)
ICRS(J)	ID number of the course presently being con- sidered for scheduling
MSEC(J)	Number of sections offered by the Jth course
JS	Section index number, JS = 1,2,...,MSEC(J)
IDF(J)	Index of difficulty of the Jth course
ITYP(J)	Type of the Jth course (0 = Locked-in, >1 = Several-hour, =1, others)
NJ	Maximum number of students allowable in a section of the Jth course
ISEM(J)	Semester in which the Jth course is offered (0 = all-year, 1 = first-semester, 2 = second semester)
ITIME(J)	Course restriction (1 = there is course restriction, 0 = otherwise)
ITIME1(J)	Time restriction (1 = there is time restric- tion, 0 = otherwise)
ITIME2(J)	The starting time period if there is time restriction
ITIME3(J)	The final time period if there is time restric- tion
ITIME4(J)	Sum of the number of sections of those courses by which the Jth course is restricted
NCR(K)	ID of the Kth course by which a course is restricted (K = 1,...,ITIME4(J))
NHR	Number of hours required by a several-hour course
TR	Time restriction
CR	Course restriction

KG(IP)	Time period at which a course NCR(K) is scheduled by which the present course being scheduled is restricted
KT	Time period index, $KT = 1, 2, 3, 4, 5, 6$
N(KT)	Number of courses scheduled at time period KT so far
IN	$= N(KT)$
IKP(KT, IN)	Index number of the INth course scheduled at time period KT
IKC(KT, IN)	ID number of the INth course scheduled at time period KT
IKF(KT, IN)	The section of the INth course which is scheduled at time period KT ($IKF(KT, IN) = JS$ is the section of the INth course, $IKC(KT, IN)$, scheduled at time period KT)
MSEM(KT, IN)	The semester of the INth course which is scheduled at time period KT ($MSEM(KT, IN) = MSEM(J)$ is the semester in which the INth course is offered and the INth course $IKC(KT, IN)$, scheduled at time period KT)
IT	Iteration number
KCON(J1, J2)	Number of conflicts if courses J1 and J2 are both scheduled at the same time period
LCON(IT)	Total number of conflicts at iteration IT
NOCO	Number of restricted courses in a group of single section teacher or room restricted courses (in this case $NOCO = ITIME4(J)$)
C(J, K)	Total number of conflicts if the Jth course is scheduled at time period K (with respect to all courses already scheduled)
NCAN(KT)	Number of students which could successfully be scheduled at time period KT
NCRS(KT, KS)	ID of those students which could successfully be scheduled at time period KT, $KS = 1, \dots, NCAN(KT)$
KCRS(KT, KS)	$= NCRS(KT, KS)$

KTN(NP)	Time periods KT at which a section of the present course, being considered for scheduling, could be scheduled
L12(KT)	Number of students common between two courses which are offered in different semesters and considered for scheduling at time period KT
ID12(KT,I1)	ID of those students common between two courses which are offered in different semesters and considered for scheduling at time period KT where $I1 = 1, \dots, L12(KT)$
ITC(IT)	Total number of conflicts at iteration IT
JT(IT,JS)	Time periods at which different sections of the present course, being considered for scheduling, are temporarily scheduled where $JS = 1, \dots, MSEC(J)$
NK(KT)	Number of students scheduled at time period KT so far
NST(JS)	Number of students scheduled into section JS of the present course being considered for scheduling
MCRS(JS,I2)	ID of those students scheduled into section JS of the present course being considered for scheduling, $I2 = 1, \dots, NST(JS)$

Scheduling Locked-In Courses

```

STEP 1:  Set J = 0   N(L) = 0   L = 1,6
STEP 2:  Set J = J + 1   KT = ITIME(J)
        If ITYP(J) > 0   Set J = J - 1, Go to STEP 4
STEP 3:  N(KT) = N(KT) + 1
        IKP(KT,N(KT)) = J
        IKC(KT,N(KT)) = ICRS(J)
        IKG(KT,N(KT)) = IG(J)
        MSEM(KT,N(KT)) = ISEM(J)
        IKF(KT,N(KT)) = 0
        Go to STEP 2

```

· Scheduling Several-Hour Courses

```

STEP 4:  Set J = J + 1, LI = 1, LL = 6, IT = 1

        If ITYP(J) ≤ 1 course is not several-hour,
        set J = J - 1, go to STEP 9

STEP 5:  If ITIME(J) = 0 there are no time restrictions,
        go to STEP 6

        Otherwise, there are time restrictions

        LI = ITIME2(J)

        LL = ITIME3(J) Course J must be scheduled between

                time periods LI and LL

STEP 6:  LF = LI + NHR - 1

        for LK = LI to LF  compute:

        LCON(IT) = SUM KCON(J,J1)
                   LK

                where J1 = IKP(LK,I),
                   I = 1,N(LK)

STEP 7:  Set IHOLD(IT) = LI

        LI = LI + 1

        LF = LF + 1

        IT = IT + 1

        If LF ≤ LL  go to STEP 6

STEP 8:  Find I such that LCON(I) = min{LCON(I), I=1,...,IT}

        Set LI = IHOLD(I)

        LF = LI + NHR - 1

        for LK = LI to LF Compute:  N(LK) = N(LK) + 1
                                     IKP(LK,N(LK)) = J
                                     IKC(LK,N(LK)) = ICRS(J)
                                     IKG(LK,N(LK)) = IG(J)
                                     MSEM(LK,N(LK)) = ISEM(J)
                                     IKF(LK,N(LK)) = 0
                                     Go to STEP 4

```

```
STEP 9:  Set J = J + 1  LI = 1  LL = 6  IN = J  K = 1
        If ITIME(J) = 0 the course is not restricted,
        set J = J - 1,
        Go to STEP 15
```

```

If ITIME1(J) = 0  There are no time restrictions,
                  go to STEP 11

```

LL = ITIME3(J)

For $I = 1, N(K)$ Compute:

If ISEM(J) * MSEM(K,I) = 2 Courses J and J1 are
supposed to be offered at two dif-
ferent semesters. Go to STEP 11(b)

go to STEP 12

IF $K \leq LL$ go to STEP 11

If $J \leq IN + NOCO$ go to STEP 11

STEP 13: Solve the assignment problem:

$$\text{Min}_{j,k} C(J,K) x_{jk}$$

S.T. $IN+NOCO$

$$\sum_{J=IN} x_{jk} \leq 1 \quad K = LI, \dots, LF$$

$$\sum_{K=LI}^{LF} x_{jk} = 1 \quad J = IN, \dots, IN+NOCO$$

$$x_{jk} = 0 \text{ or } 1 \quad J = IN, \dots, IN+NOCO, K = LI, \dots, LF$$

STEP 14: For courses $J = IN, \dots, IN+NOCO$ and each $x_{jk} = 1$

compute:

$$\begin{aligned} N(K) &= N(K) + 1 \\ IKP(K, N(K)) &= J \\ IKC(K, N(K)) &= ICRS(J) \\ IKG(K, N(K)) &= IG(J) \\ MSEM(K, N(K)) &= ISEM(J) \\ IKF(K, N(K)) &= 0 \end{aligned}$$

Go to STEP 9

Scheduling Single-Section Unrestricted Courses

STEP 15: Set $J = J + 1$ $LI = 1$ $LL = 6$

If $MSEC(J) > 1$ the course is multi-section,

set $J = J - 1$, go to STEP 22

If $ITIME1(J) = 0$ There is no time restriction,

go to STEP 16

$LI = ITIME2(J)$

$LL = ITIME3(J)$

STEP 16: $K = LI$, $LCON(IT) = 0$, $IT = LI, LL$, $I = 0$ Initiatization

STEP 17: For $I = I + 1$

If $I > N(K)$ go to STEP 18

$J1 = IKP(K, I)$

If $ISEM(J) * MSEM(K, I) = 2$ J and J1 are offered in
different semesters,
go to STEP 17(b)

(a) $LCON(K) = LCON(K) + KCON(J, J1)$,
go to STEP 17

(b) $LCON(K) = LCON(K) - KCON(J, J1)$,
go to STEP 17

STEP 18: $K = K + 1$

If $K \leq LL$ set $I = 0$, go to STEP 17

Find K such that

$M = LCON(K) = \text{Min}\{LCON(K), K=LI, LL\}$

STEP 19: Set $L = 0$

For $K = LI, LL$

If $LCON(K) = M$ Set $L = L + 1$

STEP 20: If $L = 1$ go to STEP 21

Find K such that

$N(K) = \text{Min}\{N(K)/LCON(K) = M\}$

STEP 21: Set

$N(K) = N(K) + 1$
 $IKP(K, N(K)) = J$
 $IKC(K, N(K)) = ICRS(J)$
 $IKG(K, N(K)) = IG(J)$
 $MSEM(K, N(K)) = ISEM(J)$
 $IKF(K, N(K)) = 0$

Go to STEP 15

Scheduling Multi-Section Courses
(2 to 5 Sections)

STEP 22: Set $J = J + 1$, $LI = 1$, $LL = 6$

$NCAN(KT) = 0$, $LL2(KT) = 0$, $KT = LI$, LL

If $MSEC(J) > 5$ the course offers 6 sections or more,
go to STEP 37

If $ITIME1(J) = 0$ There is no time restriction,
go to STEP 23

$LI = ITIME2(J)$

$LL = ITIME3(J)$

STEP 23: If $ITIME(J) = 0$ there is no course restriction,
go to STEP 24

Otherwise, there are course restrictions

Restrictions are $(NCR(K), K=1, ITIME4(J))$, set $K = 0$

(a) Set $K = K + 1$

for time periods LI to LL find that time period

KT for which $NCR(K) = IKC(KT, I), I=1, N(KT)$

Then, set $KG(K) = KT$

If $K \leq ITIME4$ go to STEP 23(a)

(b) Set $NP = 0$

for $KT = LI$ to LL

Set $NP = NP + 1$

If $KG(K) \neq KT$ for $K = 1, ITIME4$,

set $KTN(NP) = KT$

STEP 24: Set $N1 = 0$

(a) $N1 = N1 + 1$, if $N1 > K1$ go to STEP 28

(Consider student $N1$ in the list of those
students requesting the present course.)

IDS = IST1(N1) ID of student N1,
 (JCRS(L), KSEC(L), L=1, NSCO) Courses
 requested by student N1 and the corresponding
 sections to which he is scheduled, if any.

K = 0

(b) Set K = K + 1 KT = KTN(K)

If K > NP go to STEP 24(a)

IN = 0

(c) IN = IN + 1

If IN > N(KT) go to STEP 27

If there exists an L such that {IKC(KT,IN) =
 JCRS(L), L=1, NSCO} go to STEP 25

Go to STEP 24(c)

STEP 25: J1 = IKP(KT,IN)

If MSEC(J1) = 1 The student is scheduled into a
 single section course at time
 period KT. Go to STEP 26

If MSEC(J1) ≠ 1 Then KSEC(L) indicates the section
 in which the student is scheduled.

If IKF(KT,IN) = KSEC(L) The student is scheduled
 at a section of a course previously
 scheduled at time period KT,
 go to STEP 26

If IKF(KT,IN) ≠ KSEC(L) go to STEP 24(c)

STEP 26: ISEM = ISEM(J) * MSEM(KT,IN)

If ISEM = 2 The present course and J1 are supposed
 to be scheduled at different semesters

L12(KT) = L12(KT) + 1

ID12(KT, L12(KT)) = IDS

Go to STEP 24(c)

If ISEM $\neq$ 2 Go to STEP 24(b)

STEP 27: NCAN(KT) = NCAN(KT) + 1

KS = NCAN(KT)

KCRS(KT,KS) = IDS

Go to STEP 24(b)

STEP 28: K1 = ITST(J)

MO = K1/MSEC(J)

NO = PECENT*MO

NP = 0

KT = LI

(a) If NCAN(KT) < NO go to STEP 28(b)

NP = NP + 1

KTN(NP) = KT

(b) KT = KT + 1

If KT $\leq$ LL go to STEP 28(a)

STEP 29: NP = No. of feasible time periods

M = MSEC(J) No. of sections offered by the course.

Possible assignment combinations = C_M^{NP} =

Alternative choices

JS = 0, IT = 0, ITC(IT) = 0, IT = 1, ..., C_M^{NP}

(a) NK(KT) = 0, KT = LI, LL, IT = IT + 1

(b) JS = JS + 1

KT = JT(IT,JS)

ITC(IT) = ITC(IT) - L12(KT)

If JS < M go to STEP 29(b)

L1 = 0

STEP 30: JS = 1

L1 = L1 + 1

If L1 > K1 go to STEP 33

(a) KT = JT(IT,JS)

(b) If NK(KT) > NCAN(KT) go to STEP 31

If IST1(L1) > KCRS(KT,NK(KT)) go to STEP 32

If IST1(L1) < KCRS(KT,NK(KT)) go to STEP 31

NK(KT) = NK(KT) + 1

Go to STEP 30

STEP 31: JS = JS + 1

If JS ≤ M go to STEP 30(a)

ITC(IT) = ITC(IT) + 1

Go to STEP 30

STEP 32: NK(KT) = NK(KT) + 1

Go to STEP 30(b)

STEP 33: If $IT \leq C_M^{NP}$ go to STEP 29(a)

Find IT such that

$ITC(IT) = \min\{ITC(IT), IT=1, \dots, C_M^{NP}\}$

Then JT(IT,JS), JS=1,M is the best alternative choice

Set NST(K) = 0, NK(K) = 1, K = 1, ..., 6, L1 = 0

STEP 34: ITC(J) = 0

(a) L1 = L1 + 1

If L1 > K1 go to STEP 36

IDS1 = IST1(L1)

If ISEM(J) = 0 the course is all-year,

go to STEP 35

Set JS = 1

(b) $KT = JT(IT, JS)$

If $IDS1 = ID12(KT, K)$, $K=1, L12(KT)$,

go to STEP 35(b)

JS = JS + 1

If $JS \leq M$ go to STEP 34(b)

STEP 35: Consider $S = 1, 2, \dots, M$

(a) Find JS such that $NST(JS) = \min\{NST(IS), IS \in S\}$

Then $KT = JT(IT, JS)$

(b) If $IDS1 = KCRS(KT, IL)$, $IL = 1, \dots, NCAN(KT)$,

then $NST(JS) = NST(JS) + 1$,

$MCRS(JS, NST(JS)) = IDS1$

Find L such that $ICRS(J) = JCRS(L)$, $L = 1, NSCO$

Set $KSEC(L) = JS$, go to STEP 34(a)

If $IDS1 \neq KCRS(KT, IL)$, $IL = 1, \dots, NCAN(KT)$,

then $S = S - JS$

If $S = \emptyset$ Go to STEP 35(c)

Go to STEP 35(a)

(c) $ITC(J) = ITC(J) + 1$

Go to STEP 34(a)

STEP 36: JS = 0

(a) JS = JS + 1

$KT = JT(IT, JS)$

$N(KT) = N(KT) + 1$

$IKP(KT, N(KT)) = J$

$IKC(KT, N(KT)) = ICRS(J)$

$IKG(KT, N(KT)) = IG(J)$

IKF(KT,N(KT)) = JS

MSEM(KT,N(KT)) = ISEM(J)

If JS < M Go to STEP 36(a)

Go to STEP 23

Scheduling Multi-Section Courses
(6 Sections or More)

STEP 37: Set J = J + 1, LI = 1, LL = 6, NCAN(KT) = 0,

L12(KT) = 0, KT = LI,LL

If J > NOC end of scheduling, go to STEP 49

If ITIME1(J) = 0 there is no time restriction,
go to STEP 38.

LI = ITIME2(J)

LL = ITIME3(J)

STEP 38: If ITIME(J) = 0 there is no course restriction,
go to STEP 39

There are course restrictions given by:

(NCR(K), K = 1,ITIME4(J))

Set K = 0

(a) K = K + 1

for time periods KT = LI to LL

If NCR(K) = IKC(KT,I), I = 1,N(KT) Set KG(K) = KT

If K ≤ ITIME4(J) go to STEP 38(a)

STEP 39: Set N1 = 0

(a) N1 = N1 + 1

Consider N1th student in the list of those
requesting the present course

If N1 > K1 Go to STEP 43

IDS = IST1(N1) ID of N1th student

JCRS(L),KSEC(L),L=1,NSCO Course requests and
sections of courses to which N1th student is
scheduled

IN = 0

KT = LI - 1

(b) KT = KT + 1

(c) Set IN = IN + 1 If KT > LL go to STEP 39(a)

If IN > N(KT) go to STEP 42

If there exists an L such that

{IKC(KT,IN) = JCRS(L), L=1, NSCO} Go to STEP 40

Go to STEP 39(c)

STEP 40: J1 = IKP(KT,IN)

If MSEC(J1) = 1 the student is scheduled into a
single-section course at time KT.
Go to STEP 41

If MSEC(J1) ≠ 1 then KSEC(L) indicates the section
in which the student is scheduled

If IKF(KT,IN) = KSEC(L) the student is scheduled at
a section of a course previously
scheduled at time period KT.
Go to STEP 41

If IKF(KT,IN) ≠ KSEC(L), go to STEP 39(c)

STEP 41: ISEM = ISEM(J)*MSEM(KT,IN)

If ISEM = 2 The present course and J1 are supposed
to be scheduled at different semesters.

L12(KT) = L12(KT) + 1

ID12(KT,L12(KT)) = IDS

Go to STEP 39(c)

If ISEM ≠ 2 Go to STEP 39(b)

STEP 42: $\text{NCAN}(\text{KT}) = \text{NCAN}(\text{KT}) + 1$

$\text{KS} = \text{NCAN}(\text{KT})$

$\text{KCRS}(\text{KT}, \text{KS}) = \text{IDS}$

Go to STEP 39(b)

STEP 43: $\text{K1} = \text{ITST}(\text{J})$

$\text{MO} = \text{K1}/\text{MSEC}(\text{J})$

$\text{NO} = \text{PECENT} * \text{MO}$

If $\text{ITIME} = 0$ Go to STEP 44

$\text{MCAN}(\text{KT}) = \text{NCAN}(\text{KT}), \text{KT} = \text{LI}, \text{LL}$

$\text{KT} = \text{LI}$

(a) If $\text{KG}(\text{K}) = \text{KT}$ for $\text{K} = 1, \text{ITIME4}(\text{J})$,

set $\text{NCAN}(\text{KT}) = \text{NO}$, If $\text{NCAN}(\text{KT}) > \text{NO}$

Otherwise $\text{KT} = \text{KT} + 1$

If $\text{KT} \leq \text{LL}$ go to STEP 43(a)

$\text{KT} = \text{LI} - 1, \text{IA} = 0$

(b) $\text{KT} = \text{KT} + 1$

If $\text{NCAN}(\text{KT}) \geq \text{NO}$ $\text{IA} = \text{IA} + 1$ $\text{KTN}(\text{IA}) = \text{KT}$

If $\text{KT} \leq \text{LL}$ go to STEP 43(b)

$\text{LA} = 0$

Set $\text{NO2} = \text{NO}$

STEP 44: If $\text{MSEC}(\text{J}) \leq \text{IA}$ go to STEP 46

$\text{LA} = \text{LA} + 1$

$\text{KEEP}(\text{LA}) = \text{IA}$

$\text{NO2} = \text{NO2} + \text{NO}, \text{IB} = 0, \text{KT} = \text{LI} - 1$

(a) $\text{KT} = \text{KT} + 1$

If $\text{NCAN}(\text{KT}) \geq \text{NO2}$ $\text{IA} = \text{IA} + 1$ $\text{IB} = \text{IB} + 1$
 $\text{KTN}(\text{IA}) = \text{KT}$

If $KT < LL$ go to STEP 44(a)

If $IB = 0$ go to STEP 45

Go to STEP 44

STEP 45: $AO = NO$

$NO = .9 * AO$ (PERCNT = .9)

$KT = LI - 1$, $IA = 0$ go to STEP 43(b)

STEP 46: If $MSEC(J) = IA$ go to STEP 47

$K = KEEP(LA) + 1$

$I = KTN(K), \dots, KTN(IA)$

(a) Find KT such that

$N(KT) = \text{Min}\{N(I), I=KTN(K), \dots, KTN(IA)\}$

Set $KTN(K) = KT$

$K = K + 1$

If $K > MSEC(J)$ go to STEP 11

$I = \{KTN(K), \dots, KTN(IA)\} - KTN(KT)$

Go to STEP 46(a)

STEP 47: $L1 = 0$, $M = MSEC(J)$, ($NST(JS) = 0$, $JS = 1, \dots, M$)

(a) $L1 = L1 + 1$

If $L1 > K1$ go to STEP 48

$IDS = IST1(L1)$

Consider $S = 1, 2, \dots, M$

(b) Find JS such that $NST(JS) = \text{Min}\{NST(IS), IS \in S\}$,

then $KT = KTN(JS)$

If $IDS = KCRS(KT, IL)$, $IL = 1, \dots, NCAN(KT)$,

then $NST(JS) = NST(JS) + 1$

$MCRS(JS, NST(JS)) = IDS$

Find L such that $ICRS(J) = JCRS(L)$, $L = 1, NSCO$

Set $KSEC(L) = JS$ go to STEP 47(a)

If $IDS \neq KCRS(KT, IL)$, $IL = 1, \dots, NCAN(KT)$,

Then $S = S - JS$

If $S = \phi$ go to STEP 47(c)

Go to STEP 47(b)

(c) $ITC(J) = ITC(J) + 1$

Go to STEP 47(a)

STEP 48: $JS = 0$

(a) $JS = JS + 1$

$KT = KTN(JS)$

$N(KT) = N(KT) + 1$

$IKP(KT, N(KT)) = J$

$IKC(KT, N(KT)) = ICRS(J)$

$IKG(KT, N(KT)) = IG(J)$

$IKF(KT, N(KT)) = JS$

$MSEM(KT, N(KT)) = ISEM(J)$

If $JS < M$ go to STEP 48(a)

Go to STEP 37

STEP 49: End of scheduling process. All courses have been considered and scheduled.

4.4. Implementation

All the algorithms for the heuristic approach have been implemented on the IBM 370/158 computer of the University of Oklahoma.

All the programs are written in FORTRAN and compiled with IBM FORTRAN IV level G compiler.

As indicated earlier the overall solution approach could be considered in several independent modules, yet related to each other in the sense that each module or component is a logical succession of previous segments.

Therefore, an overall program is developed which is composed of several components, mainly based on the type of courses considered for scheduling, as follows:

Component 1: Program MASTERFILE

Component 2: Program SCHEDULE-1

Component 3: Program SCHEDULE-2

Component 4: Program SCHEDULE-3

To run the overall program the above components must be run successively in the order given above. The functions of each component is discussed subsequently.

4.4.1. Component 1: Program MASTERFILE

Program MASTERFILE originally sets up several files on an auxiliary disk as follows:

1. CSF: Course specification file
2. SRF: Student request file
3. CRF: Course request file
4. NRF: New request file
5. HLD: Hold intermediate results file
6. SWP: Swap intermediate results file

Information concerning courses to be scheduled are read in by this program from cards and stored in course specification file (CSF). This input information is punched on cards one card for each course as follows:

	<u>Columns</u>	<u>Format</u>
ICRS(J) = Course ID number: (IG(J), Type A or B or blank)	1-5	I4,A1
MSEC(J) = Number of sections offered by the course	7-8	I2
NJ = Maximum number of students allowed in a section	9-12	I4
ITYP(J) = Course type	13	I1
ISEM(J) = Semester offered	14	I1
ITIME(J) = Course restrictions	15	I1
ITIME1(J) = Time restrictions	16	I1
ITIME2(J) = If time restricted, beginning time	17	I1
ITIME3(J) = If time restricted, ending time	18	I1
ITIME4(J) = If course restricted, number of sections	19	I1
NCR1,NG1 = First course restriction	20-24	I4,A1
NCR2,NG2 = Second course restriction	25-29	I4,A1
NCR3,NG3 = Third course restriction	30-34	I4,A1

Conventions for the use of the above variables are as follows:

$$ITYP(J) = \begin{cases} 0 & \text{for locked in courses 2, 3 or 4} \\ & \text{for several-hour courses (based on number of hours)} \\ 1 & \text{for other types of courses} \end{cases}$$

$$ISEM(J) = \begin{cases} 0 & \text{for all-year courses} \\ 1 & \text{for first-semester courses} \\ 2 & \text{for second semester courses} \end{cases}$$

ITIME(J) = $\begin{cases} 0 & \text{if there is no course restriction} \\ 1 & \text{if there are course restrictions} \end{cases}$

ITIME1(J) = $\begin{cases} 0 & \text{if there is no time restriction} \\ 1 & \text{if there are time restrictions} \end{cases}$

ITIME2(J) = If ITIME1(J) = 1 then ITIME2(J) is the beginning time period and ITIME3(J) is the ending time periods at which course ICRS(J) is allowed to be scheduled

ITIME4(J) = If ITIME(J) = 1 (course restrictions), ITIME4(J) is the number of sections and NC1, NC2, NC3, are the corresponding courses by which ICRS(J) is restricted.

IG(J), NG1, NG2, and NG3 could be characters A, B, or blank.

Information concerning students are stored in student request file (SRF) and they are read in from cards one card for each student according to their ascending ID number as follows:

	<u>Columns</u>
I = Student identification number	1-6
NAME(L) = Student's name	7-36
IGRD = Student's grade (9th grade, 10th grade, etc.)	37-38
ISEX = Student's sex (male, female)	39
NSCO = Number of courses requested by the student	40
JSCR(L), = List of courses requested by the student (L = 1, NSCO)	41-76

The information about courses, stored in course specification file, and about students stored in student request file is used by program MASTERFILE to generate another file called Course Request File (CRF) which includes,

for each course, information about the course, and a list of ID of those students who are requesting the course. This program computes an index of difficulty for each course in the single-section and multi-section category and then arranges all the courses and their corresponding information according to the descending order of the index of difficulty in a new file called New Request File (NRF). The indices of difficulty for different groups of courses are computed as explained in previous sections. The two files HLD and SWP are set up by program MASTERFILE, but they are used by other program components when multi-section courses are being scheduled.

Information about courses and their corresponding request lists are printed such that all the information such as course number, number of sections, restrictions, index of difficulty, etc., are printed first and then the ID of all the students requesting the course are printed.

4.4.2. Component 2: Program SCHEDULE-1

SCHEDULE-1 computes a conflict matrix for all the locked-in, several-hour, single-section restricted, and single-section unrestricted courses. The conflict matrix is printed out as follows:

```
COURSE 1 (529) AND COURSE 2 (538) NUMBER OF CONFLICTS = 0
:
COURSE 1 (529) AND COURSE 7 (570) NUMBER OF CONFLICTS = 1
:
```

This means that no students have requested course 529 and 538 at the same time, i.e., if these two courses are scheduled at the same time no conflicts would result. On the other hand, if course 529 and 570 are scheduled at the same time period, one student would encounter a conflict, and so on.

SCHEDULE-1 considers locked-in courses first and after all the locked-in courses are scheduled, several-hour courses, single-section restricted, and single-section unrestricted courses are scheduled, respectively. For each of the above course categories the corresponding algorithm explained in previous sections is implemented.

Information concerning courses, the time periods at which they are scheduled, conflicts, ID of students encountering conflicts, etc. is printed out and a part of this information is stored in CSF to be used, in the future, when other program components are used to schedule multi-section courses. This program also prints out a table showing the time periods and the courses scheduled at their corresponding time periods. The input data for program SCHEDULE-1 is already stored in one of several files mentioned earlier and, therefore, there is no new input data to be read in by this program.

4.4.3. Component 3: Program SCHEDULE-2

SCHEDULE-2 is designed for scheduling multi-section courses which offer less than 6 sections employing heuristics

and models discussed in earlier sections for these types of courses. The input data to this program comes from the information stored in various files set up by program MASTERFILE and computed partially, in some cases, by program SCHEDULE-1.

Information about courses, the corresponding sections, time periods at which each section is scheduled, ID of students scheduled into each section, etc., is printed out and a part of this information is stored in CSF and CRF to be used, in the future, when other multi-section courses and other program components are considered. Finally, this program also prints out a table showing the time periods and the courses and sections of courses and their corresponding time periods.

4.4.4. Component 4: Program SCHEDULE-3

This program is designed for scheduling courses which offer 6 sections or more employing heuristics and procedures discussed earlier for these types of courses. The input-output operations of this program are very similar to that of SCHEDULE-2.

4.4.5. Setting Up the Program and Data Deck

Data deck for program MASTERFILE is shown in Figure 8.

Other programs, i.e., SCHEDULE-1, SCHEDULE-2, and SCHEDULE-3 do not require input information from cards and they are simply run, one after the other, using proper control cards.

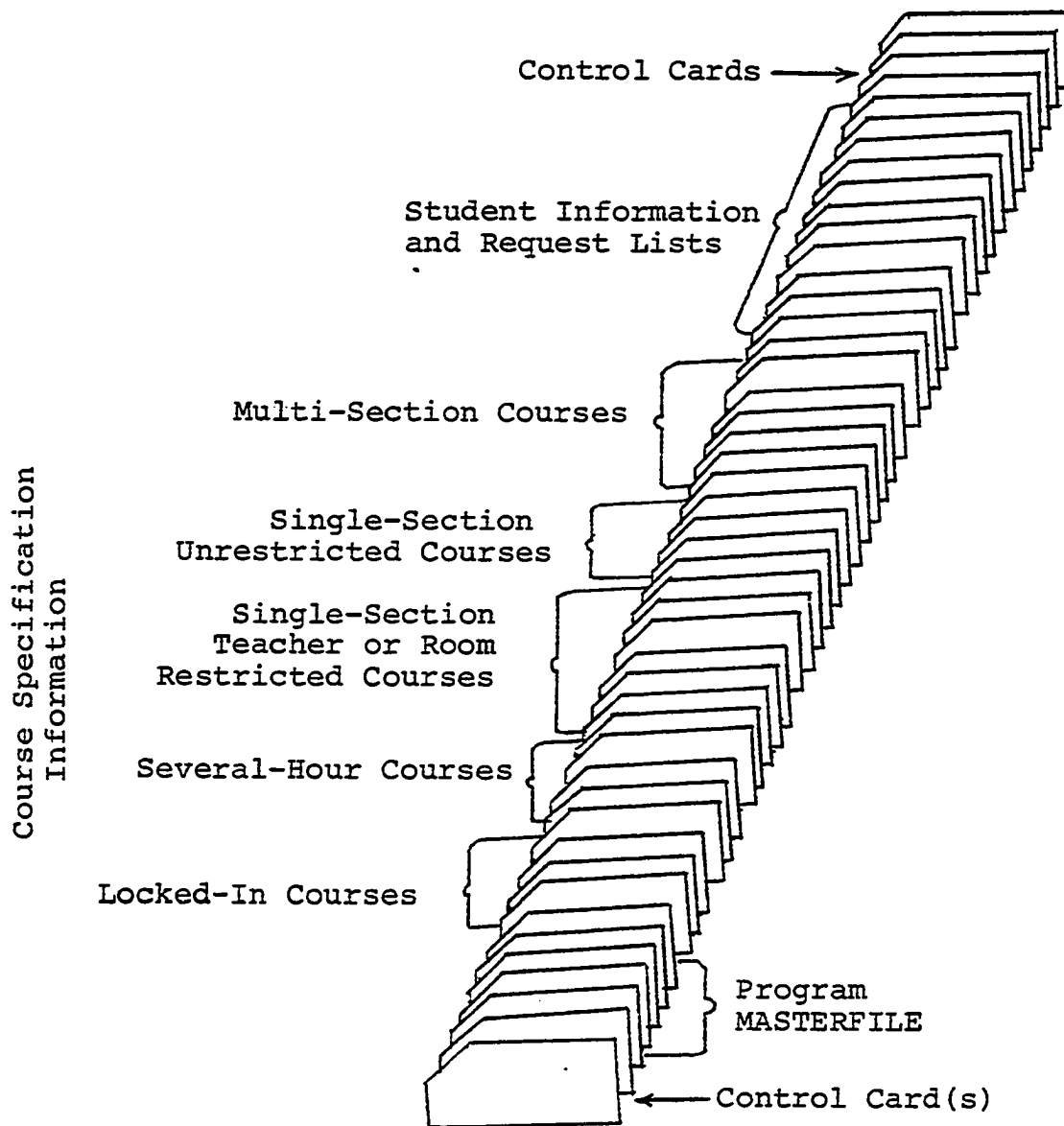


Figure 8.

4.5. Computational Experience

The algorithms described in Chapter IV and incorporated in a program were tested on a data set obtained from Central Mid-High School, in Norman, Oklahoma. This school was selected as a typical example because of its relatively complex school structure. The school, unfortunately, has no schedules generated by other systems and therefore it is not possible to compare this system with

other systems. However, a careful manual class schedule has been constructed by the school principal for the data set mentioned above.

The school had 714 students and 68 courses with a total of 157 sections. There were 11 locked-in courses, 4 several-hour courses (4-hours and 3-hours length), 6 single-section teacher or room restricted courses, 7 single-section unrestricted courses, and all the rest (40 courses) were multi-section courses offering between 2 and 10 sections. Some of the multi-section courses were restricted by other multi-section courses and/or other single-section courses. There were 10 first-semester courses (single-section and multi-section), 9 second-semester courses (single-section and multi-section), and the rest of the courses were of all-year type.

The class schedule constructed by the heuristic algorithms generated student conflicts less than one-half student conflicts resulting from the manual class schedule as shown in Appendices II and III. The number of conflicts generated by scheduling different types of courses using the heuristic approach and total number of conflicts for the second run of the manual schedule are shown below:

Course Type	Number of Courses	Total Number of Conflicts
Locked-in	11	54
Several-hour	4	
Teacher or room restricted	6	
Single section unrestricted	7	
Multi-section (2-5 sections)	27	84
Multi-section (6 sections and more)	7	47
Total number of conflicts		185
Manual schedule second run,* Total number of conflicts		359

*The corresponding first run manual schedule generated over 400 conflicts.

There was a good degree of balance between different sections of most courses (as can be seen in the sample output part) although a measure of balance or imbalance was not calculated.

The relevant cost factors for the system IBM 370/158 computer of the University of Oklahoma are shown below:

CPU time	\$.075/sec.
Storage residency	\$.003/1K-min.
Cards read	\$.75/1000 cards
Lines printed	\$.85/1000 lines
Disk I/O	\$.25/1000 Excp

Different components of the overall program were quite fast as far as CPU time is concerned. For one clean run from original input data to final complete class schedule the following information was obtained.

<u>Program</u>	<u>C.P.U. Time in Seconds</u>	<u>C.P.U. Cost Dollars</u>	<u>Total Cost Dollars</u>
MASTERFILE	25	1.90	4.50
SCHEDULE-1	21	1.56	3.40
SCHEDULE-2	193	14.50	18.89
SCHEDULE-3	<u>66</u>	<u>4.94</u>	<u>7.04</u>
Total	305	22.90	33.83

The total cost column includes the sum of CPU, storage residency, cards read, lines printed, and Disk I/O costs for each component program. Considering the above information it is seen that:

CPU time per student = $305/714 = .4$ sec.

CPU cost per student = $\$.075 \times .4 = \$.03 = 3$ cents

Total cost per student = $33.83/714 = \$.046 = 4.6$ cents

As indicated above, the costs are for one successful run of each program component. It is, however, unlikely that a single first run of each component would be successful, especially for program MASTERFILE where original errors in the input data could easily require several runs of this program. The comparison of costs between this system and others is not easily possible, since there is no way of

knowing how much it would cost to perform class scheduling for this case using other systems.

The costs per student figures, nonetheless, could be roughly compared to those of other programs or systems such as IBM's CLASS (Class Load And Student Scheduling) and ECSS (Educational Coordinates Scheduling System) which is a much improved variation of the SSSS (Stanford School Scheduling System). In the case of IBM CLASS [21] computer costs ran approximately 35 cents per student for performing scheduling for schools in the Cleveland, Ohio-Pittsburgh, Pennsylvania area. Salazar [37] indicates that traditional school scheduling work (one pass from data to final schedule, excluding final printouts) using the ECSS system represents computer expenditures of 25 to 30 cents per student. The comparison shows that the present system requires costs which are almost one tenth of those encountered when using such systems as IBM CLASS and ECSS.

Finally, from the experience of the present system implementation, it is the author's belief that the overall system composed of several components, processed successively and not at the same time, is suitable for implementation on small computers. The main reason being that each of the programs MASTERFILE, SCHEDULE-1, SCHEDULE-2, and SCHEDULE-3 are relatively small size programs. The sizes of these programs for runs using Central Mid-High data, as indicated by FORTRAN IV level G compiler is summarized below:

<u>Program</u>	<u>Size in K-bytes</u>
MASTERFILE	15.5
SCHEDULE-1	23.3
SCHEDULE-2	22.2
SCHEDULE-3	20.0

Sacrifices, however, would perhaps be necessary to make such as giving up CPU speed for lower memory requirements or reducing a component's size by eliminating some of its non-essential options. To keep main memory requirements relatively low, auxiliary storage such as magnetic disks for setting up different files and storing intermediate and final results would be necessary.

4.6. Summary

In this chapter a heuristic scheduling algorithm was developed and coded in FORTRAN which is suitable for generating class schedules for traditional and modified school structures.

The overall algorithm was composed of several components. Each component incorporated heuristics, mathematical optimization techniques, or a combination of both and was specifically developed for scheduling a specific type of course being considered.

For scheduling groups of teacher or room restricted courses, an assignment problem was constructed for each group of courses restricted by a teacher or room. The

solution to this assignment problem was used to schedule the group of restricted courses simultaneously and optimally with respect to all other courses already scheduled.

For multi-section courses an integer programming model was presented, and an enumerative procedure used to schedule a multi-section course optimally with respect to all other courses already scheduled.

The heuristic algorithm was implemented using actual data from a local high school. Computational results indicated that the algorithm was fast, generated relatively few conflicts, and computer memory requirements were low. Computational results were reported and compared to a manually constructed schedule using the same set of data.

CHAPTER V

CONCLUSIONS AND SUGGESTIONS FOR FURTHER RESEARCH

One important characteristic of the overall algorithm is its relatively low main memory requirements and use of auxiliary memory for the storage of data. The author believes that the algorithm could be implemented on a small computer installation, such as those often available to many secondary schools. However, some compromises would perhaps have to be made, such as the trade-off between speed and memory requirements or, in other words, sacrificing computers speed for lower memory requirements.

5.1. Conclusions

Considering the theoretical model for the problem of scheduling students into an existing class schedule it is seen that:

- (a) The model based on network flow concepts could be used to obtain an overall optimum solution by formulating and solving a small network flow problem for each student.
- (b) The model is flexible in the sense that different

types of courses, time periods, cycles, etc., could be incorporated in it.

(c) Other applications of the model could be considered such as resolving student conflicts, evaluating a heuristic procedure, evaluating a class schedule constructed by a heuristic procedure, etc.

(d) The model could be quite practical since there are extremely efficient techniques for solution of network flow problems of the type encountered here.

Considering the heuristic system, designed for traditional and modified school structures, developed, implemented, and tested, it is concluded that:

(a) The system is composed of several components each designed for scheduling one or several special types of courses. The courses are considered for scheduling successively and each course, or group of courses, theoretically, is scheduled optimally with respect to all other courses which are already scheduled.

(b) The system is quite fast in terms of computer cpu time and therefore economical in terms of costs. In the case of a complex schedule constructed for a local school cpu times of .4 seconds per student and total costs of approximately 5 cents per student were observed.

(c) The system is very flexible in the sense that it is capable of handling several types of courses and more

importantly could be revised and improved to make it able to handle still more complex situations in terms of courses, time periods, cycles, etc.

- (d) The system is efficient in terms of number of student conflicts generated. In the case of class schedule mentioned in part (b) above the number of student conflicts generated using this system was less than one-half the number of student conflicts generated by a carefully manually constructed class schedule for the same data.
- (e) In most cases there was a good degree of balance between different sections of each course for Central Mid-High final class schedule.
- (f) Since the system is composed of several small components it is suitable for implementation on small computers which are often available to secondary schools.

5.2. Suggestions for Further Research

Further research is recommended in three main areas. These include theoretical investigations, improvements on the existing programs and algorithms, and development of new programs and algorithms. More specifically, the following areas are suggested for further research:

- (a) Development of models similar to those investigated in this research, possibly using network flow theory, for other types of scheduling problems such

as multi-section scheduling, etc.

- (b) Design, development, implementation and testing of algorithms and programs based on the network flow model discussed here to schedule students into an existing class schedule, resolve conflicts generated by a heuristic approach, and evaluate a heuristic approach or its corresponding class schedule.
- (c) Computational experience on several other data sets from different schools is highly recommended to gain insight into weak and strong points of the present system and obtain more reliable information such as speed, efficiency, flexibility, section balancing, etc.
- (d) Improve the present system for reduced cpu time (or cost) and memory requirements possibly by program modifications to represent arrays in bit strings for more efficiency. This may be achieved by having parts of different programs written in assembler language. Other improvements include allowing for computation of a measure of balance or imbalance between different sections of a course, refined output information, etc.
- (e) Revise the present system to allow for more scheduling flexibility such as use of weekly or other types of cycles in addition to daily cycles, use of conflicting time periods, etc. It should be pointed out

that the present system allows for these types of revisions and modifications and all the algorithms utilized by the present system could easily be modified and used in more complex situations indicated above.

- (f) Finally, assuming a class schedule is constructed using a heuristic technique, possibilities of improvements on the number of conflicts for this class schedule should be investigated. These types of improvements could be obtained in one of two ways: One way is to attempt to find feasible schedules for some of the students who have encountered a conflict. This is the problem of finding a feasible schedule for a set of students when a class schedule exists. It was shown, in previous chapters, that this problem could be solved optimally using a network flow model. The second way of improving on a present schedule could be to deschedule and reschedule one or a set of courses which produce a reduction on the total number of students encountering conflicts. In this case existing methods such as piecewise interchange or other optimization methods could be promising.

REFERENCES

1. Gilmore, P. C. "Optimal and Suboptimal Algorithms for the Quadratic Assignment Problem." SIAM Journal. Vol. 10, No. 2. 1962.
2. Land, A. H. "A Problem of Assignment with Inter-related Costs." Operational Research Quarterly. Vol. 14. 1963.
3. Lawler, Eugene. "The Quadratic Assignment Problem." Management Science. Vol. 9, No. 4. July, 1963.
4. Laughhunn, D. J. "Quadratic Binary Programming with Application to Capital-Budgeting Problems." Journal of the Operations Research Society of America. Vol. 18, No. 3. May-June, 1970.
5. Carlson, R. C., and Nemhauser, G. L. "Scheduling to Minimize Interaction Cost." Journal of the Operations Research Society of America. Vol. 14, No. 1. January-February, 1966.
6. Broder, S. "Final Examination Scheduling." Communications of the ACM. Vol. 7, No. 8. August, 1964. pp. 494-498.
7. Davis, F. E., Devine, M. D. and Lutz, R. P. "Scheduling Activities among Conflicting Facilities to Minimize Conflict Cost." Mathematical Programming. Vol. 6, No. 2. April, 1974.
8. Hall, A. D., Jr., Acton, F. S. "Scheduling University Course Examination by Computer." Communications of the ACM. Vol. 10, No. 4. April, 1967. pp. 235-238.
9. Weil, Roman L. and Kettler, Paul C. "Rearranging Matrices to Block-Angular Form for Decomposition (and Other) Algorithms." Management Science. Vol. 18, No. 1. September, 1971.

10. Gotlieb, C. C. "The Construction of Class-Teacher Time-Tables." Proceedings of the International Federation for Information Processing Congress 62, Munich. Published by North-Holland Publishing Company, Amsterdam, 1963. pp. 73-77.
11. Dempster, M. A. H. "Two Algorithms for the Time-Table Problem." Reprinted from D. G. A Wish, Col: Combinational Mathematics and the Applications. Academic (1971).
12. Lions, J. "Matrix Reduction Using the Hungarian Method for the Generation of School Timetables." Communications of the ACM. Vol. 9, No. 5. May, 1966. pp. 349-354.
13. Almond, M. "An Algorithm for Constructing University Timetables." The Computer Journal. January 8, 1966. pp. 331-340.
14. Berghuis, J., Van Der Heiden, A. J. and Bakner, R. "The Preparation of School Timetables by Electronic Computer." BIT, Norwegian Journal of Information Processing. No. 4. 1964.
15. Barraclough, E. D. "Application of a Digital Computer to the Construction of Timetables." The Computer Journal. July 8, 1965. pp. 136-46.
16. Balas, Egon. "An Additive Algorithm for Solving Linear Programs with Zero-One Variables." Opr. Res. Vol. 13, No. 4. pp. 517-546.
17. Glover, Fred. "A Multiphase-Dual Algorithm for the Zero-One Integer Programming Problem." Operations Research. Vol. 13, No. 6. 1965.
18. Petersen, Clifford C. "Computational Experience with Variants of the Balas Algorithm Applied to the Selection of R&D Projects." Management Science. Vol. 13, No. 9. May, 1967.
19. Geoffrion, A. M. and Marsten, R. E. "Integer Programming Algorithms: A Framework and State-of-the-Art Survey." Management Science. Vol. 18, No. 9. May, 1972.
20. Holz, R. E. "Computer-Assisted Scheduling." Journal of Educational Data Processing. Vol. 7, No. 2. Spring, May 1964. pp. 36-41.
21. Smith, E. T. "The Computer in Student Scheduling (IBM "CLASS")." Journal of Educational Data Processing. Vol. 1, No. 2. Spring, May 1964. pp. 53-60.

22. Wilkes, C. F. "SOCRATES: A System for Student Scheduling." Journal of Educational Data Processing. Vol. 1, No. 2. Spring, May 1964. pp. 46-52.
23. Anderson, G. E., Jr. "An Algorithm for Assigning Pupils to Classes." Journal of Educational Data Processing. Vol. 1, No. 2. May 1964.
24. Oakford, Robert V., Allen, Dwight W. and Chatterton, Lynne A. "School Scheduling--Practice and Theory." Journal of Educational Data Processing. Vol. 4, No. 1. Winter 1966-67.
25. Van Dusseldrop, Ralph A. and Richardson, Duane E. "Class Scheduling by Computer." Journal of Educational Data Processing. Vol. 4, No. 1. Winter, 1966-67.
26. Faulkner, M. A. "Computer Sectioning and Class Scheduling for a University." Automated Educational System Enoch Haga (Ed.) The Business Press, Elmhurst, Illinois, 1967.
27. Busam, V. A. "An Algorithm for Class Scheduling with Section Preference." Communications of the ACM. Vol. 10, No. 9. September, 1967. pp. 567-569.
28. Macon, N. and Walker, E. E. "A Monte Carlo Algorithm for Assigning Students to Classes." Communications of the ACM. Vol. 9, No. 5. May, 1966. pp. 339-340.
29. Lutz, R. P., Devine, M. D., Kumin, H. L. and Smith, W. C. "An Application of Operations Research to School Desegregation." Management Science. Vol. 19, No. 4. December, Part 2, 1972.
30. Lawler, E. and Bell, M. "A Method for Solving Discrete Optimization Problems." Journal of the Operations Research Society of America. Vol. 14, No. 6. November-December, 1966.
31. Salazar, A. and Oakford, R. V. "A Graph Formulation of a School Scheduling Algorithm." Communications of the ACM. December, 1974. Vol. 17, No. 12.
32. Steen, L. A. "Computational Unsolvability." Science News 109 (1976), pp. 298-301.
33. Glover, F. and Klingman, D. "Real World Applications of Network Related Problems and Breakthroughs Is Solving Them Efficiently." ACM Transactions on Mathematical Software. Vol. 1, No. 1. March 1975. pp. 47-55.

34. Oakford, R. V. and Salazar, A. "Effective Use of the Conflict Matrix in Automated Scheduling." Journal of Educational Data Processing. Vol. 12, No. 1. 1975.
35. Holzman, G. A. "Master Time Schedule Generator--An Algorithm Approach." Paper presented at the Joint National Meeting, Operations Research Society of America and the Institute of Management Sciences, Las Vegas, November 18, 1975.
36. Heller, Nelson B. "Mathematical Methods of Examination Scheduling." Masters Thesis, M.I.T., Sept. 1963.
37. Salazar-Garcia A. "School Scheduling on an Interactive Computer System." Doctoral Dissertation, Department of Industrial Engineering, Stanford University, 1976.
38. Glover, F., Klingman, D. "Network Application in Industry and Government." AIIE Transactions, Vol. 9, No. 4. December, 1977.
39. Fulkerson, D. R. "An Out-of-Kilter Method for Solving Minimal Cost Flow Problems." J. Soc. Indust. Appl. Math, 9, 18-27 (1961).
40. Analysis, Research and Computation, Inc., Development and Computational Testing on a Capacitated Primal Simplex Transshipment Cole. ARC Technical Research Report, P.O. Box 4067, Austin, Texas 78765.
41. Barr, R. S., Glover, F., and Klingman, D. "An Improved Version of the Out-of-Kilter Method and a Comparative Study of Computer Codes." Mathematical Programming 7, 1, 60-87 (1974).
42. Clasen, R. J. The Numerical Solution of Network Problems Using the Out-of-Kilter Algorithm, Rand Corporation Memorandum RM-5456-PR, Santa Monica, California (1968).
43. Edmonds, J. and Karp, R. M. "Theoretical Improvements in Algorithmic Efficiency for Network Flow Problems," Journal of the Association for Computing Machinery, 19, 248-264 (1972).
44. Lutz, R. P.; Devine, M.D.; and Kumin, H. J. "Interactive Master Scheduling Program." Unpublished report, Department of Industrial Engineering, University of Oklahoma (1970).
45. Murikipudi, Madhavarao. "Computational Comparisons of Algorithms for Scheduling Activities among Conflicting Facilities to Minimize Conflict Cost." Unpublished Masters Thesis, Department of Industrial Engineering, University of Oklahoma (1975).

46. Hillier, F. S. and M. M. Connors. "Quadratic Assignment Problem Algorithms and the Location of Indivisible Facilities." Management Science. Vol. 13, No. 1. 1966.
47. Nugent, C. E., Vollmann, T. E., and Ruml, J. "An Experimental Comparison Technique for the Assignment of Facilities to Locations." Operations Research. Vol. 16, No. 1. 1968.
48. Pierce, J. F., and Crowston, W. B. "Tree-Search Algorithms for Quadratic Assignment Problems." Naval Research Logistics Quarterly. Vol. 18, No. 1. 1971.
49. Hanan, M., and Kurtzberg, J. "A Review of the Placement and Quadratic Assignment Problems." SIAM Review. Vol. 14. 1972.
50. Francis, R. L., and White, J. A. "Facility Layout and Location, An Analytical Approach." Prentice-Hall, 1974.
51. Geoffrion, A. M., and Graves, G. W. "Scheduling Parallel Production Lines with Changeover Costs: Practical Application of a Quadratic Assignment/LP Approach." Operations Research. Vol. 24, No. 4. July-Aug. 1976.
52. Garfinkel, R. S., and Nemhauser, G. L. "Integer Programming." John Wiley & Sons, 1972.
53. Karney, D., and Klingman, D. "Implementation and Computational Study on an In-Core, Out-of-Core Primal Network Code." Operations Research. Vol. 24, No. 6. Nov.-Dec. 1976.
54. Geoffrion, A. M. "A Guide Tour of Recent Practical Advances in Integer Linear Programming." SIAMP Newsletter, 17, 22-32, 1974.
55. Shih, W. "A Branch-and-Bound Method for the Multi-Constraint Zero-One Knapsack Problem." The Journal of Operational Research Society. Vol. 30, No. 4. 1979.
56. Talbot, F. B., and Patterson, J. H. "An Efficient Integer Programming Algorithm with Network Cuts for Solving Resource-Constrained Scheduling Problems." Management Science. Vol. 24, No. 11. 1978.
57. Zanakakis, S. H. "Heuristic 0-1 Linear Programming: An Experimental Comparison of Three Methods." Management Science. Vol. 24, No. 1. 1977.

58. Karp, R. M. "Complexity of Computer Computations."
Plenum Press, N.Y. pp. 85, 1972.

APPENDIX I

TERMINOLOGY

APPENDIX I

TERMINOLOGY

Master Schedule: Time schedule of courses

Conflict = student conflict: When two single-section courses requested by a student are scheduled at the same time period the student is encountering a conflict.

Time period = Time block: A time interval at which courses could be scheduled such as 8-9 A.M., 9-11 A.M., etc.

Conflicting Time Blocks: Time blocks or time periods which overlap such as 8-9 A.M., Monday and 8-11 A.M., Monday.

Cycle: A time interval during which all courses are repeated once, such as a day where all courses are repeated every day at the same time or a week where all the courses meet T Th or MWF at the same time every week.

Heuristic: An intuitive rule which works well in real life and is usually based on past experience. For example "more crowded single-section courses are harder to schedule than (and therefore should be scheduled before) less crowded single-section courses."

Index of Difficulty: A number assigned to each course based on some heuristics such as number of sections that the course is offering, number of students requesting the course, types of courses that these students are requesting, etc.

This number represents the relative difficulty of scheduling this course with respect to all other courses.

Algorithm: A step by step solution procedure for a problem, specially a mathematical problem, frequently involving repetition of one or more operations.

Student Request List: A list containing information about students and courses that they are requesting. For each student the list contains students, ID, name, grade, sex, and courses requested.

Course Specification List: A list containing information about courses offered. For each course the list contains course number, recommended number of sections, recommended index of difficulty, semester in which the course is offered, any course or time restrictions, etc.

Conflict Matrix:

1. Course Conflict Matrix: A matrix constructed by listing courses both row wise and column wise and each element of the matrix represents number of students which would encounter a conflict if the corresponding row and column courses are scheduled at the same time: example

	20	21	22
20	0	1	2
21	1	0	2
22	2	3	0

Suppose there are only three courses 20, 21, and 22. The conflict matrix above (considering the first row) indicates that if course 20 and course 21 are scheduled at the same time one student would encounter a conflict and if course 20 and 22 are scheduled at the same time two students would encounter conflicts and so on.

The conflict matrix would be constructed for single-section courses only.

2. Time Block Conflict Matrix: A matrix constructed by listing time blocks both row wise and column wise and each element of the matrix is a zero or one. A one indicates that the corresponding time blocks overlap and a zero indicates that the corresponding time blocks do not overlap. For example, consider:

Time Block	Time	
1	8-9	M.W.F.
2	9-10	M.W.F.
3	8-10	M.W.F.

then the conflict matrix is:

		Time Block		
		1	2	3
Time blocks	1	1	0	1
	2	0	1	1
	3	1	1	1

A "1" in row 2 column 3 means that time blocks 2 and 3 overlap. A "0" in row 2 column 1 means that time blocks 2 and 1 do not overlap.

APPENDIX II

SAMPLE MANUAL SCHEDULE, CENTRAL MID-HIGH SCHOOL

STUDENT COURSE CONFLICT LISTING 1/11/78

STUDENT NO.	NAME	CLASS	C O U R S E S	
			ACCEPTABLE NO SECT	UNACCEPTABLE NO SECT
ADA300	ADAMS, CINDY S	09	5441 50 5461 20 5801 10 5521 60 5331 60 5071 30	
AKR350	AKERMAN, RALPH T	09	5733 10 5931 60 5681 50 5531 60	
AKI300	AKIN, PATRICIA S	10	5611 60 5431 30 0121 60 5071 10 5691 50 5501 20	399140 537111
ALE350	ALEXANDER, LORNA	10	0001 30 5311 50 3091 31 5621 61 5651 41 5371 10	5581 20 5501 37 SECTION CANCELED

~~5501 37~~ SECTION CANCELED

~~5501 37~~ SECTION CANCELED

STUDENT COURSE CONFLICT LISTING 12/31/10

STUDENT NO. NAME CLASS ACCEPTABLE UNACCEPTABLE NO SECT NO SECT

ALE960 ALEXANDER, MICHAEL 09

5733 10 5701 40 5701 50 5521 60

AL1300 ALINGER, DANIEL 10

5691 30 3031 42 5691 10 5691 50 5531 60

ALL300 ALLEN, TONY M 09

5781 10 5591 60 5601 50 5571 20

ALL695 ALLISON, JALYN D 10

5713 10 5511 40 5541 50 5581 60

5801 31 5341 40
~~SECTION CONFLICT~~

0231 40
~~SECTION CONFLICT~~

5661 20

STUDENT COURSE CONFLICT LISTING 7/31/78

STUDENT NU.	NAME	CLASS	-----C O U R S E S-----		
			ACCEPTABLE NO SECT	UNACCEPTABLE NO SECT	
BAB600	BABICH, KRISTEN E	09	5733 10 5311 40		5311 60 5311 40 SECTION COMPLETE 0211 40 SECTION COMPLETE
			5521 50		
BAC100	BACHRACH, CASSIE L	10	5713 10 5391 50 0211 40 3091 40 5581 60		
BAC400	BACON, PEGGY S	09	5441 50 3241 30 0001 30 5011 40 5081 60		5801 21 5341 40 5341 40 SECTION COMPLETE
			5531 10		
BAC400	BACON, RANDALL D	10	5591 60 0241 30 3091 12 5671 50 5651 40 5581 20		3991 12 5321 11
					5311 40 SECTION COMPLETE

STUDENT COURSE CONFLICT LISTING 7731270

STUDENT NO.	NAME	CLASS	-----C O U R S E S-----	
			ACCEPTABLE NO SECT	UNACCEPTABLE NO SECT

BAR460	BARNARD, GERALD R	09	5411 00	5341 50
			5511 21	5411 00 UNACCEPTABLE
			5611 30	5491 10
			5711 40	5811 40
			5111 30	5681 60

BAR490	BARNARD, MARVIN W	10	5723 40
			5661 31
			5481 10
			3801 20
			0091 20

JAR630	BARNETT, ROBERT D	10	5701 10
			5723 40
			0601 20
			3091 20
			5671 30

BAR710	BARNHILL, RITA K	09	5801 30	5531 30
			5811 40	5311 60
			5311 21	5671 10
			5331 21	5671 00 UNACCEPTABLE
			5251 50	

STUDENT COURSE CONFLICT LISTING 7/31/78

STUDENT NO.	NAME	CLASS	-----C O U R S E S-----	
			ACCEPTABLE NO SECT	UNACCEPTABLE NO SECT
JUT260	BUTLER, DANNY J	10	5591 60 3051 10 0091 10 5691 40 5651 30 5581 20 5371 50	
BUT300	BUTLER, GLENNA J	10	5713 10 5321 40	
			5581 50	5831 77 SECTION COMPLETE 5001 04 SECTION COMPLETE 5461 60
JUT600	BUTLER, JINA M	09	5611 60 5611 40 5681 50	5801 31 5341 40 5311 77 SECTION COMPLETE 5461 20
			5531 10 5511 30	
BYR175	BYRD, BARBARA M	10	5791 10 5561 20 5311 30 5691 51 5651 60 5371 40	

STUDENT COURSE CONFLICT LISTING 11/17/78

STUDENT NAME CLASS ACCEPTABLE UNACCEPTABLE
 NO SECT NO SECT

CAN597 CANO, LISA L 09

5621 10 5811 50 5801 40 5571 21 5541 60 5311 30

CAP255 CAPE, CHESTER D 10

5591 60 0081 30 5521 10 5091 11 5691 40 5651 50

CAR370 CARPENTER, DAVID 10

5123 40 5081 10 5091 10 5081 10 5081 10

CAN587 CARSON, DAN 10

5751 50 5561 10 5511 30 5091 00 5091 11 5311 20

5591 60

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

5691 40

STUDENT COURSE CONFLICT LISTING 7/31/78

STUDENT NO.	NAME	CLASS	C O U R S E S	
			ACCEPTABLE NO SECT	UNACCEPTABLE NO SECT
CLE430	CLEAR, KEAL WAYNE	10	5591 61 5571 20 5561 10 5541 10 5531 21 5521 40	3791 10 0091 12 5561 30 5331 21 5801 40 5591 77 SECTION CONFLICT 5671 51
CL1570	CLINKENBEARD, LISA	10	5751 50 5713 10 5561 60 5501 40	Had Back in Summer. 5311 30 (out of town) 5511 77 SECTION CONFLICT
CLO700	CLOUSE, CYNTHIA L.	09	5311 60 5011 60 5601 50 5571 21 5311 10	5301 30 5341 50 5541 00 SECTION CONFLICT
CLO920	CLOUSE, DAVID A	09	5591 60 5301 30 5011 40 5601 50 5571 10	5801 21 5341 40 5541 93 SECTION CONFLICT

A

STUDENT COURSE CONFLICT LISTING 1/31/10

STUDENT NO. NAME CLASS CLASS NO SECT NO SECT

DAN305 DANIEL, DEBORAH C 09 5611 60
5501 50
5801 10
5521 40
5331 20
5671 30

DAN307 DANIEL, JANIE I 09 5611 60
3211 40
0211 40
5801 10
5521 50
5331 20
5671 30

DAN325 DANIELS, GEORGE D 09 5591 60
5301 30
5601 50
5571 10

DAN330 DANIELS, RON D 10 5591 60
0121 50
5691 30
5651 20
5501 40

5801 21
5301 40

SECTION COMPLETE

SECTION COMPLETE

STUDENT COURSE CONFLICT LISTING 7/31/78

STUDENT NO.	NAME	CLASS	-----C O U R S E S-----	
			ACCEPTABLE NO SECT	UNACCEPTABLE NO SECT

SECTION CONFLICT 292

CAPACITY CONFLICT 61

COURSE UNAVAILABLE 0

MUN300	Munday, Anthony	9	5733	10	
			5521	50	5991 40
			5611	60	

BUT350	Butler, James L	10	5651	60	
			0091	11	
			3991	10	
			5581	20	
			5381	30	
			5371	40	
			5691	50	

CAR700	Carter, William H	10	5691	10	
			5371	21	
			5381	30	
			5551	40	
			5501	50	
			5651	61	

MAL400	Malone, Catina L	9	5801	11	
			5291	20	
			5551	30	
			5881	40	
			5811	50	
			5521	60	

NOR700	North, Terry	10	5521	10	
			5371	20	
			0091	30	
			3241	30	
			5651	41	
			5691	51	
			101	10	

APPENDIX III

PROGRAM LISTINGS AND SAMPLE OUTPUTS

APPENDIX III

PART A. INPUT LISTINGS

COURSE	SECTION	SEATS	TYPE	SEMESTER	1TIME	1TIME1	1TIME2	1TIME3	1TIME4	NCR1	NCR2	NCR3
529	1	0	0	0	0	0	0	0	0	0	0	0
530	1	0	0	0	0	3	0	0	0	0	0	0
540	1	0	0	0	0	4	0	0	0	0	0	0
541	1	0	0	0	0	4	0	0	0	0	0	0
550	1	0	0	0	0	5	0	0	0	0	0	0
551	1	0	0	0	0	4	0	0	0	0	0	0
559	1	0	0	0	0	6	0	0	0	0	0	0
561	1	0	0	0	0	6	0	0	0	0	0	0
578	1	0	0	0	0	1	0	0	0	0	0	0
579	1	0	0	0	0	1	0	0	0	0	0	0
12	1	0	0	0	0	3	0	0	0	0	0	0
570	1	0	4	0	0	1	1	4	0	0	0	0
571	1	0	3	0	0	1	1	3	0	0	0	0
572	1	0	3	0	0	1	4	6	0	0	0	0
573	1	0	3	0	0	1	1	3	0	0	0	0
19	1	0	1	1	1	1	4	5	2	0	0	0
21	1	0	1	1	0	1	4	5	2	0	0	0
320	1	0	1	2	1	1	4	5	2	0	0	0
321	1	0	1	2	0	1	4	5	2	0	0	0
575	1	0	1	0	1	0	0	0	0	0	0	0
577	1	0	1	0	0	0	0	0	0	0	0	0
532	1	0	1	0	0	0	0	0	0	0	0	0
547	1	0	1	0	0	0	0	0	0	0	0	0
555	1	0	1	0	0	0	0	0	0	0	0	0
7	1	0	1	1	0	0	0	0	0	0	0	0
16	1	0	1	1	0	0	0	0	0	0	0	0
306	1	0	1	2	0	0	0	0	0	0	0	0
316	1	0	1	2	0	0	0	0	0	0	0	0
8	2	0	1	1	0	1	1	4	0	0	0	0
308	2	0	1	2	0	1	1	4	0	0	0	0
562	2	0	1	0	1	0	0	0	0	0	0	0
564	2	0	1	0	1	0	0	0	3	561	562	0
530	2	0	1	0	1	0	0	0	1	7	0	0
546	2	0	1	0	1	0	0	0	1	547	0	0
550	2	0	1	0	1	0	0	0	1	559	0	0
5	2	0	1	1	1	0	0	0	3	7	530	0
23	2	0	1	1	1	0	0	0	2	575	577	0
24	2	0	1	1	1	0	0	0	3	547	546	0
305	2	0	1	2	1	0	0	0	3	306	530	0
323	2	0	1	2	1	0	0	0	2	575	577	0
324	2	0	1	2	1	0	0	0	3	547	546	0
549	2	0	1	0	0	0	0	0	0	0	0	0
534	3	0	1	0	0	1	4	0	0	0	0	0
548	3	0	1	0	1	0	0	0	2	549	0	0
568	3	0	1	0	0	1	4	6	0	0	0	0
501	3	0	1	0	1	0	4	0	1	532	0	0
531	4	0	1	0	1	0	1	4	0	0	0	0
553	4	0	1	0	0	1	1	4	0	0	0	0
557	4	0	1	0	1	0	0	4	2	8	0	0
9	4	0	1	1	1	1	1	4	0	0	0	0
309	4	0	1	2	0	1	1	4	0	0	0	0
525	5	0	1	0	0	0	0	0	0	0	0	0
535	5	0	1	0	0	0	0	0	0	0	0	0
556	5	0	1	0	0	0	0	0	0	0	0	0
550	5	0	1	0	0	0	0	0	0	0	0	0
1	5	0	1	1	0	0	0	0	0	0	0	0

2	5	0	1	1	0	0	0	0	0	0	0	0
301	5	0	1	2	0	0	0	0	0	0	0	0
302	5	0	1	2	0	0	0	0	0	0	0	0
552	6	0	1	0	0	0	0	0	0	0	0	0
569	7	0	1	0	1	0	0	0	2	16	316	0
533	7	0	1	0	0	0	0	0	0	0	0	0
537	7	0	1	0	0	0	0	0	0	0	0	0
565	7	0	1	0	0	0	0	0	0	0	0	0
567	7	0	1	0	0	0	0	0	0	0	0	0
580	8	0	1	0	0	0	0	0	0	0	0	0
14	10	0	1	1	0	0	0	0	0	0	0	0
314	10	0	1	2	0	0	0	0	0	0	0	0

STUDENT ID	NAME	GRADE	SEX	REQUESTS	REQUEST LIST							
1001	ADAMS CINDY SUE	9	F	6	533	544	546	552	567	580		
1002	CHRISTIAN JEFF S	9	M	6	534	535	557	559	568	581		
1003	AKERNAN KALPH TALUE	9	M	4	548	553	559	573				
1004	ALEXANDER MICHAEL A	9	M	4	540	548	552	573				
1005	ALLEN TONY W	9	M	6	534	557	559	568	578	581		
1006	ANDERSON DIETRA JEA	9	F	7	24	321	531	534	552	568	581	
1007	ANDERSON ALBERT	9	M	6	531	534	542	557	560	581		
1008	ANNESLEY JAMES	9	M	6	530	534	548	557	568	581		
1009	ARGO GEORGIA LYNN	9	F	6	531	533	550	553	567	580		
1010	ARGO MARK WAYNE	9	M	6	533	548	553	559	567	580		
1011	ATNIP CHUCK RAY	9	M	6	534	535	557	559	568	581		
1012	AYRES KAREN RENEE	9	F	6	531	533	552	564	567	580		
1013	BOBICH KRISTEN ELIZ	9	F	5	21	321	551	552	573			
1014	BACON PEGGY SUE	9	F	7	8	324	534	544	553	568	581	
1015	BAKER LEAH ROBIN	9	F	3	530	546	570					
1016	BALLOU CATHERINE L	9	F	6	531	533	553	561	567	580		
1017	BALSBAUGH TOM JAY	9	M	6	531	533	552	560	567	580		
1018	BARELA LORI ANN	9	F	6	533	551	553	567	577	580		
1019	BARNARD GERALD RAY	9	M	6	531	534	548	557	568	581		
1020	BARNHILL RITA KAYE	9	F	6	525	531	533	553	567	580		
1021	BARR SHARI KAY	9	F	6	530	531	534	557	568	581		
1022	BELDEN CHERYL LYNN	9	F	6	531	533	552	567	575	580		
1023	BELL DEBRA LEE	9	F	4	531	538	552	573				
1024	BEST FRANK	9	M	7	23	399	534	557	568	578	581	
1025	BEVERLY ROUSCO SCOTT	9	M	4	531	552	559	573				
1026	BICKHAM GARY	9	M	6	534	535	548	557	568	581		
1027	BICKHAM GWENDOLYN	9	F	6	534	546	557	568	578	581		
1028	BLAKE MARICE GUDS	9	M	6	533	548	553	560	567	580		
1029	BOYD MARGARET ELEAN	9	F	4	530	552	573	578				
1030	BOYCE KARLA LEE	9	F	7	8	324	531	533	552	567	580	
1031	BOYD TRACEY ALISON	9	F	6	530	533	552	564	567	580		
1032	BRADLEY JOSEPH ANDERSON	9	M	4	548	550	550	573				
1033	BRATCHER TERRI LYNN	9	F	4	531	546	553	573				
1034	BRAWNER DANA LEA	9	F	5	19	320	551	552	573			
1035	BRESHEARS DEBRA ANN	9	F	6	531	534	546	557	568	581		
1036	BROCK TIMOTHY JAY	9	M	6	530	533	546	557	567	580		
1037	BROMS DONALD AUSTIN	9	M	6	531	533	552	559	567	580		
1038	BROWN TRACY DAWN	9	F	6	531	533	546	553	567	580		
1039	BRUCE TRACY LYNN	9	F	7	5	321	533	558	564	567	580	
1040	BRUEHL KIM ANN	9	F	7	24	308	531	533	552	567	580	
1041	BURKETT WILLIAM SID	9	M	7	5	306	534	551	557	568	581	
1042	BURLESON TEDDI ELIZ	9	F	6	533	550	552	564	567	580		
1043	BURNETT MICHAEL RAY	9	M	4	530	550	552	573				
1044	BURNETT TIMOTHY ALL	9	M	4	531	542	552	573				
1045	BURNSWORTH JUDI AN	9	F	7	24	305	533	546	557	567	580	
1046	BUTLER TINA MARIE	9	F	6	531	534	553	561	568	581		
1047	CADY JANANN L	9	F	6	531	533	550	553	567	580		
1048	CAMPBELL DOUGHLAS M	9	M	4	534	557	568	581				
1049	CANG DANNY LOZANO	9	M	6	531	534	557	559	568	581		
1050	CANO LISA LOZANO	9	F	6	531	534	557	562	568	581		
1051	CASEY JERRY GLENN	9	M	6	533	548	552	559	567	580		
1052	CASKEY JAMES WINDE	9	M	6	531	533	552	559	567	580		
1053	CAVENDISH LISA MARIE	9	F	7	21	321	533	552	561	567	580	

1054	CEDENO CHOURIO YAME	9	F	6	323	533	538	553	567	580	
1055	CHEVES KATHRYN MICH	9	F	6	531	533	557	567	575	580	
1056	CHOU SHOU YU SCOTT	9	M	4	8	531	558	573			
1057	CLARK RICHARD LEE	9	M	6	530	534	548	558	568	581	
1058	CLOUSE CYNTHIA LOU	9	F	6	530	531	534	557	568	581	
1059	COLLINS DAVID PATRI	9	M	4	544	558	559	573			
1060	CORLEY RAY	9	M	4	548	552	559	573			
1061	CORNELL PAUL	9	M	6	534	548	553	559	568	581	
1062	CLOUSE DAVID ALVIN	9	M	6	530	534	557	559	568	581	
1063	CUNNER DAVID BRYAN	9	M	4	531	552	559	573			
1064	CRAIG BALINDA JILL	9	F	4	538	551	552	573			
1065	CRAIG ROBYN BEA	9	F	6	530	533	552	561	567	580	
1066	CRANDELL JUANITA M	9	F	6	531	533	553	567	577	580	
1067	CROWN TAMARA DAWN	9	F	7	8	316	533	552	561	567	580
1068	CURREN EARL LEROY	9	M	4	533	557	567	580			
1069	DAIGNAULT ROBERT	9	M	6	530	533	550	553	567	580	
1070	DANIEL DEBORAH CAR	9	F	6	533	550	552	561	567	580	
1071	DANIEL TANIE IRENE	9	F	7	21	321	533	552	561	567	580
1072	DANIELS GEORGE DANE	9	M	6	530	534	557	559	568	581	
1073	DANNER KEVEN CHELT	9	M	6	23	533	544	553	567	580	
1074	DARNER KEITH	9	M	6	530	533	548	557	567	580	
1075	DAVUSON CARRIE LYN	9	F	7	8	324	531	533	552	567	580
1076	DAVIDSON JILL ANN	9	F	4	530	531	552	573			
1077	DAVIS DANA RUBIN	9	F	5	24	308	531	553	573		
1078	DAVIS FRANK JAMES	9	M	6	533	548	552	559	567	580	
1079	DAVIS LAWRENCE JR	9	M	7	16	308	533	552	559	567	580
1080	DELHEIMER JILL LYNN	9	F	6	531	533	546	557	567	580	
1081	DENSMORE JIMMY ROYC	9	M	6	531	533	557	567	575	580	
1082	DICKSON DORELLA JE	9	F	7	19	320	531	533	553	567	580
1083	DIRHAM BARFIE ELLEN	9	F	6	530	533	546	553	567	580	
1084	DOMINGUEZ DANIEL JA	9	M	6	533	548	552	559	567	580	
1085	DONATHAN MICHAEL EV	9	M	7	19	320	531	533	553	567	580
1086	DOWNY FLAVIL LEE	9	M	6	533	557	560	568	578	581	
1087	DRUM JONATHAN GILUE	9	M	6	530	533	550	553	567	580	
1088	DUNCAN BILL LEVI	9	M	4	544	553	559	573			
1089	DUNN KELLI JO	9	F	6	531	534	553	561	568	581	
1090	EARLS MARTIN LYNCH	9	M	8	5	21	321	323	533	553	567 580
1091	ELLIOT TAMMY KAY	9	F	7	8	324	534	557	564	568	581
1092	ELLIS BRADLEY MART	9	M	6	531	533	550	559	567	580	
1093	EMERY MICHAEL LEE	9	M	7	5	306	534	552	559	568	581
1094	EMERLIN PATRICA LYNN	9	F	7	5	324	534	557	564	568	581
1095	ESTES RICHARD LEE	9	M	6	533	540	550	553	567	580	
1096	EWING VANA LEE	9	F	7	5	306	533	546	557	567	580
1097	FARRIS TERRI GALE	9	F	6	533	553	564	567	577	580	
1098	FIELDS BARANDI EAYLE	9	F	6	530	531	533	557	567	580	
1099	FOX JEFFREY KENT	9	M	4	531	544	552	573			
1100	FRAUDIELE KIMBERLY	9	F	7	5	308	531	533	552	567	580
1101	FRANCIS KIRK LELAND	9	M	6	533	548	553	559	567	580	
1102	FRANKS SHAWN E	9	M	6	530	531	533	552	567	580	
1103	FRANKLIN KERI ALLIS	9	M	7	19	308	531	533	553	567	580
1104	FREEMAN JERRY THOM	9	M	7	23	308	533	553	559	567	580
1105	FREEZE HARI	9	F	6	533	553	561	567	575	580	
1106	FRENSLEY JANE ELIZA	9	F	4	531	546	552	573			
1107	FROST JIMMY RAY	9	M	7	5	323	533	552	560	567	580
1108	FRYAR REBECCA MARIE	9	F	6	534	557	564	568	575	581	
1109	GALLAHER THOMAS HER	9	M	4	531	552	573	575			
1110	GAMMON JANA LYN	9	F	4	546	558	573	575			
1111	GARDNER BRENDA KAY	9	F	3	546	564	570				
1112	GIADENS KIMBERLY RENAE	9	F	7	21	324	533	546	552	567	580
1113	GILL DAVID ALLEN	9	M	6	533	548	550	552	567	580	
1114	GILL JIMMIE JEAN	9	F	6	531	533	551	552	567	580	
1115	GOAD SONAI LAVARNE	9	F	6	533	540	553	561	567	580	
1116	GORDON ERIC LINN	9	M	6	534	548	557	559	568	581	
1117	GRADY TOMMY JAMES	9	M	3	548	570	585				
1118	GRAY JAMES WARREN	9	M	6	531	534	548	557	568	581	
1119	GREEN CARL ALLEN	9	M	4	533	552	567	580			

1120	GRISHAM TERRY VAN	9	M	6	533	548	557	561	567	580	
1121	GRIZZLE GARY DAVID	9	M	7	8	323	530	534	567	568	581
1122	GRIZZLE RICK JOE	9	M	4	531	548	552	573			
1123	GRICSE SUZANNE DICKI	9	F	5	19	320	538	552	573		
1124	GUINN LIA DAWN	9	F	7	19	320	533	557	564	567	580
1125	HAGGARD DERRICK WAY	9	M	7	16	308	534	553	559	568	581
1126	HALL BRENDA KAY	9	F	6	533	550	563	564	567	580	
1127	HALL JANICE MARIE	9	F	6	531	533	553	561	567	580	
1128	HALL LOUIS RAY	9	M	6	533	548	557	560	567	580	
1129	HALL ROBERT MATHEW	9	M	6	300	533	548	553	567	580	
1130	HAMES CHRISTIE HAYL	9	F	6	531	533	546	552	567	580	
1131	HAMES SHARON LYNN	9	F	6	531	534	546	557	568	581	
1132	HAMES STEVE LYNN	9	M	7	23	308	533	548	553	567	580
1133	HAMM BEVERLY F	9	F	4	533	552	567	580			
1134	HAMMON JERRY BRYCE	9	M	4	527	550	558	573			
1135	HARDCASTLE JACK	9	M	7	5	306	533	550	557	567	580
1136	HARDIN JILLY S	9	M	6	534	546	553	560	568	581	
1137	HART CHRISTIE E	9	F	6	531	533	550	553	567	580	
1138	HAWKINS BOBBY MICH	9	M	6	534	548	557	559	568	581	
1139	HAUSAM JAMES LESTER	9	M	6	323	531	534	553	568	581	
1140	HAWKINS NEAL RAY	9	M	6	531	533	548	553	567	580	
1141	HAYES JIMMY RAY	9	M	6	530	534	536	557	568	581	
1142	HEATON DONALD ALLE	9	M	6	24	530	533	552	567	580	
1143	HEBRON LUIS ALFONSO	9	M	4	551	552	559	573			
1144	HEDRICK JAMES ROBERT	9	M	6	534	548	557	560	568	581	
1145	HELMS FREDONNA LYNN	9	F	6	531	534	546	557	568	581	
1146	HENDERSON KELLY JAM	9	M	7	8	323	533	552	559	567	580
1147	HENDRICKS TERRY LYNNE	9	F	6	533	546	553	561	567	580	
1148	HENNING AMY LYNNE	9	F	5	5	324	531	553	573		
1149	HENSON GAVIN LEE	9	M	6	530	533	557	567	575	580	
1150	HENRY BUDDY LEE	9	M	6	534	535	548	567	568	581	
1151	HICKS CHERYL ANN	9	F	4	531	542	552	573			
1152	HICKSON DOUGLAS SCO	9	M	7	23	308	534	548	553	568	581
1153	HIGGINBOTHAM MAEIND	9	F	4	531	542	552	573			
1154	HOLLADAY DAVID RAND	9	M	4	531	552	573	575			
1155	HORNBACK JOHN PAUL	9	M	6	533	546	548	553	567	580	
1156	HORNBACK THOMAS DE	9	M	7	5	324	530	533	553	567	580
1157	HUDSON CHERYL KAY	9	F	4	533	552	567	580			
1158	HYDE MICHAEL LEE	9	M	6	534	548	557	568	578	581	
1159	JACKSON JAMES ALLEN	9	M	7	19	320	533	553	560	567	580
1160	JAMESON BOBBY J	9	M	7	5	324	534	557	559	568	581
1161	JAMESON DESIRA ANN	9	F	6	531	534	552	561	568	581	
1162	JENSEN GREG DEAN	9	M	6	24	533	548	552	567	580	
1163	JENNINGS KEN AUSTEN	9	M	6	533	548	552	559	567	580	
1164	JOHNSON CONNIE LY	9	F	4	544	558	562	573			
1165	JOLLS FRANK A	9	M	4	531	550	552	573			
1166	JOHNSON VALERIE ANN	9	F	7	24	308	531	533	553	567	580
1167	JUAREZ KATHY L	9	F	6	534	546	557	568	578	581	
1168	KADRI BANU NORA	9	M	4	530	552	559	573			
1169	KARNPHAU TRACY LYNN	9	F	7	8	324	533	538	552	569	580
1170	KIRBY KENNETH DON	9	M	6	534	548	557	560	567	580	
1171	KIRBY DONALD WAYNE	9	M	6	534	548	557	568	578	581	
1172	KIVLER KENNETH	9	M	4	8	540	557	573			
1173	KNIGHT NORA KRISTINE	9	F	6	531	534	546	557	568	581	
1174	KONDONAGSIS YGLANDA	9	F	5	19	324	538	558	573		
1175	KRAUSE BRIAN LEE	9	M	7	5	306	533	550	552	567	580
1176	KUHLMAN KATHY JOYC	9	F	4	546	550	552	573			
1177	LAMB CLINTON IKE	9	M	7	21	321	530	533	557	567	580
1178	LAMPRECHT SHERRY K	9	F	4	546	551	552	573			
1179	LANE BRIAN CODY	9	M	6	530	533	548	553	567	580	
1180	LANE KAREN KATHRYN	9	F	7	21	321	533	553	561	567	580
1181	LARKINS KELLY JOE	9	M	6	534	535	557	559	568	581	
1182	LARIMORE CHRISTINE	9	F	6	534	546	557	568	577	581	
1183	LAY JAMES TODD	9	M	6	530	531	533	557	567	580	
1184	LAZALIER KRISTEN AN	9	F	7	21	321	533	550	557	567	580
1185	LEONARD PEGGY LYNN	9	F	4	531	551	552	573			

1186	LEPPERT LYNOA LOUIS	9	F	4	550	558	561	573		
1187	LEWIS MARK ALAM	9	M	6	19	530	533	553	567	580
1188	LIBERTY STUART SPEN	9	M	6	530	533	553	559	567	580
1189	LOCKE CARLENE RHAE	9	F	5	24	316	538	552	573	
1190	LOVEDAY BILLIE LAV	9	F	4	531	552	573	575		
1191	LOUINS EARL DEAN	9	M	6	8	534	557	560	568	581
1192	LUNDY KIMIKO M	9	F	4	533	552	567	580		
1193	LYLES KARLA KAY	9	F	6	534	535	546	557	568	581
1194	MAHANES KIMBERLY ANN	9	F	6	533	546	552	564	567	580
1195	MARTIN KATHY DIANE	9	F	6	530	534	557	564	568	581
1196	MARTIN MARK	9	M	6	534	557	559	566	578	581
1197	MAST PAULA SUE	9	F	5	5	308	531	552	571	
1198	MASTERS DONNA MARIE	9	F	6	530	533	557	567	578	580
1199	MCCAIN LINDA LEA	9	F	6	533	546	552	561	567	580
1200	MCCOLLUM JANNINE SUE	9	F	6	530	533	540	552	567	580
1201	MCCOOL RICHARD IVY	9	M	6	533	552	559	567	575	580
1202	MCCRACKEN GLENDA KE	9	F	6	533	551	557	561	567	580
1203	MCCRAY PHYLLIS LYNN	9	F	7	21	321	533	557	567	575
1204	MCCURLEY GREGORY N	9	M	6	533	540	548	557	567	580
1205	MC GARITY JOE	9	M	6	530	533	550	553	567	580
1206	MCKENZIE DONNA SUE	9	F	5	8	324	550	552	573	
1207	MCLEHANEY TOMMY	9	M	6	533	540	548	553	567	580
1208	MCILLEN PAUL GENE	9	M	6	531	533	557	567	575	580
1209	MEEKS PAT DAVID	9	M	6	533	550	552	559	567	580
1210	MEINERS DONDY CARL	9	M	4	5	307	570	575		
1211	MELTON TROY	9	M	6	534	548	557	559	568	581
1212	MEYER WILLIAM PAUL	9	M	6	534	548	557	568	578	581
1213	MILLER ROSA DIANA	9	F	6	531	534	546	557	568	581
1214	MILLER TRACY SCOTT	9	M	6	533	548	557	560	567	580
1215	MITCHELL DRAIN LEE	9	M	6	534	548	553	560	568	581
1216	MITCHUM CINDY	9	F	6	534	550	557	564	568	581
1217	MITCHELL JANICE FAY	9	F	6	531	534	557	564	568	581
1218	MITCHELL LEE ANN	9	F	6	531	534	546	553	567	580
1219	MITTON DIANNE NARI	9	F	6	21	531	534	557	568	581
1220	MOORE IVAN RAY	9	M	6	533	548	552	559	567	580
1221	MONTGOMERY GLENN ALLEN	9	M	6	534	550	557	560	568	581
1222	MOORE JAMIE LOUISE	9	F	4	530	546	552	573		
1223	MOORE TOMMY GUY	9	M	6	533	548	552	559	567	580
1224	MURRIS HOLLY JEAN	9	F	7	8	305	531	533	553	567
1225	MORRISON CYNTHIA	9	F	6	531	533	552	567	578	580
1226	MORRISON BICKY DALE	9	M	7	5	323	534	548	557	568
1227	MYERS GREGORY GID	9	M	4	531	558	559	573		
1228	NEAL GLENN MITCELL	9	M	7	8	323	533	548	552	567
1229	NELSON ROBERT NATHAN	9	M	6	533	548	552	559	567	580
1230	NEWBY TERRY LYNN	9	M	6	531	533	553	564	567	580
1231	NEWTON JEFF ALLEN	9	M	6	534	551	553	559	568	581
1232	NCLAN SONYA MICHELL	9	F	7	21	321	531	533	553	567
1233	ODAM DAVID JEROME	9	M	6	530	533	538	557	567	580
1234	OWEN JEFFERY	9	M	6	530	533	540	553	567	580
1235	OWEN JOANNA L	9	F	7	21	321	533	546	553	567
1236	PARKS WADE ALLEN	9	M	6	530	533	538	552	567	580
1237	PARNELL CHARLENE ELLEN	9	F	6	531	533	553	561	567	580
1238	PEERY DEANNA LYNN	9	F	6	531	533	557	564	567	580
1239	PERSONS JEANNE ANNE	9	F	5	19	320	552	573	577	
1240	PESHEL ROBERT LEE	9	M	7	60	324	534	535	557	568
1241	PETERSON JAMES BLAK	9	M	6	534	548	557	559	568	581
1242	PETTYS DENPA ANN	9	F	6	533	552	562	567	577	580
1243	PLUARNIK MIKEL	9	M	6	5	533	548	553	567	580
1244	PHAM GIAO NGOC	9	M	5	7	323	552	559	573	
1245	PIERCE MAX RAY	9	M	5	8	316	558	559	573	
1246	PIGG DEBORAH LEANN	9	F	5	24	308	558	564	573	
1247	PGLASEK TIM P	9	M	6	530	533	540	552	567	580
1248	POWELL KENNY DAVID	9	M	6	533	544	548	557	567	580
1249	PRICE KIMBERLY SUE	9	F	6	534	546	553	561	568	581
1250	PRIMROSE JOHN GREGU	9	M	4	531	558	559	573		
1251	PROCTOR JOSEPH CLYD	9	M	6	533	552	560	567	578	580

1252	PUCKETT J LYNNETTE	9	F	6	531	533	553	564	567	580	
1253	QUINTON STANLEY GEN	9	M	4	531	550	552	573			
1254	RAINER JEFFREY SCOT	9	M	3	23	308	558	559	573		
1255	RANEY HOWARD JAMES	9	M	7	5	323	534	551	557	568	581
1256	RANEY JULIE ESTHER	9	F	5	19	320	531	552	573		
1257	REA JACK TRACY	9	M	8	5	23	306	324	534	557	568 581
1258	REDCORN YANCEY GAR	9	M	4	550	552	561	573			
1259	REDMON LISA ANN	9	F	6	530	533	557	567	577	580	
1260	REICH RONNIE DEAN	9	M	6	530	533	553	559	567	580	
1261	REICH RONNIE DEAN	9	M	6	530	533	553	559	567	580	
1262	REID LAURA ALLENE	9	F	6	533	542	553	567	575	580	
1263	REID MICHAEL FORRES	9	M	4	548	551	550	573			
1264	RHEA LISA ANN	9	F	6	533	552	564	567	577	580	
1265	RILEY TERESA DALE	9	F	4	531	538	553	573			
1266	RINGWALD JOHN ALLEN	9	M	6	531	534	538	557	568	581	
1267	ROBERTS GURDON CLIFFORD	9	M	6	534	535	557	559	568	581	
1268	ROBERSON BETH	9	F	6	531	533	553	561	567	580	
1269	ROBERTS PETER NICH	9	M	7	19	320	533	538	553	567	580
1270	ROCKETT TAMMY K	9	F	6	531	534	551	557	568	581	
1271	ROGERS JASON LEWIS	9	M	6	533	540	557	560	567	580	
1272	ROLLER GERRI LYNN	9	F	6	24	533	550	552	567	580	
1273	RUSSELL GURDON PAUL	9	M	6	531	533	548	557	567	580	
1274	RUSSELL MITCH RAY	9	M	4	5	306	570	585			
1275	SANDERS CERRY JEAN	9	F	6	531	533	553	561	567	580	
1276	SANDERS JOHN JOSEPH	9	M	6	530	533	540	553	567	580	
1277	SANDERS LORANNA MA	9	F	7	8	305	531	533	553	567	580
1278	SCHLIENZ DARLEENA S	9	F	6	534	546	557	564	568	580	
1279	SCHNERINGER S RENEE	9	F	7	5	324	531	533	552	567	580
1280	SETTLE DAVID HARRY	9	M	6	533	544	553	559	567	580	
1281	SETTLE FRED WESLEY	9	M	6	533	544	553	559	567	580	
1282	SHAFFER SHERRY KATH	9	F	7	5	324	533	557	561	567	580
1283	SHAVER DUNNA CAROL	9	F	6	533	553	564	567	577	580	
1284	SHAW DONALD RAY	9	M	7	8	316	533	548	553	567	580
1285	SHAW ROBERT WAYNE	9	M	6	533	540	557	560	567	580	
1286	SHEARER KATHY ANN	9	F	5	5	324	550	552	573		
1287	SHELL JUDY ANN	9	F	6	531	533	553	567	577	580	
1288	SHERIDAN LYNDA KAYE	9	F	7	5	324	533	546	553	567	580
1289	SHINERT THERESA LOU	9	F	6	533	553	561	567	577	580	
1290	SIMONS HAROLD ALLEN	9	M	7	5	306	534	548	553	568	581
1291	SIMS SHANNON DEE	9	F	4	533	552	567	580			
1292	SMITH FLOYD ALLEN	9	M	7	23	308	533	548	552	567	580
1293	SMITH KIMBERLY ANN	9	F	6	534	557	562	568	577	581	
1294	SMITH STEVEN CHARLE	9	M	6	533	548	553	559	567	580	
1295	SMITHSON FRANKLIN DELANA	9	M	6	530	533	542	552	567	580	
1296	SHELL ALATHEA JOY	9	F	5	544	552	573	577	580		
*** INFEASABLE REQUEST LIST ***											
1297	SPINK PAUL SANKER	9	M	4	533	552	567	580			
1298	SPOR LISA ANN	9	F	4	531	546	552	573			
1299	STACY BET WAYNE	9	M	6	533	548	552	559	567	580	
1300	STAUBER MORGEN GLAR CLAR	9	F	6	531	533	546	557	567	580	
1301	STEPHENS CAROL LOUI	9	F	6	530	534	546	557	568	581	
1302	STJOHN MATTHEW CARL	9	M	6	533	548	553	559	567	580	
1303	STRATTON BRADLEY DO	9	M	7	19	308	533	557	559	567	580
1304	STRAUGHN JOY LYN	9	F	5	24	308	531	552	573		
1305	TALBOT DAVID	9	M	6	534	548	557	559	568	581	
1306	TATE STEVEN BLAIR	9	M	7	5	323	534	553	560	568	581
1307	TAYLOR TODD NOLAN	9	M	6	533	548	553	559	567	580	
1308	TILLER SHARLA ELAIN	9	F	4	531	552	561	573			
1309	TOWNLEY ELLEN ELIZA	9	F	6	531	533	551	552	567	580	
1310	TUCCI LOUIS JOHN	9	M	6	533	544	552	559	567	580	
1311	TULLIUS BETTY JANE	9	F	4	530	558	561	573			
1312	TURNER JEFF MICHAEL	9	M	7	5	324	534	546	557	568	581
1313	UNDERWOOD TERRY LEE	9	M	7	5	324	534	548	553	568	581
1314	UPTEGRAFT NATHAN KR	9	M	4	531	552	559	571			
1315	URDARL RICHARD D	9	M	5	21	321	552	559	573		
1316	VALOUCH JANIE LOUIS	9	F	6	24	308	531	533	553	580	

1317	VANDERHELM MICHAEL	9	M	6	533	538	553	560	567	580	
1318	VANNASTRAM DARLA DE	9	F	7	19	321	533	552	567	577	580
1319	VESELY STEVEN MURE	9	M	5	24	308	553	559	573		
1320	VIS DANIEL RAYMOND	9	M	4	24	531	558	573			
1321	WAGGONER FRANK GREG	9	M	6	534	535	548	553	568	581	
1322	WAGNER MARY MARGARET	9	F	4	531	558	561	573			
1323	WALKER CHARLES EUGENE	9	M	6	530	533	553	560	567	580	
1324	WALKER KENT LEE	9	M	7	5	306	533	552	559	567	580
1325	WALTNER KELLY ANN	9	F	4	531	558	561	573			
1326	WARDIUS GEORGE SCOT	9	M	4	548	558	559	573			
1327	WARREN PAUL EUGENE	9	M	6	533	557	560	567	578	580	
1328	WEAVER RANDAL SCOTT	9	M	5	24	323	540	558	573		
1329	WEBB WARREN DOUGLAS	9	M	6	533	548	552	560	567	580	
1330	WEBER KARL BERNARD	9	M	7	5	324	533	557	560	567	580
1331	WEED TINA MARIE	9	F	6	530	534	546	557	568	581	
1332	WEIRICH AMY LOUISE	9	F	6	531	533	558	564	567	580	
1333	WESNER JACOB C	9	M	6	16	19	308	320	558	573	
1334	WHEELER BRAD LEE	9	M	7	23	308	530	534	557	568	581
1335	WHEELER KEVIN SCOT	9	M	7	5	306	533	552	559	567	580
1336	WHITE SHEILA KAY	9	F	6	531	534	538	557	568	581	
1337	WHITE MARE S EDWARD	9	M	6	534	551	557	567	575	580	
1338	WIGGINS RONNIE G	9	M	6	534	535	548	557	568	581	
1339	WILKINSON E CHARLES	9	M	6	530	533	553	559	567	580	
1340	WILLCOX TIOMMIE JOE	9	M	6	5	534	544	557	568	581	
1341	WILLIAMS MICHELLE D	9	F	6	531	533	553	561	567	580	
1342	WILLIAMS RUBYNN LEE	9	F	7	5	308	531	533	553	567	580
1342	WILLIAMS ROOGN LEE	9	F	7	5	308	531	533	553	567	580
1343	WILLIAMS SCOTTIE RA	9	M	7	8	324	534	548	553	568	581
1344	WITCHER JAMES CHRISTOPHER	9	M	7	8	323	533	548	552	567	580
1345	WOLF KIM RENEE	9	F	7	5	308	531	533	553	567	580
1346	WOLVERTON LISA ANN	9	F	6	534	557	564	568	575	581	
1347	YATES LARRY TODD	9	M	6	533	544	557	559	567	580	
1348	YOUNG JOE ARVINE	9	M	4	534	557	568	580			
1349	YOUNG MARK GREGORY	9	M	6	530	531	533	552	567	580	
1350	YOUNG VAN	9	M	6	534	548	557	560	568	581	
1351	ZANGLA TONI DENISE	9	F	6	533	553	564	567	577	580	
1352	ZASTE MARLIN	9	M	7	5	306	534	557	560	568	581
1353	ZIBULSKY SUZANNE	9	F	7	21	321	531	533	552	567	580
1354	ZIMMER RICKY DEAN	9	M	3	548	570	585				
1355	MCDONNELL DOLAN JAM	9	M	6	531	534	538	553	568	581	
2001	AKIN PATRICIA SUE	10	F	7	12	309	537	538	558	561	569
2002	ALEXANDER O LORNA GAI	10	F	7	8	309	531	537	558	565	569
2003	ALINGER DANIEL SCOT	10	M	7	23	309	533	549	556	565	569
2004	ALLISON JALYN DENIS	10	F	4	539	551	558	571			
2005	AMCS DENISE ANN	10	F	2	556	571					
2006	ANDERSON ROBERT MIC	10	M	7	12	309	537	558	559	565	569
2007	ARGO MICKEY E	10	M	7	8	323	537	558	559	565	569
2008	ARMSTRONG JAMI JUI	10	F	4	309	539	558	571			
2009	AUGHTRY JONNIE LYNN	10	F	7	12	309	537	558	561	565	569
2010	AVILES RICHARD	10	M	3	8	305	572				
2011	BAABE JIMMY	10	M	6	530	537	549	556	566	569	
2012	BAIRD THOMAS EDWARD	10	M	6	16	557	560	565	569	575	
2013	BARCOCK SHARI LYNN	10	F	6	9	531	537	552	565	569	
2014	BACHARCH CASSIE LAA	10	F	5	21	309	539	558	571		
2015	BACON RANDALL DEWAY	10	M	7	24	309	537	558	559	565	569
2016	BAIRD ROBERT MARC	10	M	7	9	308	537	556	560	566	569
2017	BAKER DAVID SHAWN	10	M	7	24	309	537	552	559	565	569
2018	BANKS RENITA FAYE	10	F	5	16	309	531	567	572		
2019	BARKSDALE GARY LYNN	10	M	7	9	323	537	544	552	565	569
2020	BARNARD MARVIN WAYNE	10	M	5	9	317	548	566	572		
2021	BARNETT ROBERT DALE	10	M	4	309	567	572	578			
2022	BATES SCOTTY NOLAN	10	M	5	5	309	566	572	578		
2023	BELKNAP HAL R	10	M	6	537	538	558	560	565	569	
2024	BELYEA DAVID J	10	M	5	9	308	549	558	571		
2025	BENTON CHRIS STEPH	10	M	7	9	323	537	550	558	565	569
2026	BEVILLS RUSSELL WAY	10	M	5	9	316	560	566	572		

2027	BLACK DECKY LYNDELL	10	F	5	16	309	561	566	572		
2028	BLAYLOCK DONNA LUU	10	F	7	23	309	537	556	565	569	575
2029	BLAYLOCK KARLA K	10	F	4	309	550	558	571			
2030	BELVINS ALICIA MAE	10	F	4	541	558	571	575			
2031	BOGGS KELLIE JO	10	F	5	9	305	539	550	571		
2032	BOND CORINA L	10	F	5	5	309	531	564	572		
2033	BUONE BRIAN KEITH	10	M	7	23	309	537	558	560	565	569
2034	BOWEN RICHARD	10	M	5	9	305	308	323	572		
2035	BOX MICHAEL ALBERT	10	M	7	23	309	537	549	552	565	569
2036	BRESHEARS DONNA SUE	10	F	5	9	316	531	566	572		
2037	BRITT RENA KAY	10	F	7	21	309	532	537	556	565	569
2038	BROWN JEANETTE CARO	10	F	7	9	308	537	556	565	569	575
2039	BROWN MELUDY ANN	10	F	7	9	305	537	556	565	569	575
2040	BROENLEE MARQUETTA	10	F	7	24	309	537	556	565	569	577
2041	BRAYANT DANA CARLO	10	F	7	19	309	537	550	558	565	569
2042	BRAYANT DARREL JAMES	10	M	6	534	535	557	560	568	581	
2043	BURGETT WILMA JEAN	10	F	5	9	305	544	566	572		
2044	BUSSEY TERRY LYNN	10	M	6	9	537	549	556	565	569	
2045	BUTLER DANNY JO	10	M	7	9	305	537	558	559	565	569
2046	BUTLER GLENNA JEAN	10	F	5	23	309	532	558	571		
2047	BYRD BARBARA MARIE	10	F	6	531	537	556	565	569	579	
2048	CAMPHELL KIMBERLY R	10	F	7	9	308	537	558	565	569	577
2049	CAPE CHESTER DEWAYN	10	M	7	8	309	537	552	559	565	569
2050	CARPENETR DAVID M	10	M	5	21	309	566	572	580		
2051	CARSON DAN WALLACE	10	M	6	531	537	556	565	569	575	
2052	CARTER JOHN CHRIS	10	M	5	17	309	566	572	579		
2053	CARY DAVID A	10	M	5	9	540	550	558	571		
*** INFEASABLE REQUEST LIST ***											
2054	CASH CARLENE'D	10	F	5	5	324	562	566	572		
2055	CAVES LISA MARIE	10	F	6	531	537	558	565	569	575	
2056	CHRISTIAN IAN MARK	10	M	4	537	550	555	559			
2057	CLAEYS CAROL GAIL	10	F	7	7	323	537	556	561	566	569
2058	CLARK LUMMIE	10	F	7	9	324	537	556	561	565	569
2059	CLARK TRACY LEIGH	10	M	7	7	309	537	539	558	565	569
2060	CLEAR KEAL WAYNE	10	M	6	17	309	556	559	567	569	
2061	CLINKENBEARD LISA J	10	F	5	531	546	558	571	575		
*** INFEASABLE REQUEST LIST ***											
2062	CLASON LISA ANNETTE	10	F	7	8	309	537	550	552	567	569
2063	COLWELL STEVE ED	10	M	7	24	309	537	558	559	565	569
2064	COMPTON STEVEN SC	10	M	4	309	549	566	572			
2065	CONNOR MICHAEL STEVE	10	M	7	8	309	537	552	559	565	569
2066	COSSEY TIM J	10	M	5	5	317	560	565	572		
2067	COUCH STEVEN DALE	10	M	7	9	305	537	558	559	565	569
2068	COX ANNETTE	10	F	5	309	324	558	564	571		
*** INFEASABLE REQUEST LIST ***											
2069	CRAIG BRENDA ANN	10	F	5	8	309	550	556	571		
2070	CRAIG TWESSIE ROSS	10	F	4	539	551	558	571			
2071	CRANDELL ROBERT D	10	M	4	559	566	572	578			
2072	CROWN MARC HOWARD	10	M	7	9	308	537	558	559	565	569
2073	CURRY SHARON KAYE	10	F	3	567	572	580				
2074	DANG NGUYET MINH	10	F	7	7	309	537	539	552	565	569
2075	DANIEL BARBARA ELAINE	10	F	4	531	566	572	575			
2076	DANIEL RON DALE	10	M	7	19	309	537	558	559	565	569
2077	O DAVIS BRENDA KAE	10	F	8	16	24	308	309	537	552	565 569
2078	DAVIS DEE ANN	10	F	4	9	325	570	575			
2079	DAVIS MELUDY ANN	10	F	6	9	531	537	558	565	569	
2080	DAVIS REGINA	10	F	5	7	537	553	565	569		
2081	DAWSON DEBRA LYNN	10	F	5	9	316	531	558	571		
2082	DETTAVEN ROSANNA MAR	10	F	7	9	305	537	552	561	565	569
2083	DEHN TERRI LYNN	10	F	6	9	537	550	556	565	569	
2084	DELAUGHTER JOHN	10	M	4	530	543	558	571			
2085	DELHEIMER THOMAS STEVEN	10	M	6	9	537	549	566	565	569	
2086	DENSMORE TERESA JO	10	F	6	531	537	552	564	565	569	
2087	DELVEES JOHN	10	M	5	8	537	548	565	569		
2088	DIERS GANE G	10	M	4	9	550	558	571			
2089	DONALDSON WES SCOTT	10	M	6	9	537	556	565	569	579	

2090	DOYLE TRACY EUGENE	10	M	6	309	537	556	565	569	579		
2091	DRAUGHN DANNY RAY	10	M	4	548	559	566	572				
2092	DRISKELL CINDY LEN	10	F	7	12	309	537	550	558	569	577	
2093	DISWELL KARLA ANN	10	F	8	23	24	309	316	537	556	567	569
2094	DUMP TAMMY GAIL	10	F	7	12	309	537	550	558	565	569	
2095	DURHAM KELLY ELIZA	10	F	5	12	309	558	571	675			
2096	DURKEE GARY RHFA	10	M	6	531	537	541	552	565	569		
2097	DUVALL MICHAEL LEE	10	M	6	309	545	552	559	565	569		
2098	EADES DORIAN SUE	10	F	4	9	532	566	571				
2099	EDINGTON GARD DON	10	M	5	21	309	558	559	571			
2100	ELLIOT LORI H	10	F	4	9	325	547	570				
2101	FIELDS JEFFREY ALAN	10	M	6	309	537	549	558	565	569		
2102	FISHBECK LESLIE MA	10	F	6	8	309	532	551	558	571		
*** INFEASABLE REQUEST LIST ***												
2103	FISHER LAWRENCE MIC	10	M	6	8	309	550	558	559	571		
*** INFEASABLE REQUEST LIST ***												
2104	FLOYD BRAIN	10	M	4	9	305	570	585				
2105	FOWLER TOMMY RAY	10	M	7	21	309	537	556	565	569	575	
2106	FRAYAR WILLIAM LEE	10	M	4	9	305	570	585				
2107	GAGE SHELLEY DIANA	10	F	7	9	308	537	564	565	569	575	
2108	GASS LOUISE ALMA	10	F	6	532	537	556	565	569	578		
2109	GATES RONALD JAY	10	M	6	537	541	551	556	565	569		
2110	GENTRY COLIN LYNN	10	M	7	23	309	537	550	558	565	569	
2111	GHASEMI NAFISEH	10	F	6	9	537	538	556	566	569		
2112	GIBBONS MELISSA DEA	10	F	5	9	308	531	558	571			
2113	GOAD MICHAEL	10	M	2	566	572						
2114	GRADY HAROLD WAYNE	10	M	4	9	585	548	570				
*** INFEASABLE REQUEST LIST ***												
2115	GRAF KARIE MARIE	10	F	7	9	308	537	558	565	569	575	
2116	GREEN KENNY JOE	10	M	6	537	549	552	565	569	579		
2117	GRAFFIN ANNA CECIL	10	F	7	8	309	537	552	565	569	575	
2118	GROVE GLORIA	10	F	6	535	537	556	566	569	577		
2119	HAAS LORI MICHELE	10	F	5	12	309	550	558	571			
2120	HAMBY ROSS RANDALL	10	M	7	7	309	537	548	556	565	569	
2121	HAMILTON DICK H	10	M	7	5	309	537	544	549	558	569	
2122	HANGER STEVEN HUGO	10	M	7	7	309	537	545	558	566	569	
2123	HANSMEYER TODD EUGE	10	M	6	537	552	565	569	575	579		
2124	HARDEN KAREN KYM	10	F	4	309	531	558	571				
2125	HARDMAN MELEIA LYNN	10	F	5	24	309	525	558	571			
2126	HARKEY BEVERLY LOUI	10	F	6	532	537	556	564	565	569		
2127	HARRI PATRICK	10	M	7	7	309	533	556	560	565	569	
2128	HARRIS HAREN	10	F	7	24	317	537	556	564	565	569	
2129	HARRIS SHERRY JANE	10	F	4	550	558	561	571				
2130	HARRIS JERRY HOWARD	10	M	4	531	551	558	571				
2131	HART TERRI ALANA	10	F	5	9	308	550	556	571			
2132	HARTMAN RAYMOND EAR	10	M	6	537	549	552	565	569	575		
2133	HAUSAM BETTY JEAN	10	F	7	24	309	531	537	556	565	569	
2134	HEARD STEVEN ISAAAC	10	M	7	60	309	537	552	565	569	575	
2135	HEITZ MARCIE LYNN	10	F	7	9	364	537	546	556	565	569	
2136	HELMIT KATHI LYNN	10	F	5	23	309	551	566	572			
2137	HENDRICKS LINDA MAR	10	F	7	5	323	531	537	556	565	569	
2138	HENNINGS BILL ALEX	10	M	7	23	309	537	549	556	565	569	
2139	HENSLEE BILLYE G	10	F	4	537	556	565	569				
2140	HEWITT LISA ANN	10	F	5	24	309	532	558	571			
2141	HICKS JIMMY	10	M	5	9	323	549	566	572			
2142	HODSON DAVID MICHA	10	M	4	539	558	559	571				
2143	HOLMAN STEVEN DUNAL	10	M	4	525	558	559	571				
2144	HOPKINS KAREN DENIS	10	F	5	8	309	558	571	578			
2145	HORN DECKY DEEANNE	10	F	7	9	324	537	552	565	569	575	
2146	HORNBACK LISA ANN	10	F	4	550	556	561	571				
2147	HOSKINSON	10	F	5	24	399	562	566	572			
2148	HOWLAND KATHLEEN AN	10	F	7	24	309	537	558	561	565	569	
2149	HUGHES DANIEL	10	M	7	9	308	535	549	556	568	569	
2150	HUGHES DON RAY	10	M	5	12	309	558	559	571			
2151	HUGHES JOHN CHARLES	10	M	6	8	16	309	324	566	572		
2152	HUNT CARRIE BETH	10	F	6	531	537	538	558	569	577		

2153	HUNTER BEN	10	M	6	537	548	556	565	569	579	
2154	HUNTER BOBBY LYNN	10	M	4	537	556	566	569			
2155	INGELS CAROL SUSANN	10	F	6	537	543	550	558	565	569	
2156	IRWIN ROBERT LOUIS	10	M	6	19	309	541	550	558	571	
*** INFEASABLE REQUEST LIST ***											
2157	JACOBS JUDY LYNN	10	F	7	16	17	21	308	531	552	564
2158	JAMAR ROSEMARY CARO	10	F	7	16	309	537	551	552	565	569
2159	JENNINGS SHEILA L	10	F	7	9	305	531	537	558	565	569
2160	JESSOP SAMUEL ZANE	10	M	4	549	560	566	572			
2161	JEWELL ALUCE V	10	F	8	5	8	309	323	537	556	565 569
2162	JOHNSON PATTY JANN	10	F	7	9	306	531	537	556	565	569
2163	JOHNSON SPENCER COLE	10	M	6	537	542	558	559	565	569	
2164	JONES KENT MARUICE	10	M	7	60	309	537	556	565	569	578
2165	JOHNSON MARK	10	M	7	16	308	531	537	558	565	569
2166	JONES JERRI LYNN	10	F	6	530	537	539	556	565	569	
2167	KEITH LELIA SHELLY	10	F	7	9	324	537	556	564	565	569
2168	KELLEY LADONNA SUE	10	F	7	9	338	530	537	556	565	569
2169	KENNEDY GEWEN ELAINE	10	F	7	21	308	537	551	556	565	569
2170	KENNEDY LEA ANN	10	F	4	532	538	552	571			
2171	KERRY KARI DAWN	10	F	7	25	309	537	552	565	569	575
2172	KERR KIMHERLY ANNE	10	F	6	531	537	552	561	565	569	
2173	KERSEY TONY DALE	10	M	6	531	537	558	565	569	579	
2174	KESLER CYNTHIA ANN	10	F	5	9	324	532	558	571		
2175	KLIMA CRAIG W	10	M	7	5	309	537	558	559	565	569
2176	KLIMA CRAIG W	10	M	7	5	309	537	558	559	565	569
2177	KYURIGER SHARON LOUI	10	F	7	9	537	543	550	558	565	569
*** INFEASABLE REQUEST LIST ***											
2178	LACK PAUL EDWARD	10	M	7	9	308	537	544	558	565	569
2179	LANDRY SHERRY LYNN	10	F	7	24	309	537	556	564	565	569
2180	LANE COLLEEN ANN	10	F	6	537	545	552	565	567	569	
2181	LANGE KIMBERLY ELA	10	F	5	8	309	538	558	571		
2182	LASATER TERRY D	10	M	7	8	309	537	556	565	569	575
2183	LAWRENCE LISA GAIL	10	F	6	530	537	556	565	569	575	
2184	LEMONS TONY COLE	10	M	7	9	317	533	552	560	565	569
2185	LESSLY MARK EDWARDS	10	M	7	7	323	537	557	559	565	569
2186	LEWIS DANE D	10	M	4	549	556	560	572			
2187	LOEWEN JANICE MARIE	10	F	4	531	544	558	571			
2188	LOVINS DONNA JO	10	F	7	9	399	537	552	561	565	569
2189	LOYD VICKI ANNETTE	10	F	6	531	537	552	565	569	575	
2190	LUCAS MARY ELIZABETH	10	F	4	9	305	547	570			
2191	LUMPKINS JONATHAN R	10	M	4	531	549	555	571			
2192	LUNDY DAPHNE LOU	10	F	5	9	308	550	558	571		
2193	LYNN DONALD B	10	M	5	9	323	558	559	571		
2194	LYONS JILL M	10	F	8	8	24	305	309	537	556	565 569
2195	LYONS JILL M	10	F	5	537	556	561	565	569		
2196	MANGUS DEBBIE LYNN	10	F	7	9	308	537	556	565	569	579
2197	MANNING JAY DEE	10	M	4	549	560	566	572			
2198	MASTERS KAEN ANN	10	F	5	12	309	558	561	571		
2199	MAULDIN PATRICIA M	10	F	5	9	308	531	566	572		
2200	MAYBERRY LANDA DAW	10	F	4	538	555	561	571			
2201	MAYES CYNTHIA ANNE	10	F	6	525	537	556	564	565	569	
2202	MAYES KELLY MICHELLE	10	F	6	530	532	537	552	565	569	
2203	MCALISTER ROBIN CLA	10	F	5	8	309	558	559	571		
2204	MCCARREL TAMMY LYN	10	F	6	537	557	564	565	569	575	
2205	MCCAUGHTRY DARR LYN	10	F	6	531	537	552	564	565	569	
2206	MCCLELLAN MITCH GAR	10	M	5	9	360	549	566	572		
2207	MCDONALD BRYAN WAYN	10	M	7	9	308	537	550	559	565	569
2208	MCGRATH JENNIE KAY	10	F	4	530	564	565	572			
2209	MCKENZIE LINDA KAY	10	F	7	5	309	531	537	556	565	569
2210	MENUTT JIMMY RAY	10	M	7	9	308	537	550	558	565	569
2211	MEADOWS SHAWNA GAY	10	F	5	9	324	532	558	571		
2212	MEISSNER PATRICIA L	10	F	7	9	308	532	537	558	565	569
2213	MESSERLE SUSANNE B	10	F	4	531	544	558	571			
2214	MEYER DENISE KAE	10	F	8	7	9	305	325	531	537	556 569
2215	MEYER DONNA RAE	10	F	7	9	325	530	532	537	552	569
2216	MEYER JODIE LEANN	10	F	5	9	324	531	537	572		

2217	MILLER DAVID RAY	10	M	7	24	309	537	558	559	565	569
2218	MILLER JAMES M	10	M	6	531	537	558	559	565	569	
2219	MITCHELL RICK ALLEN	10	M	7	24	309	537	556	560	565	569
2220	MITTEN LORI ANNE	10	F	7	12	309	532	537	558	565	569
2221	MIZE VALERIE JEAN	10	F	4	547	566	572	577			
2222	MONTAYNE RODIN N	10	F	7	9	325	537	556	569	575	577
2223	MONTGOMERY DA K	10	M	7	9	323	530	537	556	565	569
2224	MONTGOMERY SHELIA L	10	F	6	537	543	558	565	569	577	
2225	MUORE KEVIN WADE	10	M	5	60	309	549	558	571		
2226	MORREN THOMAS R	10	M	7	7	309	537	552	559	565	569
2227	MUSES DARRYL	10	M	7	16	309	537	556	559	565	569
2228	MURPHY CYNTHIA ANN	10	F	5	9	361	545	558	571		
2229	MUSGRAVE SHERRI LYN	10	F	7	62	309	537	550	558	565	569
2230	MYERS NANCY DARLENE	10	F	6	532	537	550	558	565	569	
2231	NEFF MARY ANN	10	F	7	9	323	537	552	565	569	575
2232	NELMS BRENT RUSSELL	10	M	5	8	309	559	572	578		
2233	NEWBERG GREG	10	M	2	566	572					
2234	NEWTON KARI LISA	10	F	4	558	564	571	577			
2235	NICKERSON DOUGLAS A	10	M	7	59	309	537	550	552	565	569
2236	NOTTENSMEYER EVELYN	10	F	5	8	309	531	566	572		
2237	OVERMILLER MARK GUN	10	M	5	8	309	558	559	571		
2238	ODOM ANGELA G	10	F	7	9	316	537	547	552	565	569
2239	OLIPHANT MELISSA JA	10	F	7	21	309	537	558	564	565	569
2240	ORTH DANIEL LEO	10	M	7	9	324	537	552	559	565	569
2241	PATTERSON LARRY E	10	M	7	9	308	537	552	560	565	569
2242	PAYNE PAMELA J	10	F	5	9	308	532	567	572		
2243	PEACE ROBERT WAYNE	10	M	7	9	308	537	549	552	565	569
2244	PELLEY KATHERYNE R	10	F	6	9	308	531	551	558	571	
*** INFEASABLE REQUEST LIST ***											
2245	PHILLIPAS SUSAN LYNN	10	F	5	19	309	566	572	578		
2246	PLASTER SHELLY LYN	10	F	7	9	308	537	556	565	569	578
2247	POLK DAVID DWAIN	10	F	7	8	323	537	552	559	565	569
2248	POTTS JERRY	10	M	2	566	572					
2249	POTTS LORI L	10	F	7	64	309	532	537	552	565	569
2250	POWELL JAMIE W	10	M	6	537	549	556	565	569	579	
2251	PRIDEMORE CAROLYN G	10	F	7	9	308	537	558	559	565	569
2252	PUGH GREGORY WILLIAM	10	M	7	24	309	537	544	558	565	569
2253	RAEL ORLANDO FRANCIS	10	M	7	6	305	531	537	558	565	569
2254	RAHAY RUSSELL C	10	M	5	8	399	560	572	579		
2255	RAMSEY TODD M	10	M	6	537	544	550	558	565	569	
2256	RANEY JOHN DAVID	10	M	7	7	309	537	551	556	565	569
2257	REED SHELLEY LYNN	10	F	6	537	556	564	565	569	577	
2258	REID NATFAN LINN	10	M	7	19	309	537	543	550	552	569
2259	REYNOLDS DALE C	10	M	6	531	537	549	558	565	569	
2260	REYNOLDS ERIC EUGEN	10	M	7	16	309	530	537	558	565	569
2261	REYNOLDS RALPH THOM	10	M	7	24	309	537	556	559	565	569
2262	RICHARDSON JANEY	10	F	5	12	309	558	564	571		
2263	RIDGE WAY BURNEY GENE	10	M	5	8	309	559	572	578		
2264	RIEGER ROBERT HENRY	10	M	6	9	323	531	548	555	571	
*** INFEASABLE REQUEST LIST ***											
2265	RITCHEY KAREN CLARK	10	F	7	64	309	531	537	552	565	569
2266	ROBBINS BRYAN KEITH	10	M	7	8	324	537	552	559	565	569
2267	ROPER JOE JERAL	10	M	7	23	309	537	549	556	565	569
2268	ROSS JILL ELAINE	10	F	6	9	300	537	552	565	569	
2269	ROSS SHERYL DENISE	10	F	5	8	309	550	555	571		
2270	RUSSELL GREG THOMAS	10	M	6	9	306	539	556	559	571	
*** INFEASABLE REQUEST LIST ***											
2271	RUSSELL MARK BROOKS	10	M	7	7	309	537	550	558	565	569
2272	RUSSELL ROBERT LEON	10	M	6	537	549	556	559	565	569	
2273	SANCHEZ DORA DELISA	10	F	6	537	539	556	561	565	569	
2274	SANDERS RALPH MARK	10	M	5	9	308	545	556	571		
2275	SCHMIDT TROY KEVIN	10	M	7	9	359	537	549	556	565	569
2276	SHAW MARGUERITE	10	F	5	8	309	558	561	571		
2277	SHETTON JOHN TRIY	10	M	5	9	308	543	555	571		
2278	SHERIDAN LORI L	10	F	4	550	558	561	571			
2279	HSIPMAN GENE ALLEN	10	M	7	9	323	537	556	565	569	578

2280	SHOUBERT KRISTI LEIGH	10	F	6	537	556	564	565	569	577		
2281	SIDES VICKY ANN	10	F	7	8	309	537	552	564	565	569	
2282	SIMMONS HELEN S	10	F	8	9	24	323	399	537	556	565	569
2283	SIMSON DAVID W	10	M	7	9	323	537	548	556	565	569	
2284	SKIRLPCE JANESE GAL	10	F	5	8	323	546	558	571			
2285	SMITH DRENDA EILEEN	10	F	6	531	537	544	552	565	569		
2286	SMITH GLORIA DEAN	10	F	6	9	16	306	399	566	572		
2287	SMITH TAMARA KAY	10	F	4	531	558	561	571				
2288	SMITH YVONNE N	10	F	7	8	309	532	537	558	565	569	
2289	SNIDER CYNTHIA KAY	10	F	5	12	309	558	561	571			
2290	SNOW SHARON J	10	F	7	9	364	537	550	556	565	569	
2291	SPANN DANUI L	10	F	7	12	309	537	558	564	565	569	
2292	STANDIFER MICHAEL J	10	M	5	8	309	550	558	571			
2293	ST CLAIKE STEVE GLEN	10	M	7	60	309	530	537	556	565	569	
2294	STEELY MICHAEL E	10	M	6	537	548	556	559	565	569		
2295	STEVES CALVIN L	10	M	3	529	570	585					
2296	STEVES CLINTON C	10	M	4	9	305	549	570				
2297	STRATEGIER KIMBERLY	10	F	5	9	399	555	559	571			
2299	STREET CLINTON LEDN	10	M	6	537	549	551	556	565	569		
2299	SWANN MICHAEL K	10	M	7	16	308	537	552	565	569	575	
2300	TABU SHELVE R	10	M	7	23	309	537	549	552	565	569	
2301	TADD TERRY ELAINE	10	F	5	9	362	547	566	572			
2302	TALIFERRO KEITH LEE	10	M	4	9	385	548	570				
2303	TATE CARTER LEE	10	M	4	9	305	548	570				
2304	TAYLOR SALLY ELIZABETG	10	F	7	19	309	537	547	558	565	569	
2305	TAYLOR SANDRA KAY	10	F	7	16	367	537	552	564	565	569	
2306	TEE DOREEN JU	10	F	7	23	309	537	550	558	565	569	
2307	TEVAULT BAHDI LYNET	10	F	7	99	309	531	537	552	565	569	
2308	THOMPSON SHERI LEE	10	F	5	12	309	558	571	575			
2309	TILKENS ANNA JOYCE	10	F	7	9	308	537	550	556	565	569	
2310	TODD JOHN ERIC	10	M	7	5	309	537	552	559	565	569	
2311	TCMLIN GENA LOUISE	10	F	5	8	309	556	562	571			
2312	TOUHEY DEURIE SUE	10	F	5	64	309	550	566	572			
2313	TOWNSEND JERRY DON	10	M	7	16	309	537	550	552	565	569	
2314	TRAVIS JAN THARESA	10	F	6	530	534	635	557	568	581		
2315	TURNER JERRY D	10	M	5	24	309	558	560	571			
2316	VAIL DIANA LYNNE	10	F	7	12	309	537	558	559	565	569	
2317	VALENTINE MARK S	10	M	7	7	399	537	549	556	565	569	
2318	VANDERHELM BECKY S	10	F	7	9	323	537	541	558	565	569	
2319	VANHORN THOMAS H	10	M	2	566	572						
2320	VAUGHAN HOLLY JENIS	10	F	6	537	539	552	561	565	569		
2321	VO HCANE ANH THI	10	F	6	531	537	539	558	565	569		
2322	WADDLLE MICHAEL D	10	M	7	9	323	537	549	552	565	569	
2323	WADDLE SHELLEY SUE	10	F	7	64	309	531	537	558	565	569	
2324	WALKER HAS BRETT	10	M	4	560	567	572	575				
2325	WALKUP SCOTT DANIEL	10	M	7	23	309	537	556	565	569	579	
2326	WALLACE DAVID EAVEN	10	M	7	7	309	537	545	556	565	569	
2327	WALSH DAVID GLENN	10	M	7	7	309	537	551	556	565	569	
2328	WARREN SANDY KAY	10	F	7	5	309	537	558	561	565	569	
2329	WATFORD JOHN CLAUDE	10	M	6	23	309	552	560	566	569		
2330	WEBB ROBERT L	10	M	5	99	309	560	566	572			
2331	WELCH DOUGLAS GRANT	10	M	6	537	550	558	559	565	569		
2332	WELLS MARY BETH	10	F	5	12	309	558	564	571			
2333	WHEELER TODD VER	10	M	4	559	565	572	578				
2334	WHITEHEAD ROY DALE	10	M	7	24	309	537	549	556	565	569	
2335	WHITMORE JAMES E	10	M	7	23	325	537	552	565	569	575	
2336	WICKHAM NOVELLA LYN	10	F	4	531	550	558	571				
2337	WIGGINS BICKY DALE	10	M	4	560	566	572	579				
2338	WILKINSON GINGER P	10	F	7	9	308	537	550	552	565	569	
2339	WILLCOX BILLY DCN	10	M	5	534	557	560	568	581			
2340	WILLIAMS G KENT	10	M	7	9	323	537	551	552	565	569	
2344	WILLIAMS MARVIN RAY	10	M	6	537	548	556	565	569	579		
2345	WILLIAMS ROBIN LYNN	10	F	7	12	309	531	537	558	565	569	
2346	WILLIAMSON LICIA NU	10	F	7	9	316	537	556	561	565	569	
2347	WILLIFORD DOUGLAS L	10	M	6	531	537	552	559	565	569		
2348	WILLSON MIKE LEE	10	M	5	9	323	548	566	572			

2349	WILSON ALLEN ODEAN	10	H	7	9	316	537	551	555	565	569
2350	WILSON JENNIE LLU	10	F	4	539	558	571	575			
2351	WIMSETT DENISE LEE	10	F	7	5	309	531	537	558	565	569
2352	WINES STEVE VINEYARD	10	H	6	530	533	557	559	565	569	
2353	WITT DONNA GAYLE	10	F	7	12	309	531	537	550	558	569
2354	WOLF DAVID WAYNE	10	H	6	530	531	537	556	565	569	
2355	WOMACK SUSAN RHENEA	10	F	5	9	316	558	561	571		
2356	WOOD CECILIA ANN	10	F	5	12	309	558	561	571		
2357	WOODCOCK PATRICIA D	10	F	7	9	399	537	556	565	569	577
2358	YAHN KIMBERLY DIAN	10	F	5	9	317	552	564	571		
2359	YATES THEODORE THOM	10	H	7	9	323	537	557	559	565	569
2360	ZIBULSKY LISA KAY	10	F	7	25	309	537	558	565	569	580
2361	ZWICK EVAN	10	H	7	19	309	537	557	565	569	575

APPENDIX III

PART B. OUTPUT LISTINGS

NUMBER OF COURSES OFFERED	68	NUMBER OF STUDENTS	714
NO. OF COURSES HAVING NO REQUESTS	6	NO OF STUDENTS WITH INFEASIBLE REQUEST LISTS	12
NUMBER OF COURSES TO BE SCHEDULED	62	NUMBER OF STUDENTS TO BE SCHEDULED	702

COURSES ARRANGED ACCORDING TO INDICES OF DIFFICULTY

COURSE	SECTIONS	TYPE	SEMESTER	INDEX OF DIFFICULTY	COURSE AND TIME RESTRICTIONS
529	MSEC= 1 ISEC= 1 NJ= 0 NCR1 NG1 2295	ITYP= 0 ISEM= 0 IDF= 0 NCR2 NG2 NCR3 NG3	0 0 0	ITME= 0 ITIME1= 6 ITIME2= 0 ITIME3= 0 ITIME4= 0	K1= 1
538	MSEC= 1 ISEC= 1 NJ= 0 NCR1 NG1 1023 1054 1064 1317 1336 1355	ITYP= 0 ISEM= 0 IDF= 0 NCR2 NG2 NCR3 NG3	0 0 0	ITME= 0 ITIME1= 3 ITIME2= 0 ITIME3= 0 ITIME4= 0	K1= 22 1265 1266 1269 2170 2181 2200
540	MSEC= 1 ISEC= 1 NJ= 0 NCR1 NG1 1004 1095 1115	ITYP= 0 ISEM= 0 IDF= 0 NCR2 NG2 NCR3 NG3	0 0 0	ITME= 0 ITIME1= 4 ITIME2= 0 ITIME3= 0 ITIME4= 0	K1= 11 1271 1276 1328
541	MSEC= 1 ISEC= 1 NJ= 0 NCR1 NG1 2030 2096 2109	ITYP= 0 ISEM= 0 IDF= 0 NCR2 NG2 NCR3 NG3	0 0 0	ITME= 0 ITIME1= 4 ITIME2= 0 ITIME3= 0 ITIME4= 0	K1= 4
550	MSEC= 1 ISEC= 1 NJ= 0 NCR1 NG1 1009 1032 1042 1135 1137 1165 1253 1258 1272 2092 2094 2110 2235 2255 2258 2331 2336 2338	ITYP= 0 ISEM= 0 IDF= 0 NCR2 NG2 NCR3 NG3	0 0 0	ITME= 0 ITIME1= 5 ITIME2= 0 ITIME3= 0 ITIME4= 0	K1= 64 1087 1095 1113 1126 1134 1205 1206 1209 1216 1221 2056 2062 2069 2083 2088 2155 2192 2210 2229 2230 2292 2306 2309 2312 2313
551	MSEC= 1 ISEC= 1 NJ= 0 NCR1 NG1 1013 1018 1034 1263 1270 1309 2298 2327 2340	ITYP= 0 ISEM= 0 IDF= 0 NCR2 NG2 NCR3 NG3	0 0 0	ITME= 0 ITIME1= 4 ITIME2= 0 ITIME3= 0 ITIME4= 0	K1= 28 1178 1185 1202 1231 1255 2130 2136 2153 2169 2256
559	MSEC= 1 ISEC= 1 NJ= 0 NCR1 NG1 1002 1003 1005 1061 1062 1063 1116 1125 1138 1209 1211 1220 1260 1261 1267 1314 1315 1319 2045 2049 2056 2099 2142 2143 2218 2226 2227 2294 2297 2310	ITYP= 0 ISEM= 0 IDF= 0 NCR2 NG2 NCR3 NG3	0 0 0	ITME= 0 ITIME1= 6 ITIME2= 0 ITIME3= 0 ITIME4= 0	K1= 117 1049 1051 1052 1059 1060 1088 1092 1093 1101 1104 1168 1181 1188 1196 1201 1241 1244 1245 1250 1254 1302 1303 1305 1307 1310 1347 2006 2007 2015 2017 2071 2072 2076 2091 2097 2105 2193 2203 2207 2217 2251 2261 2263 2266 2272
561	MSEC= 1 ISEC= 1 NJ= 0 NCR1 NG1	ITYP= 0 ISEM= 0 IDF= 0 NCR2 NG2 NCR3 NG3	0 0 0	ITME= 0 ITIME1= 6 ITIME2= 0 ITIME3= 0 ITIME4= 0	K1= 54

1016	1046	1053	1065	1067	1070	1071	1089	1105	1115	1120	1127
1147	1161	1180	1186	1199	1202	1237	1249	1258	1268	1275	1282
1289	1308	1311	1322	1325	1341	2001	2009	2027	2057	2058	2082
2129	2146	2148	2172	2188	2195	2198	2200	2273	2276	2278	2287
2289	2320	2328	2346	2355	2356						

578 MSEC= 1 ISEC= 1 NJ= 0 ITP= 0 ISEM= 0 IDF= 0 ITME= 0 ITIME1= 1 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 26
 1005 1024 1027 1029 1086 1158 1167 1171 1196 1198 1212 1225
 1251 1327 2021 2022 2071 2108 2144 2164 2232 2245 2246 2263
 2279 2333

579 MSEC= 1 ISEC= 1 NJ= 0 ITP= 0 ISEM= 0 IDF= 0 ITME= 0 ITIME1= 1 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 14
 2047 2052 2089 2090 2116 2123 2153 2173 2196 2250 2254 2325
 2337 2344

12 MSEC= 1 ISEC= 1 NJ= 0 ITP= 0 ISEM= 1 IDF= 0 ITME= 0 ITIME1= 3 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 19
 2001 2006 2009 2092 2094 2095 2119 2150 2198 2220 2262 2289
 2291 2308 2316 2332 2345 2353 2356

570 MSEC= 1 ISEC= 1 NJ= 0 ITP= 4 ISEM= 0 IDF= 0 ITME= 0 ITIME1= 1 ITIME2= 1 ITIME3= 4 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 15
 1015 1111 1117 1210 1274 1354 2078 2100 2104 2106 2190 2295
 2296 2302 2303

571 MSEC= 1 ISEC= 1 NJ= 0 ITP= 3 ISEM= 0 IDF= 0 ITME= 0 ITIME1= 1 ITIME2= 1 ITIME3= 3 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 68
 1197 1314 2004 2005 2088 2014 2024 2029 2030 2031 2046 2069
 2070 2081 2084 2088 2095 2098 2099 2112 2119 2124 2125 2129
 2130 2131 2140 2142 2143 2144 2146 2150 2170 2174 2181 2187
 2191 2192 2193 2198 2200 2203 2211 2213 2225 2228 2234 2237
 2262 2269 2274 2276 2277 2278 2284 2287 2289 2292 2297 2308
 2311 2315 2332 2336 2350 2355 2356 2358

572 MSEC= 1 ISEC= 1 NJ= 0 ITP= 3 ISEM= 0 IDF= 0 ITME= 0 ITIME1= 1 ITIME2= 4 ITIME3= 6 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 50
 2010 2018 2020 2021 2022 2026 2027 2032 2034 2036 2043 2050
 2052 2054 2064 2066 2071 2073 2075 2091 2113 2136 2141 2147
 2151 2160 2186 2197 2199 2206 2208 2216 2221 2232 2233 2236
 2242 2245 2248 2254 2263 2286 2301 2312 2319 2324 2330 2333
 2337 2344

573 MSEC= 1 ISEC= 1 NJ= 0 ITP= 3 ISEM= 0 IDF= 0 ITME= 0 ITIME1= 1 ITIME2= 1 ITIME3= 3 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 69
 1003 1004 1013 1023 1025 1029 1032 1033 1034 1043 1044 1056
 1059 1060 1063 1064 1076 1077 1088 1099 1106 1109 1110 1122
 1123 1134 1143 1148 1151 1153 1154 1164 1165 1168 1172 1174
 1176 1178 1185 1186 1189 1190 1206 1222 1227 1239 1244 1245
 1246 1250 1253 1254 1256 1258 1263 1265 1286 1298 1304 1308
 1311 1315 1319 1320 1322 1325 1326 1328 1333

19 MSEC= 1 ISEC= 1 NJ= 0 ITP= 1 ISEM= 1 IDF= 0 ITME= 1 ITIME1= 1 ITIME2= 4 ITIME3= 5 ITIME4= 2
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 21
 1034 1082 1085 1103 1123 1124 1159 1174 1187 1239 1256 1269
 1303 1318 1333 2041 2076 2245 2258 2304 2361

21 MSEC= 1 ISEC= 1 NJ= 0 ITP= 1 ISEM= 1 IDF= 0 ITME= 0 ITIME1= 1 ITIME2= 4 ITIME3= 5 ITIME4= 2
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 22
 1013 1053 1071 1090 1112 1177 1180 1184 1203 1219 1232 1235
 1315 1353 2014 2037 2050 2099 2105 2157 2169 2239

320 MSEC= 1 ISEC= 1 NJ= 0 ITP= 1 ISEM= 2 IDF= 0 ITME= 1 ITIME1= 1 ITIME2= 4 ITIME3= 5 ITIME4= 2
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 10
 1034 1082 1085 1123 1124 1159 1239 1256 1269 1333

321	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 2	IDF= 0	ITME= 0	ITIME1= 1	ITIME2= 4	ITIME3= 6	ITIME4= 2
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 15		
			1006 1013 1039	1053 1071 1090	1177 1180	1184 1203	1232 1235				
			1315 1318 1353								
575	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 0	IDF= 0	ITME= 1	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 2
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 48		
			1022 1055 1081	1105 1108 1109	1110 1149	1154 1190	1201 1203				
			1208 1210 1262	1337 1346 2012	2028 2030	2038, 2039	2051 2055				
			2075 2078 2095	2105 2107 2115	2117 2123	2132 2134	2145 2171				
			2182 2183 2189	2204 2222 2231	2299 2308	2324 2335	2350 2361				
577	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 0	IDF= 0	ITME= 0	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 26		
			1018 1066 1097	1182 1239 1242	1259 1264	1283 1287	1289 1293				
			1318 1351 2040	2040 2092 2118	2152 2221	2222 2224	2234 2257				
			2280 2357								
16	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 1	IDF= 90104	ITME= 0	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 17		
			1079 1125 1333	2012 2018 2027	2077 2151	2157 2158	2165 2227				
			2260 2286 2299	2305 2313							
532	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 0	IDF= 90086	ITME= 0	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 17		
			2037 2046 2090	2108 2126 2140	2170 2174	2202 2211	2212 2215				
			2220 2230 2242	2249 2280							
316	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 2	IDF= 90075	ITME= 0	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 12		
			1067 1189 1245	1284 2026 2036	2081 2093	2238 2346	2349 2355				
7	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 1	IDF= 90073	ITME= 0	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 16		
			1244 2057 2059	2074 2080 2120	2122 2127	2185 2214	2226 2256				
			2271 2317 2326	2327							
306	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 2	IDF= 90072	ITME= 0	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 13		
			1041 1093 1090	1135 1175 1257	1274 1290	1324 1335	1352 2162				
			2286								
555	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 0	IDF= 90056	ITME= 0	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 7		
			2056 2191 2200	2269 2277 2297	2349						
547	MSEC= 1	ISEC= 1	NJ= 0	ITYP= 1	ISEM= 0	IDF= 90034	ITME= 0	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 6		
			2100 2190 2221	2238 2301 2304							
534	MSEC= 3	ISEC= 3	NJ= 0	ITYP= 1	ISEM= 0	IDF= 64253	ITME= 0	ITIME1= 1	ITIME2= 4	ITIME3= 6	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 97		
			1002 1005 1006	1007 1008 1011	1014 1019	1021 1024	1026 1027				
			1035 1041 1046	1048 1049 1050	1057 1058	1061 1062	1072 1089				
			1091 1093 1094	1108 1116 1118	1121 1125	1131 1136	1138 1139				
			1141 1144 1145	1150 1152 1158	1160 1161	1167 1170	1171 1173				
			1181 1182 1191	1193 1195 1196	1211 1212	1213 1215	1216 1217				
			1218 1219 1221	1226 1231 1240	1241 1249	1255 1257	1266 1267				
			1270 1278 1290	1293 1301 1305	1306 1312	1313 1321	1331 1334				
			1336 1337 1338	1340 1343 1346	1348 1350	1352 1355	2042 2314				
			2339								
581	MSEC= 3	ISEC= 3	NJ= 0	ITYP= 1	ISEM= 0	IDF= 64247	ITME= 0	ITIME1= 1	ITIME2= 4	ITIME3= 6	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 93		
			1002 1005 1006	1007 1008 1011	1014 1019	1021 1024	1026 1027				

568	MSEC= 3	1SEC= 3	NJ= 0	11YP= 1	1SEM= 0	IDF= 64244	TIME= 0	TIME1= 1	TIME2= 4	TIME3= 6	TIME4= 0	1035	1041	1048	1049	1050	1057	1058	1061	1062	1072	1086	1138	1171	1216	1267	1270	1336	1338	
													1089	1091	1094	1108	1116	1158	1160	1161	1167	1136	1173	1217	1219	1270	1336	1338		
													1139	1141	1144	1150	1152	1158	1160	1161	1167	1136	1173	1217	1219	1270	1336	1338		
													1181	1182	1191	1193	1195	1196	1211	1213	1215	1216	1217	1219	1270	1336	1338	1338		
													1221	1226	1231	1240	1249	1255	1257	1266	1267	1270	1278	1278	1336	1336	1336	1338		
													1290	1293	1301	1305	1306	1312	1313	1321	1331	1334	1336	1338	1338	1338	1338	1338		
													1340	1343	1346	1348	1350	1352	1355	2042	2314	2339								
530	MSEC= 2	1SEC= 2	NJ= 0	11YP= 1	1SEM= 0	IDF= 64206	TIME= 1	TIME1= 1	TIME2= 0	TIME3= 0	TIME4= 1	1008	1015	1021	1029	1031	1036	1043	1057	1062	1065	1069								
													1008	1015	1021	1029	1031	1036	1043	1057	1062	1065	1069							
													1072	1074	1076	1083	1087	1098	1102	1141	1142	1149	1156							
													1168	1177	1179	1187	1188	1195	1200	1205	1222	1233								
													1234	1236	1247	1259	1260	1261	1276	1295	1301	1311	1323	1331						
													1334	1339	1349	2011	2084	2166	2168	2183	2202	2208	2215	2223						
													2260	2293	2314	2352	2354													
5	MSEC= 2	1SEC= 2	NJ= 0	11YP= 1	1SEM= 1	IDF= 64205	TIME= 1	TIME1= 1	TIME2= 0	TIME3= 0	TIME4= 3	1039	1041	1090	1093	1094	1096	1100	1107	1135	1148	1156	1160							
													1039	1041	1090	1093	1094	1096	1100	1107	1135	1148	1156	1160						
													1175	1197	1210	1226	1243	1255	1257	1274	1279	1286	1288							
													1290	1306	1312	1313	1324	1330	1335	1340	1342	1345	1352							
													2022	2032	2054	2066	2121	2137	2161	2175	2176	2209	2310	2328						
													2351																	
562	MSEC= 2	1SEC= 2	NJ= 0	11YP= 1	1SEM= 0	IDF= 64176	TIME= 1	TIME1= 1	TIME2= 0	TIME3= 0	TIME4= 1	1050	1164	1242	1253	2054	2147	2311												
													1050	1164	1242	1253	2054	2147	2311											
564	MSEC= 2	1SEC= 2	NJ= 0	11YP= 1	1SEM= 0	IDF= 64175	TIME= 1	TIME1= 1	TIME2= 0	TIME3= 0	TIME4= 3	1012	1031	1039	1042	1091	1094	1097	1108	1111	1124	1154								
													1012	1031	1039	1042	1091	1094	1097	1108	1111	1124	1154							
													1195	1216	1217	1230	1238	1246	1252	1264	1278	1283	1332	1346						
													1351	2032	2086	2107	2126	2128	2157	2167	2179	2201	2204	2205						
													2208	2234	2239	2257	2262	2280	2281	2291	2305	2332	2358							
546	MSEC= 2	1SEC= 2	NJ= 0	11YP= 1	1SEM= 0	IDF= 64155	TIME= 1	TIME1= 1	TIME2= 0	TIME3= 0	TIME4= 1	1001	1015	1027	1033	1035	1036	1038	1045	1080	1083	1096	1106							
													1001	1015	1027	1033	1035	1036	1038	1045	1080	1083	1096	1106						
													1110	1111	1112	1130	1131	1136	1145	1147	1155	1167	1173	1176						
													1178	1182	1193	1194	1199	1213	1218	1222	1235	1249	1278	1288						
													1258	1300	1301	1312	1331	2135	2244											
324	MSEC= 2	1SEC= 2	NJ= 0	11YP= 1	1SEM= 2	IDF= 64154	TIME= 1	TIME1= 1	TIME2= 0	TIME3= 0	TIME4= 3	1014	1030	1075	1091	1094	1112	1148	1156	1160	1169	1174	1206							
													1014	1030	1075	1091	1094	1112	1148	1156	1160	1169	1174	1206						
													1240	1257	1279	1282	1286	1312	1313	1330	1343	2054	2058							
													2145	2151	2167	2174	2211	2216	2240	2266										
24	MSEC= 2	1SEC= 2	NJ= 0	11YP= 1	1SEM= 1	IDF= 64152	TIME= 1	TIME1= 1	TIME2= 0	TIME3= 0	TIME4= 3	1006	1040	1045	1077	1142	1162	1166	1189	1246	1272	1304	1316							
													1006	1040	1045	1077	1142	1162	1166	1189	1246	1272	1304	1316						
													1319	1320	1328	2017	2040	2077	2093	2125	2128	2133								
													2140	2147	2148	2179	2194	2217	2219	2252	2261	2282	2316	2334						

305	MSEC= 2	ISEC= 2	NJ= 0	ITYP= 1	ISEM= 2	IDF= 64103	ITME= 1	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 3
			NCR1 NG1	NCR2 NG2	NCR3 NG3	306	530	0	K1= 20		
			1045 1224 1277	2010	2031 2034 2039	2043	2045	2067	2082	2104	
			2106 2159 2190	2194	2214 2253 2296	2303					
8	MSEC= 2	ISEC= 2	NJ= 0	ITYP= 1	ISEM= 1	IDF= 60268	ITME= 0	ITIME1= 1	ITIME2= 1	ITIME3= 4	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 49		
			1014 1030 1056	1067	1075 1091 1121	1146	1169	1172	1191	1206	
			1224 1228 1245	1277	1284 1343 1344	2002	2007	2010	2049	2062	
			2065 2069 2087	2117	2144 2151 2161	2181	2182	2194	2203	2232	
			2236 2237 2247	2254	2263 2266 2269	2276	2281	2284	2288	2292	
			2311								
308	MSEC= 2	ISEC= 2	NJ= 0	ITYP= 1	ISEM= 2	IDF= 60252	ITME= 0	ITIME1= 1	ITIME2= 1	ITIME3= 4	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 57		
			1040 1077 1079	1100	1103 1104 1125	1129	1132	1152	1166	1197	
			1246 1254 1292	1303	1304 1316 1319	1333	1334	1342	1342	1345	
			2016 2024 2034	2038	2040 2072 2077	2107	2112	2115	2131	2149	
			2157 2165 2169	2178	2192 2196 2199	2207	2210	2212	2241	2242	
			2243 2246 2251	2260	2274 2277 2299	2309	2338				
323	MSEC= 2	ISEC= 2	NJ= 0	ITYP= 1	ISEM= 2	IDF= 60160	ITME= 1	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 2
			NCR1 NG1	NCR2 NG2	NCR3 NG3	575	577	0	K1= 35		
			1054 1050 1107	1121	1139 1146 1226	1228	1244	1255	1306	1328	
			1344 2007 2019	2025	2034 2057 2137	2141	2161	2185	2193	2223	
			2231 2247 2279	2282	2283 2284 2318	2322	2340	2348	2359		
23	MSEC= 2	ISEC= 2	NJ= 0	ITYP= 1	ISEM= 1	IDF= 60101	ITME= 1	ITIME1= 0	ITIME2= 0	ITIME3= 0	ITIME4= 2
			NCR1 NG1	NCR2 NG2	NCR3 NG3	575	577	0	K1= 24		
			1024 1073 1104	1132	1152 1254 1257	1292	1334	2003	2028	2033	
			2035 2046 2093	2110	2136 2138 2267	2300	2306	2325	2329	2335	
309	MSEC= 4	ISEC= 4	NJ= 0	ITYP= 1	ISEM= 2	IDF= 56697	ITME= 0	ITIME1= 1	ITIME2= 1	ITIME3= 4	ITIME4= 0
			NCR1 NG1	NCR2 NG2	NCR3 NG3	0	0	0	K1= 133		
			2001 2002 2003	2006	2008 2009 2014	2015	2017	2018	2021	2022	
			2027 2028 2029	2032	2033 2035 2037	2040	2041	2046	2049	2050	
			2052 2059 2060	2062	2063 2064 2065	2069	2074	2076	2077	2090	
			2092 2093 2094	2095	2097 2099 2101	2105	2110	2117	2119	2120	
			2121 2122 2124	2125	2127 2133 2134	2136	2138	2140	2144	2148	
			2150 2151 2158	2161	2164 2171 2175	2176	2179	2181	2182	2194	
			2198 2203 2209	2217	2219 2220 2225	2226	2227	2229	2232	2235	
			2236 2237 2239	2245	2249 2252 2256	2258	2260	2261	2262	2263	
			2265 2267 2269	2271	2276 2281 2288	2289					

1234	1235	1237	1243	1249	1252	1260	1261	1262	1265	1268	1269
1275	1276	1277	1280	1281	1283	1284	1287	1288	1289	1290	1294
1302	1306	1307	1313	1316	1317	1319	1321	1323	1339	1341	1342
1342	1343	1345	1351	1355	2080						

549 MSEC= 2 ISEC= 2 NJ= 0 ITP= 1 ISEM= 0 IDF= 56205 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 32
 2003 2011 2024 2035 2044 2064 2085 2101 2116 2121 2132 2138
 2141 2149 2160 2186 2191 2197 2206 2225 2243 2250 2259 2267
 2272 2275 2296 2298 2300 2317 2322 2334

548 MSEC= 3 ISEC= 3 NJ= 0 ITP= 1 ISEM= 0 IDF= 56204 ITME= 1 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 2
 NCR1 NG1 NCR2 NG2 NCR3 NG3 549 0 0 K1= 83
 1003 1004 1008 1010 1019 1026 1028 1032 1051 1057 1060 1061
 1074 1078 1084 1101 1113 1116 1117 1118 1120 1122 1128 1129
 1132 1138 1140 1144 1150 1152 1155 1158 1162 1163 1170 1171
 1179 1204 1207 1211 1212 1214 1215 1220 1223 1226 1228 1229
 1241 1243 1247 1248 1263 1273 1284 1285 1290 1292 1294 1299
 1302 1305 1307 1313 1321 1326 1329 1338 1343 1344 1350 1354
 2020 2087 2091 2120 2153 2283 2294 2302 2303 2344 2348

560 MSEC= 2 ISEC= 2 NJ= 0 ITP= 1 ISEM= 0 IDF= 56163 ITME= 1 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 1
 NCR1 NG1 NCR2 NG2 NCR3 NG3 559 0 0 K1= 45
 1017 1028 1086 1107 1128 1136 1144 1159 1170 1191 1214 1215
 1221 1251 1271 1285 1306 1317 1323 1327 1329 1330 1350 1352
 2012 2016 2023 2026 2033 2042 2046 2127 2160 2184 2186 2197
 2219 2241 2254 2315 2324 2329 2330 2337 2339

557 MSEC= 4 ISEC= 4 NJ= 0 ITP= 1 ISEM= 0 IDF= 52370 ITME= 1 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 2
 NCR1 NG1 NCR2 NG2 NCR3 NG3 8 0 0 K1= 119
 1002 1035 1007 1008 1011 1019 1021 1024 1026 1027 1035 1036
 1041 1045 1048 1049 1050 1055 1058 1062 1068 1072 1074 1080
 1081 1086 1091 1094 1096 1098 1108 1116 1118 1120 1121 1124
 1128 1131 1135 1138 1141 1144 1145 1149 1150 1158 1160 1167
 1170 1171 1172 1173 1177 1181 1182 1183 1184 1191 1193 1195
 1196 1198 1202 1203 1204 1208 1211 1212 1213 1214 1216 1217
 1219 1221 1226 1233 1238 1240 1241 1248 1255 1257 1259 1266
 1267 1270 1271 1273 1278 1282 1285 1293 1300 1301 1303 1305
 1312 1327 1330 1331 1334 1336 1337 1338 1340 1346 1347 1348
 1350 1352 2012 2042 2185 2204 2314 2339 2352 2359 2361

531 MSEC= 4 ISEC= 4 NJ= 0 ITP= 1 ISEM= 0 IDF= 48516 ITME= 1 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 1
 NCR1 NG1 NCR2 NG2 NCR3 NG3 532 0 0 K1= 161
 1006 1007 1009 1012 1016 1017 1019 1020 1021 1022 1023 1025
 1030 1033 1035 1037 1038 1040 1044 1046 1047 1049 1050 1052
 1055 1056 1058 1063 1066 1075 1076 1077 1080 1081 1082 1085
 1089 1092 1098 1099 1100 1102 1103 1106 1109 1114 1118 1122
 1127 1130 1131 1137 1139 1140 1145 1148 1151 1153 1154 1161
 1165 1166 1173 1183 1185 1190 1197 1208 1213 1217 1218 1219
 1224 1225 1227 1230 1232 1237 1238 1250 1252 1253 1256 1265
 1266 1268 1270 1273 1275 1277 1279 1287 1298 1300 1304 1308
 1309 1314 1316 1320 1322 1325 1332 1336 1341 1342 1342 1345
 1349 1353 1355 2002 2013 2018 2032 2036 2047 2051 2055 2075
 2079 2081 2086 2096 2112 2124 2130 2133 2137 2152 2157 2159
 2162 2165 2172 2173 2187 2189 2191 2199 2205 2209 2213 2214
 2216 2218 2236 2253 2259 2265 2285 2287 2307 2321 2323 2336
 2345 2347 2351 2353 2354

558 MSEC= 5 ISEC= 5 NJ= 0 ITP= 1 ISEM= 0 IDF= 40763 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 145
 1032 1039 1056 1057 1059 1063 1092 1110 1134 1164 1174 1186
 1227 1245 1246 1250 1254 1263 1311 1320 1322 1325 1326 1328
 1332 1333 2001 2002 2004 2006 2007 2008 2009 2014 2015 2023
 2024 2025 2029 2030 2031 2033 2041 2045 2046 2048 2055 2059
 2063 2067 2070 2072 2076 2079 2081 2084 2088 2092 2094 2095
 2099 2101 2110 2112 2115 2119 2121 2122 2124 2125 2129 2130

2140	2142	2143	2144	2148	2150	2152	2155	2159	2163	2165	2173
2174	2175	2176	2178	2181	2187	2192	2193	2198	2203	2207	2210
2211	2212	2213	2217	2218	2220	2224	2225	2228	2229	2230	2234
2237	2239	2251	2252	2253	2255	2259	2260	2262	2271	2276	2278
2284	2287	2288	2289	2291	2292	2304	2306	2308	2315	2316	2318
2321	2323	2328	2331	2332	2336	2345	2350	2351	2353	2355	2356
2360											

556 MSEC= 5 ISEC= 5 NJ= 0 ITP= 1 ISEM= 0 IDF= 40358 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 90
 2003 2005 2011 2016 2028 2037 2038 2039 2040 2044 2047 2051
 2057 2058 2060 2069 2083 2085 2089 2090 2093 2105 2108 2109
 2111 2118 2120 2126 2127 2128 2131 2133 2135 2137 2138 2139
 2146 2149 2153 2154 2161 2162 2164 2166 2167 2168 2169 2179
 2182 2183 2186 2194 2195 2196 2201 2209 2214 2219 2222 2223
 2227 2246 2250 2256 2257 2261 2267 2272 2273 2274 2275 2279
 2280 2282 2283 2290 2293 2294 2298 2309 2311 2317 2325 2326
 2327 2334 2344 2346 2354 2357

535 MSEC= 5 ISEC= 5 NJ= 0 ITP= 1 ISEM= 0 IDF= 40028 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 15
 1002 1011 1026 1141 1150 1181 1193 1240 1267 1321 1338 2042
 2110 2149 2314

525 MSEC= 5 ISEC= 5 NJ= 0 ITP= 1 ISEM= 0 IDF= 40011 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 4
 1020 2125 2143 2201

569 MSEC= 7 ISEC= 7 NJ= 0 ITP= 1 ISEM= 0 IDF= 4218 ITME= 1 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 2
 NCR1 NG1 NCR2 NG2 NCR3 NG3 16 316 0 K1= 218
 1169 2001 2002 2003 2006 2007 2009 2011 2012 2013 2015 2016
 2017 2019 2023 2025 2026 2033 2035 2037 2038 2039 2040 2041
 2044 2045 2047 2048 2049 2051 2055 2057 2058 2059 2060 2062
 2063 2065 2067 2072 2074 2076 2077 2079 2080 2082 2083 2085
 2086 2087 2089 2090 2092 2093 2094 2096 2097 2101 2105 2107
 2108 2109 2110 2111 2115 2116 2117 2118 2120 2121 2122 2123
 2126 2127 2128 2132 2133 2134 2135 2137 2138 2139 2145 2148
 2149 2152 2153 2154 2155 2158 2159 2161 2162 2163 2164 2165
 2166 2167 2168 2169 2171 2172 2173 2175 2176 2178 2179 2180
 2182 2183 2184 2185 2188 2189 2194 2195 2196 2201 2202 2204
 2205 2207 2209 2210 2212 2214 2215 2217 2218 2219 2220 2222
 2223 2224 2226 2227 2229 2230 2231 2235 2238 2239 2240 2241
 2243 2246 2247 2249 2250 2251 2252 2253 2255 2256 2257 2258
 2259 2260 2261 2265 2266 2267 2268 2271 2272 2273 2275 2279
 2280 2281 2282 2283 2285 2288 2290 2291 2293 2294 2298 2299
 2300 2304 2305 2306 2307 2309 2310 2313 2316 2317 2318 2320
 2321 2322 2323 2325 2326 2327 2328 2329 2331 2334 2335 2338
 2340 2344 2345 2346 2347 2349 2351 2352 2353 2354 2357 2359
 2360 2361

537 MSEC= 7 ISEC= 7 NJ= 0 ITP= 1 ISEM= 0 IDF= 4210 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0
 NCR1 NG1 NCR2 NG2 NCR3 NG3 0 0 0 K1= 210
 2001 2002 2006 2007 2009 2011 2013 2015 2016 2017 2019 2023
 2025 2028 2033 2035 2037 2038 2039 2040 2041 2044 2045 2047
 2048 2049 2051 2055 2056 2057 2058 2059 2062 2063 2065 2067
 2072 2074 2076 2077 2079 2080 2082 2083 2085 2086 2087 2089
 2090 2092 2093 2094 2096 2101 2105 2107 2108 2109 2110 2111
 2115 2116 2117 2118 2120 2121 2122 2123 2126 2128 2132 2133
 2134 2135 2137 2138 2139 2145 2148 2152 2153 2154 2155 2158
 2159 2161 2162 2163 2164 2165 2166 2167 2168 2169 2171 2172
 2173 2175 2176 2178 2179 2180 2182 2183 2185 2188 2189 2194
 2195 2196 2201 2202 2204 2205 2207 2209 2210 2212 2214 2215
 2216 2217 2218 2219 2220 2222 2223 2224 2226 2227 2229 2230
 2231 2235 2238 2239 2240 2241 2243 2246 2247 2249 2250 2251
 2252 2253 2255 2256 2257 2258 2259 2260 2261 2265 2266 2267
 2268 2271 2272 2273 2275 2279 2280 2281 2282 2283 2285 2288

2290	2291	2293	2294	2298	2299	2300	2304	2305	2306	2307	2309
2310	2313	2316	2317	2318	2320	2321	2322	2323	2325	2326	2327
2328	2331	2334	2335	2338	2340	2344	2345	2346	2347	2349	2351
2353	2354	2357	2359	2360	2361						

565- MSEC= 7 ISEC= 7 NJ= 0 ITP= 1 ISEM= 0 IDF= 4202 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0

	NCR1	NG1	NCR2	NG2	NCR3	NG3				K1=	202
2002	2003	2006	2007	2009	2012	2013	2015	2017	2019	2023	2025
2028	2033	2035	2037	2038	2039	2040	2041	2044	2045	2047	2048
2049	2051	2055	2058	2059	2063	2065	2066	2067	2072	2074	2076
2077	2079	2080	2082	2083	2085	2086	2087	2089	2090	2094	2096
2097	2101	2105	2107	2108	2109	2110	2111	2115	2116	2117	2120
2122	2123	2126	2127	2128	2132	2133	2134	2135	2137	2138	2139
2145	2148	2153	2155	2158	2159	2161	2162	2163	2164	2165	2166
2167	2168	2169	2171	2172	2173	2175	2176	2178	2179	2180	2182
2183	2184	2185	2188	2189	2194	2195	2196	2201	2202	2204	2205
2207	2208	2209	2210	2212	2217	2218	2219	2220	2223	2224	2226
2227	2229	2230	2231	2235	2238	2239	2240	2241	2243	2246	2247
2249	2250	2251	2252	2253	2255	2256	2257	2259	2260	2261	2265
2266	2267	2268	2271	2272	2273	2275	2279	2280	2281	2282	2283
2285	2288	2290	2291	2293	2294	2298	2299	2300	2304	2305	2306
2307	2309	2310	2313	2316	2317	2318	2320	2321	2322	2323	2325
2326	2327	2328	2329	2331	2333	2334	2335	2338	2340	2344	2345
2346	2347	2349	2351	2352	2354	2357	2359	2360	2361		

567 MSEC= 7 ISEC= 7 NJ= 0 ITP= 1 ISEM= 0 IDF= 4193 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0

	NCR1	NG1	NCR2	NG2	NCR3	NG3				K1=	193
1001	1009	1010	1012	1016	1017	1018	1020	1022	1028	1030	1031
1036	1037	1038	1039	1040	1042	1045	1047	1051	1052	1053	1054
1055	1065	1066	1067	1068	1069	1070	1071	1073	1074	1075	1078
1079	1080	1081	1082	1083	1084	1085	1087	1090	1092	1095	1096
1097	1098	1100	1101	1102	1103	1104	1105	1107	1112	1113	1114
1115	1119	1120	1124	1126	1127	1128	1129	1130	1132	1133	1135
1137	1140	1142	1146	1147	1149	1155	1156	1157	1159	1162	1163
1166	1170	1175	1177	1179	1180	1183	1184	1187	1188	1192	1194
1198	1199	1200	1201	1202	1203	1204	1205	1207	1208	1209	1214
1218	1220	1223	1224	1225	1228	1229	1230	1232	1233	1234	1235
1236	1237	1238	1242	1243	1247	1248	1251	1252	1259	1260	1261
1262	1264	1268	1269	1271	1272	1273	1275	1276	1277	1279	1280
1281	1282	1283	1284	1285	1287	1288	1289	1291	1292	1294	1295
1297	1299	1300	1302	1303	1307	1309	1310	1317	1318	1323	1324
1327	1329	1330	1332	1335	1337	1339	1341	1342	1342	1344	1345
1347	1349	1351	1353	2018	2021	2060	2062	2073	2093	2180	2242
2324											

580 MSEC= 8 ISEC= 8 NJ= 0 ITP= 1 ISEM= 0 IDF= 4191 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0

	NCR1	NG1	NCR2	NG2	NCR3	NG3				K1=	191
1001	1009	1010	1012	1016	1017	1018	1020	1022	1028	1030	1031
1036	1037	1038	1039	1040	1042	1045	1047	1051	1052	1053	1054
1055	1065	1066	1067	1068	1069	1070	1071	1073	1074	1075	1078
1079	1080	1081	1082	1083	1084	1085	1087	1090	1092	1095	1096
1097	1098	1100	1101	1102	1103	1104	1105	1107	1112	1113	1114
1115	1119	1120	1124	1126	1127	1128	1129	1130	1132	1133	1135
1137	1140	1142	1146	1147	1149	1155	1156	1157	1159	1162	1163
1166	1169	1170	1175	1177	1179	1180	1183	1184	1187	1188	1192
1194	1198	1199	1200	1201	1202	1203	1204	1205	1207	1208	1209
1214	1218	1220	1223	1224	1225	1228	1229	1230	1232	1233	1234
1235	1236	1237	1238	1242	1243	1247	1248	1251	1252	1259	1260
1261	1262	1264	1268	1269	1271	1272	1273	1275	1276	1277	1278
1279	1280	1281	1282	1283	1284	1285	1287	1288	1289	1291	1292
1294	1295	1297	1299	1300	1302	1303	1307	1309	1310	1316	1317
1318	1323	1324	1327	1329	1330	1332	1335	1337	1339	1341	1342
1342	1344	1345	1347	1348	1349	1351	1353	2050	2073	2360	

533 MSEC= 7 ISEC= 7 NJ= 0 ITP= 1 ISEM= 0 IDF= 4188 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0

	NCR1	NG1	NCR2	NG2	NCR3	NG3				K1=	188
							0	0	0		

1001	1009	1010	1012	1016	1017	1018	1020	1022	1028	1030	1031
1036	1037	1038	1039	1040	1042	1045	1047	1051	1052	1053	1054
1055	1065	1066	1067	1068	1069	1070	1071	1073	1074	1075	1078
1079	1080	1081	1082	1083	1084	1085	1086	1087	1090	1092	1095
1096	1097	1098	1100	1101	1102	1103	1104	1105	1107	1112	1113
1114	1115	1119	1120	1124	1126	1127	1128	1129	1130	1132	1133
1135	1137	1140	1142	1146	1147	1149	1155	1156	1157	1159	1162
1163	1166	1169	1175	1177	1179	1180	1183	1184	1187	1188	1192
1194	1198	1199	1200	1201	1202	1203	1204	1205	1207	1208	1209
1214	1220	1223	1224	1225	1228	1229	1230	1232	1233	1234	1235
1236	1237	1238	1242	1243	1247	1248	1251	1252	1259	1260	1261
1262	1264	1268	1269	1271	1272	1273	1275	1276	1277	1279	1280
1281	1282	1283	1284	1285	1287	1288	1289	1291	1292	1294	1295
1297	1299	1300	1302	1303	1307	1309	1310	1316	1317	1318	1323
1324	1327	1329	1330	1332	1335	1339	1341	1342	1342	1344	1345
1347	1349	1351	1353	2003	2127	2184	2352				

552 MSEC= 6 ISEC= 6 NJ= 0 IYP= 1 ISEM= 0 IOF= 4172 ITME= 0 ITIME1= 0 ITIME2= 0 ITIME3= 0 ITIME4= 0

	NCR1	NG1	NCR2	NG2	NCR3	NG3					
							0	0	0	K1=	172
1001	1004	1006	1012	1013	1017	1022	1023	1025	1029	1030	1031
1034	1037	1040	1042	1043	1044	1051	1052	1053	1060	1064	1065
1067	1070	1071	1075	1076	1078	1079	1084	1093	1099	1100	1102
1106	1107	1109	1112	1113	1114	1119	1122	1123	1130	1133	1142
1143	1146	1151	1153	1154	1157	1161	1162	1163	1165	1168	1169
1175	1176	1178	1185	1189	1190	1192	1194	1197	1199	1200	1201
1206	1209	1220	1222	1223	1225	1228	1229	1236	1239	1242	1244
1247	1251	1253	1256	1258	1264	1272	1279	1286	1291	1292	1295
1297	1298	1299	1304	1308	1309	1310	1314	1315	1318	1324	1329
1335	1344	1349	1353	2013	2017	2019	2035	2049	2062	2065	2074
2077	2082	2086	2096	2097	2116	2117	2123	2132	2134	2145	2157
2158	2170	2171	2172	2180	2184	2188	2189	2202	2205	2215	2226
2231	2235	2238	2240	2241	2243	2247	2249	2258	2265	2266	2268
2281	2285	2299	2300	2305	2307	2310	2313	2320	2322	2329	2335
2338	2340	2347	2358								

THE CONFLICT MATRIX FOR SINGLE SECTION COURSES

COURSE 1 (529)	AND	COURSE 2 (538)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 3 (540)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 4 (541)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 5 (550)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 6 (551)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 7 (559)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 8 (561)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 9 (578)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	1
COURSE 1 (529)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 1 (529)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 3 (540)	NUMBER OF CONFLICTS=	0

COURSE 2 (538)	AND	COURSE 4 (541)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 5 (550)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 6 (551)	NUMBER OF CONFLICTS=	1
COURSE 2 (538)	AND	COURSE 7 (559)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 8 (561)	NUMBER OF CONFLICTS=	2
COURSE 2 (538)	AND	COURSE 9 (578)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	1
COURSE 2 (538)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	3
COURSE 2 (538)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	6
COURSE 2 (538)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	3
COURSE 2 (538)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	2
COURSE 2 (538)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 2 (538)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	1
COURSE 2 (538)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	1
COURSE 2 (538)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 2 (538)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	1
COURSE 2 (538)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 4 (541)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 5 (550)	NUMBER OF CONFLICTS=	1
COURSE 3 (540)	AND	COURSE 6 (551)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 7 (559)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 8 (561)	NUMBER OF CONFLICTS=	1
COURSE 3 (540)	AND	COURSE 9 (578)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	0

COURSE 3 (540)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	3
COURSE 3 (540)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 3 (540)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 5 (550)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 6 (551)	NUMBER OF CONFLICTS=	1
COURSE 4 (541)	AND	COURSE 7 (559)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 8 (561)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 9 (578)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	1
COURSE 4 (541)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	1

COURSE 4 (541)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 4 (541)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 5 (550)	AND	COURSE 6 (551)	NUMBER OF CONFLICTS=	0
COURSE 5 (550)	AND	COURSE 7 (559)	NUMBER OF CONFLICTS=	3
COURSE 5 (550)	AND	COURSE 8 (561)	NUMBER OF CONFLICTS=	6
COURSE 5 (550)	AND	COURSE 9 (578)	NUMBER OF CONFLICTS=	0
COURSE 5 (550)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 5 (550)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	4
COURSE 5 (550)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 5 (550)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	12
COURSE 5 (550)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	1
COURSE 5 (550)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	10
COURSE 5 (550)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	2
COURSE 5 (550)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	1
COURSE 5 (550)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 5 (550)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	1
COURSE 5 (550)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	0
COURSE 5 (550)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 5 (550)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	1
COURSE 5 (550)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	1
COURSE 5 (550)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 5 (550)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	1
COURSE 5 (550)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	2
COURSE 5 (550)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	2
COURSE 5 (550)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 6 (551)	AND	COURSE 7 (559)	NUMBER OF CONFLICTS=	2
COURSE 6 (551)	AND	COURSE 8 (561)	NUMBER OF CONFLICTS=	1

COURSE 6 (551)	AND	COURSE 9 (578)	NUMBER OF CONFLICTS=	0
COURSE 6 (551)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 6 (551)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	0
COURSE 6 (551)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 6 (551)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	3
COURSE 6 (551)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	7
COURSE 6 (551)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	2
COURSE 6 (551)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 6 (551)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	2
COURSE 6 (551)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	1
COURSE 6 (551)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 7 (559)	AND	COURSE 8 (561)	NUMBER OF CONFLICTS=	0
COURSE 7 (559)	AND	COURSE 9 (578)	NUMBER OF CONFLICTS=	6
COURSE 7 (559)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 7 (559)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	3
COURSE 7 (559)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 7 (559)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	9
COURSE 7 (559)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	5
COURSE 7 (559)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	16
COURSE 7 (559)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	2
COURSE 7 (559)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	2
COURSE 7 (559)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 7 (559)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	1
COURSE 7 (559)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	1

COURSE 7 (559)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	0
COURSE 7 (559)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	3
COURSE 7 (559)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 7 (559)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	1
COURSE 7 (559)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	3
COURSE 7 (559)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	3
COURSE 7 (559)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	2
COURSE 7 (559)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 8 (561)	AND	COURSE 9 (578)	NUMBER OF CONFLICTS=	0
COURSE 8 (561)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 8 (561)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	5
COURSE 8 (561)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 8 (561)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	10
COURSE 8 (561)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	1
COURSE 8 (561)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	6
COURSE 8 (561)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	0
COURSE 8 (561)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	3
COURSE 8 (561)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 8 (561)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	3
COURSE 8 (561)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	1
COURSE 8 (561)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 8 (561)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	1
COURSE 8 (561)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 8 (561)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	3
COURSE 8 (561)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	1
COURSE 8 (561)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 8 (561)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	1
COURSE 8 (561)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 10 (579)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	1
COURSE 9 (578)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	7

COURSE 9 (578)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	1
COURSE 9 (578)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	1
COURSE 9 (578)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	1
COURSE 9 (578)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 9 (578)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 11 (12)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	3
COURSE 10 (579)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	1
COURSE 10 (579)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 10 (579)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 12 (570)	NUMBER OF CONFLICTS=	0

COURSE 11 (12)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	9
COURSE 11 (12)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	2
COURSE 11 (12)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 11 (12)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	1
COURSE 11 (12)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 11 (12)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 13 (571)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	2
COURSE 12 (570)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	1
COURSE 12 (570)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 12 (570)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	2
COURSE 13 (571)	AND	COURSE 14 (572)	NUMBER OF CONFLICTS=	0

COURSE 13 (571)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	0
COURSE 13 (571)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	0
COURSE 13 (571)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	2
COURSE 13 (571)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 13 (571)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 13 (571)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	4
COURSE 13 (571)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 13 (571)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 13 (571)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	6
COURSE 13 (571)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	2
COURSE 13 (571)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 13 (571)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 13 (571)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	5
COURSE 13 (571)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 14 (572)	AND	COURSE 15 (573)	NUMBER OF CONFLICTS=	0
COURSE 14 (572)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	1
COURSE 14 (572)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	1
COURSE 14 (572)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 14 (572)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 14 (572)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	2
COURSE 14 (572)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 14 (572)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	4
COURSE 14 (572)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	1
COURSE 14 (572)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	2
COURSE 14 (572)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 14 (572)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	1
COURSE 14 (572)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 14 (572)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	2
COURSE 15 (573)	AND	COURSE 16 (19)	NUMBER OF CONFLICTS=	6
COURSE 15 (573)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	2
COURSE 15 (573)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	5
COURSE 15 (573)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	2
COURSE 15 (573)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	4

COURSE 15 (573)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 15 (573)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	1
COURSE 15 (573)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 15 (573)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	2
COURSE 15 (573)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	1
COURSE 15 (573)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 15 (573)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 15 (573)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 16 (19)	AND	COURSE 17 (21)	NUMBER OF CONFLICTS=	0
COURSE 16 (19)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	10
COURSE 16 (19)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	1
COURSE 16 (19)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	1
COURSE 16 (19)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	2
COURSE 16 (19)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	1
COURSE 16 (19)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 16 (19)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 16 (19)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 16 (19)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 16 (19)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 16 (19)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	1
COURSE 17 (21)	AND	COURSE 18 (320)	NUMBER OF CONFLICTS=	0
COURSE 17 (21)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	12
COURSE 17 (21)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	2
COURSE 17 (21)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	0
COURSE 17 (21)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	1
COURSE 17 (21)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	1
COURSE 17 (21)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 17 (21)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 17 (21)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 17 (21)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 17 (21)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 18 (320)	AND	COURSE 19 (321)	NUMBER OF CONFLICTS=	0
COURSE 18 (320)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	0

COURSE 18 (320)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 18 (320)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	1
COURSE 18 (320)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 18 (320)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 18 (320)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 18 (320)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 18 (320)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 18 (320)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 19 (321)	AND	COURSE 20 (575)	NUMBER OF CONFLICTS=	1
COURSE 19 (321)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 19 (321)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 19 (321)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 19 (321)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 19 (321)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 19 (321)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 19 (321)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 19 (321)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 20 (575)	AND	COURSE 21 (577)	NUMBER OF CONFLICTS=	1
COURSE 20 (575)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	2
COURSE 20 (575)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 20 (575)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 20 (575)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 20 (575)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 20 (575)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 20 (575)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 21 (577)	AND	COURSE 22 (16)	NUMBER OF CONFLICTS=	0
COURSE 21 (577)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0
COURSE 21 (577)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 21 (577)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 21 (577)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 21 (577)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 21 (577)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	1
COURSE 22 (16)	AND	COURSE 23 (532)	NUMBER OF CONFLICTS=	0

COURSE 22 (16)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 22 (16)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 22 (16)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	1
COURSE 22 (16)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 22 (16)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 23 (532)	AND	COURSE 24 (316)	NUMBER OF CONFLICTS=	0
COURSE 23 (532)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 23 (532)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 23 (532)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 23 (532)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 24 (316)	AND	COURSE 25 (7)	NUMBER OF CONFLICTS=	0
COURSE 24 (316)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 24 (316)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	1
COURSE 24 (316)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	1
COURSE 25 (7)	AND	COURSE 26 (306)	NUMBER OF CONFLICTS=	0
COURSE 25 (7)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 25 (7)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 26 (306)	AND	COURSE 27 (555)	NUMBER OF CONFLICTS=	0
COURSE 26 (306)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0
COURSE 27 (555)	AND	COURSE 28 (547)	NUMBER OF CONFLICTS=	0

SCHEDULING
LOCKED-IN COURSES

J 1	SCHEDULING COURSE 529	NO CONFLICTS GENERATED				
J 1	KT= 6 N(KT)= 1 IKF(KT,N)= 0 IKP(KT,N(KT))= 1 NH= 0 IKC(KT,N(KT))= 529 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
J 2	SCHEDULING COURSE 538	NO CONFLICTS GENERATED				
J 2	KT= 3 N(KT)= 1 IKF(KT,N)= 0 IKP(KT,N(KT))= 2 NH= 0 IKC(KT,N(KT))= 538 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
J 3	SCHEDULING COURSE 540	NO CONFLICTS GENERATED				
J 3	KT= 4 N(KT)= 1 IKF(KT,N)= 0 IKP(KT,N(KT))= 3 NH= 0 IKC(KT,N(KT))= 540 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
J 4	SCHEDULING COURSE 541	NO CONFLICTS GENERATED				
J 4	KT= 4 N(KT)= 2 IKF(KT,N)= 0 IKP(KT,N(KT))= 4 NH= 0 IKC(KT,N(KT))= 541 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
J 5	SCHEDULING COURSE 550	NO CONFLICTS GENERATED				
J 5	KT= 5 N(KT)= 1 IKF(KT,N)= 0 IKP(KT,N(KT))= 5 NH= 0 IKC(KT,N(KT))= 550 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
CONFLICT STUDENT ID 2109 J= 6 SCHEDULING COURSE 551 TIME KT= 4 CONFLICTING COURSE 541						
J 6	KT= 4 N(KT)= 3 IKF(KT,N)= 0 IKP(KT,N(KT))= 6 NH= 0 IKC(KT,N(KT))= 551 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
J 7	SCHEDULING COURSE 559	NO CONFLICTS GENERATED				
J 7	KT= 6 N(KT)= 2 IKF(KT,N)= 0 IKP(KT,N(KT))= 7 NH= 0 IKC(KT,N(KT))= 559 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
J 8	SCHEDULING COURSE 561	NO CONFLICTS GENERATED				
J 8	KT= 6 N(KT)= 3 IKF(KT,N)= 0 IKP(KT,N(KT))= 8 NH= 0 IKC(KT,N(KT))= 561 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
J 9	SCHEDULING COURSE 578	NO CONFLICTS GENERATED				
J 9	KT= 1 N(KT)= 1 IKF(KT,N)= 0 IKP(KT,N(KT))= 9 NH= 0 IKC(KT,N(KT))= 578 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					
J 10	SCHEDULING COURSE 579	NO CONFLICTS GENERATED				
J 10	KT= 1 N(KT)= 2 IKF(KT,N)= 0 IKP(KT,N(KT))= 10 NH= 0 IKC(KT,N(KT))= 579 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0					

CONFLICT STUDENT ID 2001 J= 11 SCHEDULING COURSE 12 TIME KT= 3 CONFLICTING COURSE 538
J 11 KT= 3 N(KT)= 2 IKF(KT,N)= 0 IKP(KT,N(KT))=11 NH= 0 IKC(KT,N(KT))= 12 IKG(KT,N(KT))= MSE1 = 1 MSEN(KT,N(KT))= 1

SCHEDULING
SEVERAL-HOUR COURSES 11

J 12 COURSE 570 NUMBER OF HOURS 4

SCHEDULED AT PERIODS 1 2 3 4

J 12 SCHEDULING COURSE 570 NO CONFLICTS GENERATED

J 12 KT= 1 N(KT)= 3 IKF(KT,N)= 0 IKP(KT,N(KT))=12 NH= 3 IKC(KT,N(KT))= 570 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

J 12 SCHEDULING COURSE 570 NO CONFLICTS GENERATED

J 12 KT= 2 N(KT)= 1 IKF(KT,N)= 0 IKP(KT,N(KT))=12 NH= 3 IKC(KT,N(KT))= 570 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

J 12 SCHEDULING COURSE 570 NO CONFLICTS GENERATED

J 12 KT= 3 N(KT)= 3 IKF(KT,N)= 0 IKP(KT,N(KT))=12 NH= 3 IKC(KT,N(KT))= 570 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

J 12 SCHEDULING COURSE 570 NO CONFLICTS GENERATED

J 12 KT= 4 N(KT)= 4 IKF(KT,N)= 0 IKP(KT,N(KT))=12 NH= 3 IKC(KT,N(KT))= 570 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

J 13 COURSE 571 NUMBER OF HOURS 3

SCHEDULED AT PERIODS 1 2 3

CONFLICT STUDENT ID 2144 J= 13 SCHEDULING COURSE 571 TIME KT= 1 CONFLICTING COURSE 578

J 13 KT= 1 N(KT)= 4 IKF(KT,N)= 0 IKP(KT,N(KT))=13 NH= 2 IKC(KT,N(KT))= 571 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

J 13 SCHEDULING COURSE 571 NO CONFLICTS GENERATED

J 13 KT= 2 N(KT)= 2 IKF(KT,N)= 0 IKP(KT,N(KT))=13 NH= 2 IKC(KT,N(KT))= 571 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

CONFLICT STUDENT ID 2170 J= 13 SCHEDULING COURSE 571 TIME KT= 3 CONFLICTING COURSE 538

CONFLICT STUDENT ID 2181 J= 13 SCHEDULING COURSE 571 TIME KT= 3 CONFLICTING COURSE 538

CONFLICT STUDENT ID 2200 J= 13 SCHEDULING COURSE 571 TIME KT= 3 CONFLICTING COURSE 538

STUDENT ID	J	SCHEDULING COURSE	TIME KT	CONFLICTING COURSE
2095	13	571	3	12
2119	13	571	3	12
2150	13	571	3	12
2190	13	571	3	12
2202	13	571	3	12
2289	13	571	3	12
2300	13	571	3	12
2332	13	571	3	12
2356	13	571	3	12

J 13 KT= 3 N(KT)= 4 IKF(KT,N)= 0 IKP(KT,N(KT))=13 NH= 2 IKC(KT,N(KT))= 571 IKG(KT,N(KT))= MSE1 = 1 MSEH(KT,N(KT))= 0

J 14 COURSE 572 NUMBER OF HOURS 3

SCHEDULED AT PERIODS 4 5 6

STUDENT ID	J	SCHEDULING COURSE	TIME KT	CONFLICTING COURSE
2136	14	572	4	551

J 14 KT= 4 N(KT)= 5 IKF(KT,N)= 0 IKP(KT,N(KT))=14 NH= 2 IKC(KT,N(KT))= 572 IKG(KT,N(KT))= MSE1 = 1 MSEH(KT,N(KT))= 0

STUDENT ID	J	SCHEDULING COURSE	TIME KT	CONFLICTING COURSE
2312	14	572	5	550

J 14 KT= 5 N(KT)= 2 IKF(KT,N)= 0 IKP(KT,N(KT))=14 NH= 2 IKC(KT,N(KT))= 572 IKG(KT,N(KT))= MSE1 = 1 MSEH(KT,N(KT))= 0

STUDENT ID	J	SCHEDULING COURSE	TIME KT	CONFLICTING COURSE
2071	14	572	6	559
2091	14	572	6	559
2232	14	572	6	559
2263	14	572	6	559
2333	14	572	6	559
2027	14	572	6	561

J 14 KT= 6 N(KT)= 4 IKF(KT,N)= 0 IKP(KT,N(KT))=14 NH= 2 IKC(KT,N(KT))= 572 IKG(KT,N(KT))= MSE1 = 1 MSEH(KT,N(KT))= 0

J 15 COURSE 573 NUMBER OF HOURS 3

SCHEDULED AT PERIODS 1 2 3

STUDENT ID	J	SCHEDULING COURSE	TIME KT	CONFLICTING COURSE
1029	15	573	1	578

J 15 KT= 1 N(KT)= 5 IKF(KT,N)= 0 IKP(KT,N(KT))=15 NH= 2 IKC(KT,N(KT))= 573 IKG(KT,N(KT))= MSE1 = 1 MSEH(KT,N(KT))= 0

J 15 SCHEDULING COURSE 573 NO CONFLICTS GENERATED

J 15 KT= 2 N(KT)= 3 IKF(KT,N)= 0 IKP(KT,N(KT))=15 NH= 2 IKC(KT,N(KT))= 573 IKG(KT,N(KT))= MSE1 = 1 MSEH(KT,N(KT))= 0

CONFLICT	STUDENT ID	1023	J= 15	SCHEDULING COURSE	573	TIME KT= 3	CONFLICTING COURSE	538
CONFLICT	STUDENT ID	1064	J= 15	SCHEDULING CCOURSE	573	TIME KT= 3	CONFLICTING COURSE	538
CONFLICT	STUDENT ID	1123	J= 15	SCHEDULING COURSE	573	TIME KT= 3	CONFLICTING COURSE	538
CONFLICT	STUDENT ID	1174	J= 15	SCHEDULING COURSE	573	TIME KT= 3	CONFLICTING COURSE	538
CONFLICT	STUDENT ID	1189	J= 15	SCHEDULING CCOURSE	573	TIME KT= 3	CONFLICTING COURSE	538
CONFLICT	STUDENT ID	1265	J= 15	SCHEDULING COURSE	573	TIME KT= 3	CONFLICTING COURSE	538

J 15 KT= 3 N(KT)= 5 IKF(KT,N)= 0 IKP(KT,N(KT))=15 NH= 2 IKC(KT,N(KT))= 573 IKG(KT,N(KT))= MSE1 = 1 MSEH(KT,N(KT))= 0

SCHEDULING
RESTRICTED
SINGLE-SECTION COURSES

MINIMUM COST = 4

CONFLICT	STUDENT ID	1034	J= 16	SCHEDULING COURSE	19	TIME KT= 4	CONFLICTING COURSE	551
CONFLICT	STUDENT ID	2245	J= 16	SCHEDULING COURSE	19	TIME KT= 4	CONFLICTING COURSE	572

J 16 KT= 4 N(KT)= 6 IKF(KT,N)= 0 IKP(KT,N(KT))=16 NH= 0 IKC(KT,N(KT))= 19 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 1

CONFLICT	STUDENT ID	1184	J= 17	SCHEDULING COURSE	21	TIME KT= 5	CONFLICTING COURSE	550
CONFLICT	STUDENT ID	2050	J= 17	SCHEDULING COURSE	21	TIME KT= 5	CONFLICTING COURSE	572

J 17 KT= 5 N(KT)= 3 IKF(KT,N)= 0 IKP(KT,N(KT))=17 NH= 0 IKC(KT,N(KT))= 21 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 1

MINIMUM COST = -20

CONFLICT	STUDENT ID	1034	J= 18	SCHEDULING COURSE	320	TIME KT= 4	CONFLICTING COURSE	551
----------------	------------	------	-------	-------------------	-----	------------	--------------------	-----

J 18 KT= 4 N(KT)= 7 IKF(KT,N)= 0 IKP(KT,N(KT))=18 NH= 0 IKC(KT,N(KT))= 320 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 2

CONFLICT	STUDENT ID	1184	J= 19	SCHEDULING COURSE	321	TIME KT= 5	CONFLICTING COURSE	550
----------------	------------	------	-------	-------------------	-----	------------	--------------------	-----

J 19 KT= 5 N(KT)= 4 IKF(KT,N)= 0 IKP(KT,N(KT))=19 NH= 0 IKC(KT,N(KT))= 321 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 2

MINIMUM COST = 6

CONFLICT	STUDENT ID	1201	J= 20	SCHEDULING COURSE	575	TIME KT= 6	CONFLICTING COURSE	559
CONFLICT	STUDENT ID	1105	J= 20	SCHEDULING COURSE	575	TIME KT= 6	CONFLICTING COURSE	561
CONFLICT	STUDENT ID	2075	J= 20	SCHEDULING COURSE	575	TIME KT= 6	CONFLICTING COURSE	572
CONFLICT	STUDENT ID	2324	J= 20	SCHEDULING COURSE	575	TIME KT= 6	CONFLICTING COURSE	572

J 20 KT= 6 N(KT)= 5 IKF(KT,N)= 0 IKP(KT,N(KT))=20 NH= 0 IKC(KT,N(KT))= 575 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

CONFLICT	STUDENT ID	2234	J= 21	SCHEDULING COURSE	577	TIME KT= 1	CONFLICTING COURSE	571
CONFLICT	STUDENT ID	1239	J= 21	SCHEDULING COURSE	577	TIME KT= 1	CONFLICTING COURSE	573

J 21 KT= 1 N(KT)= 6 IKF(KT,N)= 0 IKP(KT,N(KT))=21 NH= 0 IKC(KT,N(KT))= 577 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

SCHEDULING
SINGLE-SECTION COURSES

CONFLICT STUDENT ID 1333 J= 22 SCHEDULING COURSE 16 TIME KT= 2 CONFLICTING COURSE 573
J 22 KT= 2 N(KT)= 4 IKF(KT,N)= 0 IKP(KT,N(KT))=22 NH= 0 IKC(KT,N(KT))= 16 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 1

CONFLICT STUDENT ID 2242 J= 23 SCHEDULING COURSE 532 TIME KT= 6 CONFLICTING COURSE 572
J 23 KT= 6 N(KT)= 6 IKF(KT,N)= 0 IKP(KT,N(KT))=23 NH= 0 IKC(KT,N(KT))= 532 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

CONFLICT STUDENT ID 1189 J= 24 SCHEDULING COURSE 316 TIME KT= 3 CONFLICTING COURSE 538
CONFLICT STUDENT ID 2081 J= 24 SCHEDULING COURSE 316 TIME KT= 3 CONFLICTING COURSE 571
CONFLICT STUDENT ID 2355 J= 24 SCHEDULING COURSE 316 TIME KT= 3 CONFLICTING COURSE 571
CONFLICT STUDENT ID 1189 J= 24 SCHEDULING COURSE 316 TIME KT= 3 CONFLICTING COURSE 573
CONFLICT STUDENT ID 1245 J= 24 SCHEDULING COURSE 316 TIME KT= 3 CONFLICTING COURSE 573
J 24 KT= 3 N(KT)= 6 IKF(KT,N)= 0 IKP(KT,N(KT))=24 NH= 0 IKC(KT,N(KT))= 316 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 2

CONFLICT STUDENT ID 2271 J= 25 SCHEDULING COURSE 7 TIME KT= 5 CONFLICTING COURSE 550
J 25 KT= 5 N(KT)= 5 IKF(KT,N)= 0 IKP(KT,N(KT))=25 NH= 0 IKC(KT,N(KT))= 7 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 1

CONFLICT STUDENT ID 1274 J= 26 SCHEDULING COURSE 306 TIME KT= 3 CONFLICTING COURSE 570
J 26 KT= 3 N(KT)= 7 IKF(KT,N)= 0 IKP(KT,N(KT))=26 NH= 0 IKC(KT,N(KT))= 306 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 2

CONFLICT STUDENT ID 2349 J= 27 SCHEDULING COURSE 555 TIME KT= 4 CONFLICTING COURSE 551
J 27 KT= 4 N(KT)= 8 IKF(KT,N)= 0 IKP(KT,N(KT))=27 NH= 0 IKC(KT,N(KT))= 555 IKG(KT,N(KT))= MSE1 = 1 MSEM(KT,N(KT))= 0

CONFLICT	STUDENT ID	2100	J = 20	SCHEDULING COURSE	547	TIME KT = 2	CONFLICTING COURSE	570
CONFLICT	STUDENT ID	2190	J = 20	SCHEDULING COURSE	547	TIME KT = 2	CONFLICTING COURSE	570

J 20 KT = 2 N(KT) = 5 IKF(KT,N) = 0 IKP(KT,N(KT)) = 20 NH = 0 IKC(KT,N(KT)) = 547 IKG(KT,N(KT)) = MSE1 = 1 MSEN(KT,N(KT)) = 0

TOTAL NUMBER OF CONFLICTS = 54

	TIME PERIODS					
	1	2	3	4	5	6
	*****	*****	*****	*****	*****	*****
LOCKED-IN COURSES						
	578	0	538	540	550	529
	579	0	12	541	0	559
	0	0	0	551	0	561
SEVERAL-HOUR COURSES						
	570	570	570	570	0	0
	571	571	571	0	0	0
	0	0	0	572	572	572
	573	573	573	0	0	0
RESTRICTED SINGLE-SECTION COURSES						
	0	0	0	19	21	0
	0	0	0	320	321	0
	577	0	0	0	0	575
SINGLE-SECTION COURSES						
	0	16	316	555	7	532
	0	547	306	0	0	0

J= 29 SCHEDULING COURSE 534 NO CONFLICTS GENERATED

J= 29 COURSE= 534 NO. OF STUDENTS = 97 NO. OF SECTIONS = 3 SCHEDULED AT TIME PERIODS 4 5 6

J= 29 KT= 4 N(KT)= 9 IKP(KT,NI)= 29 IKC(KT,NI)= 534 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 29 TIME PERIOD 4 SECTION 1 NO. OF STUDENTS 33
ID OF STUDENTS
1002 1007 1011 1024
1035 1046 1057 1061
1072 1093 1116 1136
1141 1150 1160 1170
1181 1193 1196 1215
1218 1221 1241 1257
1267 1293 1305 1321
1336 1340 1346 1355
2339

J= 29 KT= 5 N(KT)= 6 IKP(KT,NI)= 29 IKC(KT,NI)= 534 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 29 TIME PERIOD 5 SECTION 2 NO. OF STUDENTS 32
ID OF STUDENTS
1005 1008 1019 1026
1041 1049 1058 1062
1089 1108 1125 1138
1144 1152 1161 1171
1182 1195 1211 1217
1226 1231 1249 1266
1278 1301 1312 1331
1337 1343 1350 2042

J= 29 KT= 6 N(KT)= 7 IKP(KT,NI)= 29 IKC(KT,NI)= 534 IKG(KT,NI)= IKF(KT,NI)=SECTION 3

J= 29 TIME PERIOD 6 SECTION 3 NO. OF STUDENTS 32
ID OF STUDENTS
1006 1014 1021 1027
1048 1050 1091 1094
1118 1121 1131 1139
1145 1158 1167 1173
1191 1212 1213 1216
1219 1240 1255 1270
1290 1306 1313 1334
1338 1348 1352 2314

J= 30 COURSE 581 CONFLICT 1 CONFLICTING STUDENT ID 1231 REQUEST LIST 534 551 553 559 568 581

J= 30 COURSE= 581 NO. OF STUDENTS = 93 NO. OF SECTIONS = 3 SCHEDULED AT TIME PERIODS 4 5 6

J= 30 KT= 4 N(KT)= 10 IKP(KT,NI)= 30 IKC(KT,NI)= 581 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 30 TIME PERIOD 4 SECTION 1 NO. OF STUDENTS 31 ID OF STUDENTS
1005 1006 1014 1026
1027 1048 1049 1062
1085 1094 1108 1118
1125 1138 1145 1161
1167 1191 1211 1212
1216 1219 1249 1290
1301 1306 1313 1338
1343 1352 2314

J= 30 KT= 5 N(KT)= 7 IKP(KT,NI)= 30 IKC(KT,NI)= 581 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 30 TIME PERIOD 5 SECTION 2 NO. OF STUDENTS 31 ID OF STUDENTS
1002 1011 1021 1024
1046 1050 1061 1072
1091 1093 1116 1121
1131 1139 1158 1160
1173 1181 1196 1213
1240 1241 1255 1267
1270 1305 1334 1336
1346 1355 2339

J= 30 KT= 6 N(KT)= 8 IKP(KT,NI)= 30 IKC(KT,NI)= 581 IKG(KT,NI)= IKF(KT,NI)=SECTION 3

J= 30 TIME PERIOD 6 SECTION 3 NO. OF STUDENTS 30 ID OF STUDENTS
1027 1038 1019 1035
1041 1057 1058 1086
1136 1141 1144 1150
1152 1171 1182 1193
1195 1215 1217 1221
1226 1257 1266 1293
1312 1321 1331 1340
1350 2042

J= 31	COURSE 568	CONFLICT 1	CONFLICTING STUDENT ID	1002	REQUEST LIST 534	535	557	559	568	581
J= 31	COURSE 568	CONFLICT 2	CONFLICTING STUDENT ID	1005	REQUEST LIST 534	557	559	568	570	581
J= 31	COURSE 568	CONFLICT 3	CONFLICTING STUDENT ID	1006	REQUEST LIST 24	321	531	534	552	568 581
J= 31	COURSE 568	CONFLICT 4	CONFLICTING STUDENT ID	1011	REQUEST LIST 534	535	557	559	568	581
J= 31	COURSE 568	CONFLICT 5	CONFLICTING STUDENT ID	1041	REQUEST LIST 5	306	534	551	557	568 581
J= 31	COURSE 568	CONFLICT 6	CONFLICTING STUDENT ID	1046	REQUEST LIST 531	534	553	561	568	581
J= 31	COURSE 568	CONFLICT 7	CONFLICTING STUDENT ID	1049	REQUEST LIST 531	534	557	559	568	581
J= 31	COURSE 568	CONFLICT 8	CONFLICTING STUDENT ID	1061	REQUEST LIST 534	548	553	559	568	581
J= 31	COURSE 568	CONFLICT 9	CONFLICTING STUDENT ID	1062	REQUEST LIST 530	534	557	559	568	581
J= 31	COURSE 568	CONFLICT 10	CONFLICTING STUDENT ID	1072	REQUEST LIST 530	534	557	559	568	581
J= 31	COURSE 568	CONFLICT 11	CONFLICTING STUDENT ID	1089	REQUEST LIST 531	534	553	561	568	581
J= 31	COURSE 568	CONFLICT 12	CONFLICTING STUDENT ID	1093	REQUEST LIST 5	306	534	552	559	568 581
J= 31	COURSE 568	CONFLICT 13	CONFLICTING STUDENT ID	1108	REQUEST LIST 534	557	564	568	575	581
J= 31	COURSE 568	CONFLICT 14	CONFLICTING STUDENT ID	1116	REQUEST LIST 534	548	557	559	568	581
J= 31	COURSE 568	CONFLICT 15	CONFLICTING STUDENT ID	1125	REQUEST LIST 16	308	534	553	559	568 581
J= 31	COURSE 568	CONFLICT 16	CONFLICTING STUDENT ID	1138	REQUEST LIST 534	548	557	559	568	581
J= 31	COURSE 568	CONFLICT 17	CONFLICTING STUDENT ID	1160	REQUEST LIST 5	324	534	557	559	568 581
J= 31	COURSE 568	CONFLICT 18	CONFLICTING STUDENT ID	1161	REQUEST LIST 531	534	552	561	568	581
J= 31	COURSE 568	CONFLICT 19	CONFLICTING STUDENT ID	1181	REQUEST LIST 534	535	557	559	568	581
J= 31	COURSE 568	CONFLICT 20	CONFLICTING STUDENT ID	1211	REQUEST LIST 534	540	557	559	568	581
J= 31	COURSE 568	CONFLICT 21	CONFLICTING STUDENT ID	1216	REQUEST LIST 534	550	557	564	568	581
J= 31	COURSE 568	CONFLICT 22	CONFLICTING STUDENT ID	1219	REQUEST LIST 21	531	534	557	568	581
J= 31	COURSE 568	CONFLICT 23	CONFLICTING STUDENT ID	1221	REQUEST LIST 534	550	557	560	568	581
J= 31	COURSE 568	CONFLICT 24	CONFLICTING STUDENT ID	1231	REQUEST LIST 534	551	553	559	568	581
J= 31	COURSE 568	CONFLICT 25	CONFLICTING STUDENT ID	1241	REQUEST LIST 534	548	557	559	568	581
J= 31	COURSE 568	CONFLICT 26	CONFLICTING STUDENT ID	1249	REQUEST LIST 534	546	553	561	568	581
J= 31	COURSE 568	CONFLICT 27	CONFLICTING STUDENT ID	1255	REQUEST LIST 5	323	534	551	557	568 581
J= 31	COURSE 568	CONFLICT 28	CONFLICTING STUDENT ID	1267	REQUEST LIST 534	535	557	559	568	581
J= 31	COURSE 568	CONFLICT 29	CONFLICTING STUDENT ID	1270	REQUEST LIST 531	534	551	557	568	581
J= 31	COURSE 568	CONFLICT 30	CONFLICTING STUDENT ID	1305	REQUEST LIST 534	548	557	559	568	581
J= 31	COURSE 568	CONFLICT 31	CONFLICTING STUDENT ID	1346	REQUEST LIST 534	557	564	568	575	581

J= 31 COURSE= 568 NO. OF STUDENTS = 95 NO. OF SECTIONS = 3 SCHEDULED AT TIME PERIODS 4 5 6

J= 31 KT= 4 N(KT)= 11 IKP(KT,NI)= 31 IKC(KT,NI)= 568 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 31 TIME PERIOD 4 SECTION 1 NO. OF STUDENTS 28 ID OF STUDENTS

1008	1019	1021	1050
1056	1086	1091	1121
1131	1139	1144	1152
1158	1171	1173	1182
1195	1213	1217	1226
1240	1266	1312	1331
1334	1348	1350	2042

J= 31 KT= 5 N(KT)= 8 IKP(KT,NI)= 31 IKC(KT,NI)= 568 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 31 TIME PERIOD 5 SECTION 2 NO. OF STUDENTS 27 ID OF STUDENTS

1007	1014	1027	1035
1048	1057	1094	1118
1136	1141	1145	1150
1167	1191	1193	1212
1215	1257	1290	1293
1306	1313	1321	1338
1340	1352	2314	

J= 31 KT= 6 N(KT)= 9 IKP(KT,NI)= 31 IKC(KT,NI)= 568 IKG(KT,NI)= IKF(KT,NI)=SECTION 3

J= 31 TIME PERIOD 6 SECTION 3 NO. OF STUDENTS 9 ID OF STUDENTS

1024	1026	1278	1301
1336	1343	1355	2149
2339			

J= 32 COURSE 530 CONFLICT 1 CONFLICTING STUDENT ID 1015 REQUEST LIST 530 546 570

J= 32 COURSE= 530 NO. OF STUDENTS = 65 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 1 4

J= 32 KT= 1 N(KT)= 7 IKP(KT,NI)= 32 IKC(KT,NI)= 530 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 32 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 32
ID OF STUDENTS
1008 1021 1036 1057
1058 1062 1072 1083
1090 1121 1141 1156
1177 1183 1187 1195
1200 1233 1234 1260
1261 1276 1301 1331
1334 2011 2166 2183
2208 2223 2293 2314

J= 32 KT= 4 N(KT)= 12 IKP(KT,NI)= 32 IKC(KT,NI)= 530 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 32 TIME PERIOD 4 SECTION 2 NO. OF STUDENTS 32
ID OF STUDENTS
1029 1031 1043 1065
1069 1074 1076 1087
1102 1142 1149 1168
1179 1188 1198 1205
1222 1236 1247 1259
1295 1311 1323 1339
1349 2024 2168 2202
2215 2260 2352 2354

J= 33 COURSE 5 CONFLICT 1 CONFLICTING STUDENT ID 1210 REQUEST LIST 5 307 570 575

J= 33 COURSE= 5 NO. OF STUDENTS = 49 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 3 6

J= 33 KT= 3 N(KT)= 8 IKP(KT,NI)= 33 IKC(KT,NI)= 5 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 33 TIME PERIOD 3 SECTION 1 NO. OF STUDENTS 28
ID OF STUDENTS
1039 1041 1093 1094
1096 1135 1160 1175
1226 1255 1257 1282
1290 1306 1312 1313
1324 1335 1340 1352
2022 2032 2054 2066
2175 2176 2310 2328

J= 33 KT= 6 N(KT)= 10 IKP(KT,NI)= 33 IKC(KT,NI)= 5 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 33 TIME PERIOD 6 SECTION 2 NO. OF STUDENTS 20
ID OF STUDENTS
1090 1100 1107 1140
1156 1197 1243 1274
1275 1286 1288 1330
1342 1342 1345 2121
2137 2161 2209 2351

J= 34 SCHEDULING COURSE 562 NO CONFLICTS GENERATED

J= 34 COURSE= 562 NO. OF STUDENTS = 7 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 2 4
J= 34 KT= 2 N(KT)= 6 IKP(KT,NI)= 34 IKC(KT,NI)= 562 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 34 TIME PERIOD 2 SECTION 1 NO. OF STUDENTS 5 ID OF STUDENTS
1050 1242 1293 2054
2147

J= 34 KT= 4 N(KT)= 13 IKP(KT,NI)= 34 IKC(KT,NI)= 562 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 34 TIME PERIOD 4 SECTION 2 NO. OF STUDENTS 2 ID OF STUDENTS
1164 2311

J= 35 COURSE 564 CONFLICT 1 CONFLICTING STUDENT ID 1195 REQUEST LIST 530 534 557 564 568 581
J= 35 COURSE 564 CONFLICT 2 CONFLICTING STUDENT ID 2208 REQUEST LIST 530 564 565 572

J= 35 COURSE= 564 NO. OF STUDENTS = 47 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 1 5
J= 35 KT= 1 N(KT)= 8 IKP(KT,NI)= 35 IKC(KT,NI)= 564 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 35 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 22 ID OF STUDENTS
1012 1039 1042 1091
1094 1100 1126 1216
1217 1270 1332 1346
2032 2107 2120 2157
2201 2205 2239 2281
2291 2305

J= 35 KT= 5 N(KT)= 9 IKP(KT,NI)= 35 IKC(KT,NI)= 564 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 35 TIME PERIOD 5 SECTION 2 NO. OF STUDENTS 23 ID OF STUDENTS
1031 1097 1111 1124
1194 1230 1238 1246
1252 1264 1283 1351
2086 2126 2167 2179
2204 2234 2257 2262
2280 2332 2358

J= 36 COURSE 546 CONFLICT 1 CONFLICTING STUDENT ID 1110 REQUEST LIST 546 558 573 575
J= 36 COURSE 546 CONFLICT 2 CONFLICTING STUDENT ID 1312 REQUEST LIST 5 324 534 546 557 568 581

J= 36 COURSE= 546 NO. OF STUDENTS = 43 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 3 6
J= 36 KT= 3 N(KT)= 9 IKP(KT,NI)= 36 IKC(KT,NI)= 546 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 36 TIME PERIOD 3 SECTION 1 NO. OF STUDENTS 22 ID OF STUDENTS
1001 1027 1035 1038
1080 1112 1130 1131
1136 1145 1147 1167
1173 1182 1193 1199
1213 1249 1278 1288
1301 1331

J= 36 KT= 6 N(KT)= 11 IKP(KT,NI)= 36 IKC(KT,NI)= 546 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 36 TIME PERIOD 6 SECTION 2 NO. OF STUDENTS 19 ID OF STUDENTS
1015 1033 1036 1045
1083 1096 1106 1111
1155 1176 1178 1194
1218 1222 1235 1298
1300 2135 2284

J= 37 COURSE 324 CONFLICT 1 CONFLICTING STUDENT ID 1091 REQUEST LIST 8 324 534 557 564 568 581
 J= 37 COURSE 324 CONFLICT 2 CONFLICTING STUDENT ID 1094 REQUEST LIST 5 324 534 557 564 568 581

J= 37 COURSE= 324 NO. OF STUDENTS = 32 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 1 4
 J= 37 KT= 1 N(KT)= 9 IKP(KT,NI)= 37 IKC(KT,NI)= 324 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 37 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 15 ID OF STUDENTS
 1014 1075 1160 1169
 1240 1257 1279 1288
 1312 1313 1343 2054
 2151 2216 2240

J= 37 KT= 4 N(KT)= 14 IKP(KT,NI)= 37 IKC(KT,NI)= 324 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 37 TIME PERIOD 4 SECTION 2 NO. OF STUDENTS 15 ID OF STUDENTS
 1030 1112 1148 1156
 1174 1206 1282 1286
 1330 2058 2145 2167
 2174 2211 2266

J= 38 COURSE 24 CONFLICT 1 CONFLICTING STUDENT ID 1246 REQUEST LIST 24 308 558 564 573

J= 38 COURSE= 24 NO. OF STUDENTS = 36 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 1 5
 J= 38 KT= 1 N(KT)= 10 IKP(KT,NI)= 38 IKC(KT,NI)= 24 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 38 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 18 ID OF STUDENTS
 1040 1045 1142 1166
 1272 1316 2015 2017
 2063 2077 2133 2147
 2148 2179 2217 2252
 2282 2334

J= 38 KT= 5 N(KT)= 10 IKP(KT,NI)= 38 IKC(KT,NI)= 24 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 38 TIME PERIOD 5 SECTION 2 NO. OF STUDENTS 17 ID OF STUDENTS
 1006 1077 1162 1189
 1304 1319 1320 1328
 2040 2093 2125 2128
 2140 2194 2219 2261
 2315

J= 39 SCHEDULING COURSE 305 NO CONFLICTS GENERATED

J= 39 COURSE= 305 NO. OF STUDENTS = 20 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 2 5
 J= 39 KT= 2 N(KT)= 7 IKP(KT,NI)= 39 IKC(KT,NI)= 305 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 39 TIME PERIOD 2 SECTION 1 NO. OF STUDENTS 9 ID OF STUDENTS
 1045 1277 2010 2034
 2043 2082 2159 2253

J= 39 KT= 5 N(KT)= 11 IKP(KT,NI)= 39 IKC(KT,NI)= 305 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 39 TIME PERIOD 5 SECTION 2 NO. OF STUDENTS 12 ID OF STUDENTS
 1224 2031 2039 2045
 2067 2104 2106 2190
 2154 2214 2296 2303

J= 40 COURSE 8 CONFLICT 1 CONFLICTING STUDENT ID 1172 REQUEST LIST 8 540 557 573
 J= 40 COURSE 8 CONFLICT 2 CONFLICTING STUDENT ID 2269 REQUEST LIST 8 309 550 555 571
 J= 40 COURSE 8 CONFLICT 3 CONFLICTING STUDENT ID 2311 REQUEST LIST 8 309 556 562 571

J= 40 COURSE= 8 NO. OF STUDENTS = 49 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 3 4
 J= 40 KT= 3 N(KT)= 10 IKP(KT,NI)= 40 IKC(KT,NI)= 8 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 40 TIME PERIOD 3 SECTION 1 NO. OF STUDENTS 23 ID OF STUDENTS
 1014 1067 1075 1091
 1121 1191 1228 1277
 1284 1343 2010 2062
 2087 2117 2151 2182
 2194 2232 2236 2254
 2263 2281 2288

J= 40 KT= 4 N(KT)= 15 IKP(KT,NI)= 40 IKC(KT,NI)= 8 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 40 TIME PERIOD 4 SECTION 2 NO. OF STUDENTS 23 ID OF STUDENTS
 1030 1056 1146 1169
 1206 1224 1245 1344
 2002 2007 2049 2065
 2069 2144 2161 2181
 2203 2237 2247 2266
 2276 2284 2292

J= 41 COURSE 308 CONFLICT 1 CONFLICTING STUDENT ID 1333 REQUEST LIST 16 19 308 320 558 573
 J= 41 COURSE 308 CONFLICT 2 CONFLICTING STUDENT ID 2034 REQUEST LIST 9 305 308 323 572
 J= 41 COURSE 308 CONFLICT 3 CONFLICTING STUDENT ID 2277 REQUEST LIST 9 308 543 555 571

J= 41 COURSE= 308 NO. OF STUDENTS = 57 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 2 4
 J= 41 KT= 2 N(KT)= 8 IKP(KT,NI)= 41 IKC(KT,NI)= 308 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 41 TIME PERIOD 2 SECTION 1 NO. OF STUDENTS 27 ID OF STUDENTS
 1040 1079 1104 1125
 1132 1152 1292 1316
 1334 1342 1342 1345
 2038 2048 2077 2115
 2149 2157 2165 2169
 2199 2210 2241 2242
 2251 2299 2309

J= 41 KT= 4 N(KT)= 16 IKP(KT,NI)= 41 IKC(KT,NI)= 308 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 41 TIME PERIOD 4 SECTION 2 NO. OF STUDENTS 27 ID OF STUDENTS
 1077 1100 1103 1129
 1146 1197 1246 1254
 1303 1304 1319 2016
 2024 2072 2107 2112
 2131 2178 2192 2196
 2207 2212 2243 2246
 2268 2274 2338

J= 42 COURSE 323 CONFLICT 1 CONFLICTING STUDENT ID 1320 REQUEST LIST 24 323 540 550 573

J= 42 COURSE= 323 NO. OF STUDENTS = 35 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 3 4
J= 42 KT= 3 N(KT)= 11 IKP(KT,NI)= 42 IKC(KT,NI)= 323 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 42 TIME PERIOD 3 SECTION 1 NO. OF STUDENTS 10 ID OF STUDENTS
1090 1107 1121 1139
1226 1228 1255 1306
2034 2141 2185 2223
2279 2282 2318 2322
2340 2348

J= 42 KT= 4 N(KT)= 17 IKP(KT,NI)= 42 IKC(KT,NI)= 323 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 42 TIME PERIOD 4 SECTION 2 NO. OF STUDENTS 16 ID OF STUDENTS
1054 1146 1244 1344
2007 2019 2025 2057
2137 2161 2193 2231
2247 2283 2284 2359

J= 43 SCHEDULING COURSE 23 NO CONFLICTS GENERATED

J= 43 COURSE= 23 NO. OF STUDENTS = 24 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 2 4
J= 43 KT= 2 N(KT)= 9 IKP(KT,NI)= 43 IKC(KT,NI)= 23 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 43 TIME PERIOD 2 SECTION 1 NO. OF STUDENTS 12 ID OF STUDENTS
1024 1104 1132 1152
1257 1292 1334 2093
2136 2267 2306 2329

J= 43 KT= 4 N(KT)= 18 IKP(KT,NI)= 43 IKC(KT,NI)= 23 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 43 TIME PERIOD 4 SECTION 2 NO. OF STUDENTS 12 ID OF STUDENTS
1073 1254 2003 2028
2033 2035 2046 2110
2138 2300 2325 2335

J= 44 COURSE 309 CONFLICT 1 CONFLICTING STUDENT ID 2269 REQUEST LIST 8 309 550 555 571
 J= 44 COURSE 309 CONFLICT 2 CONFLICTING STUDENT ID 2311 REQUEST LIST 8 309 556 562 571

J= 44 COURSE= 309 NO. OF STUDENTS =133 NO. OF SECTIONS = 4 SCHEDULED AT TIME PERIODS 1 2 3 4
 J= 44 KT= 1 N(KT)= 11 IKP(KT,NI)= 44 IKC(KT,NI)= 309 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 44 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 30

ID OF STUDENTS
 2001 2015 2017 2037
 2050 2060 2063 2077
 2097 2105 2121 2133
 2148 2161 2179 2209
 2217 2229 2235 2249
 2252 2265 2271 2307
 2312 2323 2326 2330
 2334 2360

J= 44 KT= 2 N(KT)= 10 IKP(KT,NI)= 44 IKC(KT,NI)= 309 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 44 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 29

ID OF STUDENTS
 2018 2021 2027 2040
 2052 2064 2074 2093
 2101 2120 2122 2134
 2136 2151 2158 2219
 2226 2227 2239 2245
 2256 2260 2267 2293
 2306 2313 2327 2329
 2351

J= 44 KT= 3 N(KT)= 12 IKP(KT,NI)= 44 IKC(KT,NI)= 309 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
 J= 44 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 30

ID OF STUDENTS
 2006 2009 2022 2032
 2059 2062 2090 2092
 2094 2117 2127 2164
 2171 2175 2176 2182
 2194 2220 2232 2236
 2261 2263 2281 2288
 2291 2310 2316 2328
 2345 2353

J= 44 KT= 4 N(KT)= 19 IKP(KT,NI)= 44 IKC(KT,NI)= 309 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
 J= 44 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 42

ID OF STUDENTS
 2002 2003 2008 2014
 2028 2029 2033 2035
 2041 2046 2049 2065
 2069 2076 2095 2099
 2110 2119 2124 2125
 2138 2140 2144 2150
 2181 2198 2203 2225
 2237 2258 2262 2276
 2289 2292 2300 2304
 2308 2315 2325 2332
 2356 2361

J= 45 COURSE 9 CONFLICT 1 CONFLICTING STUDENT ID 2076 REQUEST LIST 9 325 570 575
 J= 45 COURSE 9 CONFLICT 2 CONFLICTING STUDENT ID 2100 REQUEST LIST 9 325 547 570
 J= 45 COURSE 9 CONFLICT 3 CONFLICTING STUDENT ID 2104 REQUEST LIST 9 305 570 505
 J= 45 COURSE 9 CONFLICT 4 CONFLICTING STUDENT ID 2106 REQUEST LIST 9 305 570 505
 J= 45 COURSE 9 CONFLICT 5 CONFLICTING STUDENT ID 2150 REQUEST LIST 9 305 547 570
 J= 45 COURSE 9 CONFLICT 6 CONFLICTING STUDENT ID 2277 REQUEST LIST 9 308 543 555 571
 J= 45 COURSE 9 CONFLICT 7 CONFLICTING STUDENT ID 2296 REQUEST LIST 9 305 549 570
 J= 45 COURSE 9 CONFLICT 8 CONFLICTING STUDENT ID 2247 REQUEST LIST 9 399 555 559 571
 J= 45 COURSE 9 CONFLICT 9 CONFLICTING STUDENT ID 2302 REQUEST LIST 9 305 548 570
 J= 45 COURSE 9 CONFLICT 10 CONFLICTING STUDENT ID 2303 REQUEST LIST 9 305 548 570

J= 45 COURSE= 9 NO. OF STUDENTS = 98 NO. OF SECTIONS = 4 SCHEDULED AT TIME PERIODS 1 2 3 4
 J= 45 KT= 1 N(KT)= 12 IKP(KT,NI)= 45 IKC(KT,NI)= 9 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 45 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 16 ID OF STUDENTS
 2013 2039 2044 2067
 2083 2111 2135 2168
 2104 2206 2215 2216
 2240 2275 2290 2301

J= 45 KT= 2 N(KT)= 11 IKP(KT,NI)= 45 IKC(KT,NI)= 9 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 45 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 16 ID OF STUDENTS
 2020 2034 2038 2043
 2048 2002 2115 2149
 2159 2195 2210 2241
 2242 2251 2309 2357

J= 45 KT= 3 N(KT)= 13 IKP(KT,NI)= 45 IKC(KT,NI)= 9 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
 J= 45 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 22 ID OF STUDENTS
 2026 2036 2045 2079
 2085 2089 2141 2162
 2188 2214 2222 2223
 2230 2279 2202 2286
 2318 2322 2340 2346
 2348 2349

J= 45 KT= 4 N(KT)= 20 IKP(KT,NI)= 45 IKC(KT,NI)= 9 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
 J= 45 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 34 ID OF STUDENTS
 2016 2019 2024 2025
 2031 2050 2072 2081
 2080 2090 2107 2112
 2131 2145 2167 2174
 2178 2192 2193 2196
 2207 2211 2212 2220
 2231 2243 2246 2268
 2274 2203 2330 2355
 2358 2359

J= 46 COURSE 553 CONFLICT 1 CONFLICTING STUDENT ID 1077 REQUEST LIST 24 308 531 553 573
 J= 46 COURSE 553 CONFLICT 2 CONFLICTING STUDENT ID 1148 REQUEST LIST 5 324 531 553 573
 J= 46 COURSE 553 CONFLICT 3 CONFLICTING STUDENT ID 1309 REQUEST LIST 24 308 553 559 573

J= 46 COURSE= 553 NO. OF STUDENTS =132 NO. OF SECTIONS = 4 SCHEDULED AT TIME PERIODS 1 2 3 4

J= 46 KT= 1 N(KT)= 13 IKP(KT,NI)= 46 IKC(KT,NI)= 553 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 46 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 25 ID OF STUDENTS
 1009 1020 1038 1054
 1065 1085 1095 1103
 1125 1136 1139 1159
 1179 1205 1218 1243
 1245 1268 1280 1284
 1290 1306 1321 1342
 1355

J= 46 KT= 2 N(KT)= 12 IKP(KT,NI)= 46 IKC(KT,NI)= 553 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 46 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 25 ID OF STUDENTS
 1010 1014 1046 1061
 1073 1087 1097 1105
 1127 1137 1155 1156
 1167 1207 1224 1234
 1261 1269 1276 1287
 1294 1313 1323 1343
 1351

J= 46 KT= 3 N(KT)= 14 IKP(KT,NI)= 46 IKC(KT,NI)= 553 IKG(KT,NI)= IKF(KT,NI)=SECTION 3

J= 46 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 25 ID OF STUDENTS
 1016 1018 1047 1066
 1082 1089 1101 1115
 1129 1140 1152 1166
 1189 1215 1231 1252
 1262 1275 1281 1209
 1302 1316 1339 1342
 2080

J= 46 KT= 4 N(KT)= 21 IKP(KT,NI)= 46 IKC(KT,NI)= 553 IKG(KT,NI)= IKF(KT,NI)=SECTION 4

J= 46 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 24 ID OF STUDENTS
 1003 1028 1033 1083
 1088 1090 1104 1126
 1132 1147 1180 1230
 1232 1235 1237 1260
 1265 1277 1283 1288
 1307 1317 1341 1345

J= 47 COURSE 549 CONFLICT 1 CONFLICTING STUDENT ID 2141 REQUEST LIST 9 323 549 566 572

J= 47 COURSE= 549 NO. OF STUDENTS = 32 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 3 6

J= 47 KT= 3 N(KT)= 15 IKP(KT,NI)= 47 IKC(KT,NI)= 549 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 47 TIME PERIOD 3 SECTION 1 NO. OF STUDENTS 16 ID OF STUDENTS
2003 2035 2044 2064
2116 2121 2132 2149
2160 2186 2197 2206
2272 2298 2317 2334

J= 47 KT= 6 N(KT)= 12 IKP(KT,NI)= 47 IKC(KT,NI)= 549 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 47 TIME PERIOD 6 SECTION 2 NO. OF STUDENTS 15 ID OF STUDENTS
2011 2024 2005 2101
2130 2191 2225 2243
2250 2259 2267 2275
2296 2300 2322

J= 48 COURSE 548 CONFLICT 1 CONFLICTING STUDENT ID 1032 REQUEST LIST 548 550 558 573

J= 48 COURSE 548 CONFLICT 2 CONFLICTING STUDENT ID 1313 REQUEST LIST 5 324 534 548 553 568 581

J= 48 COURSE 548 CONFLICT 3 CONFLICTING STUDENT ID 1343 REQUEST LIST 8 324 534 548 553 568 581

J= 48 COURSE 548 CONFLICT 4 CONFLICTING STUDENT ID 2303 REQUEST LIST 9 305 548 570

J= 48 COURSE= 548 NO. OF STUDENTS = 33 NO. OF SECTIONS = 3 SCHEDULED AT TIME PERIODS 1 2 6

J= 48 KT= 1 N(KT)= 14 IKP(KT,NI)= 48 IKC(KT,NI)= 548 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 48 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 27 ID OF STUDENTS
1010 1019 1028 1061
1078 1113 1118 1128
1138 1144 1152 1162
1204 1214 1215 1226
1241 1273 1285 1294
1305 1329 1338 2020
2091 2120 2348

J= 48 KT= 2 N(KT)= 13 IKP(KT,NI)= 48 IKC(KT,NI)= 548 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 48 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 26 ID OF STUDENTS
1008 1026 1051 1057
1084 1116 1120 1129
1140 1150 1168 1171
1211 1212 1223 1229
1247 1284 1290 1299
1307 1321 1350 2087
2283 2344

J= 48 KT= 5 N(KT)= 12 IKP(KT,NI)= 48 IKC(KT,NI)= 548 IKG(KT,NI)= IKF(KT,NI)=SECTION 3

J= 48 TIME PERIOD 5 SECTION 3 NO. OF STUDENTS 26 ID OF STUDENTS
1003 1004 1060 1074
1101 1117 1122 1132
1155 1163 1170 1179
1207 1220 1228 1243
1248 1263 1292 1302
1326 1344 1354 2153
2294 2302

J= 49 COURSE 560 CONFLICT 1 CONFLICTING STUDENT ID 2315 REQUEST LIST 24 309 550 560 571

J= 49 COURSE= 560 NO. OF STUDENTS = 45 NO. OF SECTIONS = 2 SCHEDULED AT TIME PERIODS 1 2

J= 49 KT= 1 N(KT)= 15 IKP(KT,NI)= 49 IKC(KT,NI)= 560 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 49 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 22 ID OF STUDENTS
1017 1107 1170 1191
1221 1271 1317 1323
1330 1350 1352 2012
2016 2023 2033 2066
2160 2186 2219 2241
2329 2339

J= 49 KT= 2 N(KT)= 14 IKP(KT,NI)= 49 IKC(KT,NI)= 560 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 49 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 22 ID OF STUDENTS
1028 1086 1120 1136
1144 1159 1214 1215
1251 1285 1306 1327
1329 2026 2042 2127
2184 2197 2254 2324
2330 2337

J= 50 COURSE 557 CONFLICT 1 CONFLICTING STUDENT ID 1008 REQUEST LIST 530 534 548 557 568 581
 J= 50 COURSE 557 CONFLICT 2 CONFLICTING STUDENT ID 1024 REQUEST LIST 23 399 534 557 568 578 581
 J= 50 COURSE 557 CONFLICT 3 CONFLICTING STUDENT ID 1144 REQUEST LIST 534 548 557 568 581
 J= 50 COURSE 557 CONFLICT 4 CONFLICTING STUDENT ID 1158 REQUEST LIST 534 548 557 568 578 581
 J= 50 COURSE 557 CONFLICT 5 CONFLICTING STUDENT ID 1171 REQUEST LIST 534 548 557 568 578 581
 J= 50 COURSE 557 CONFLICT 6 CONFLICTING STUDENT ID 1212 REQUEST LIST 534 548 557 568 578 581
 J= 50 COURSE 557 CONFLICT 7 CONFLICTING STUDENT ID 1257 REQUEST LIST 5 23 306 324 534 557 568 581
 J= 50 COURSE 557 CONFLICT 8 CONFLICTING STUDENT ID 1293 REQUEST LIST 534 557 562 568 577 581
 J= 50 COURSE 557 CONFLICT 9 CONFLICTING STUDENT ID 1334 REQUEST LIST 23 308 530 534 557 568 581
 J= 50 COURSE 557 CONFLICT 10 CONFLICTING STUDENT ID 1350 REQUEST LIST 534 548 557 568 581

J= 50 COURSE= 557 NO. OF STUDENTS =119 NO. OF SECTIONS = 4 SCHEDULED AT TIME PERIODS 1 2 5 6
 J= 50 KT= 1 N(KT)= 16 IKP(KT,NI)= 50 IKC(KT,NI)= 557 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 50 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 29

ID OF STUDENTS
 1002 1007 1026 1035
 1041 1048 1050 1116
 1131 1145 1150 1173
 1181 1193 1203 1211
 1213 1219 1255 1266
 1267 1270 1336 1337
 1340 2042 2185 2204

J= 50 KT= 2 N(KT)= 15 IKP(KT,NI)= 50 IKC(KT,NI)= 557 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 50 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 38

ID OF STUDENTS
 1005 1011 1019 1021
 1027 1049 1058 1062
 1072 1091 1094 1108
 1110 1121 1138 1141
 1160 1167 1182 1191
 1195 1196 1216 1217
 1221 1226 1240 1241
 1278 1301 1305 1312
 1331 1338 1346 1352
 2314 2339

J= 50 KT= 5 N(KT)= 13 IKP(KT,NI)= 50 IKC(KT,NI)= 557 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
 J= 50 TIME PERIOD 5 SECTION 3 NO. OF STUDENTS 26

ID OF STUDENTS
 1036 1045 1055 1081
 1086 1096 1120 1149
 1172 1183 1190 1202
 1208 1259 1271 1273
 1282 1300 1303 1330
 1347 1348 2012 2352
 2359 2361

J= 50 KT= 6 N(KT)= 13 IKP(KT,NI)= 50 IKC(KT,NI)= 557 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
 J= 50 TIME PERIOD 6 SECTION 4 NO. OF STUDENTS 17

ID OF STUDENTS
 1068 1074 1080 1098
 1124 1128 1135 1170
 1177 1184 1204 1214
 1233 1238 1248 1285
 1327

J= 51 COURSE 531 CONFLICT 1 CONFLICTING STUDENT ID 1077 REQUEST LIST 24 308 531 553 573
 J= 51 COURSE 531 CONFLICT 2 CONFLICTING STUDENT ID 1304 REQUEST LIST 24 308 531 552 573

J= 51 COURSE= 531 NO. OF STUDENTS =161 NO. OF SECTIONS = 4 SCHEDULED AT TIME PERIODS 2 3 4 5
 J= 51 KT= 2 N(KT)= 16 IKP(KT,NI)= 51 IKC(KT,NI)= 531 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 51 TIME PERIOD 2 SECTION 1 NO. OF STUDENTS 40 ID OF STUDENTS

1006	1016	1030	1035
1038	1047	1066	1080
1082	1089	1102	1114
1130	1131	1139	1145
1173	1213	1219	1230
1252	1266	1279	1300
1332	1336	1349	1355
2032	2036	2086	2137
2162	2172	2205	2214
2236	2285	2323	2353

J= 51 KT= 3 N(KT)= 16 IKP(KT,NI)= 51 IKC(KT,NI)= 531 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 51 TIME PERIOD 3 SECTION 2 NO. OF STUDENTS 40 ID OF STUDENTS

1007	1017	1019	1021
1040	1046	1049	1050
1058	1098	1103	1110
1137	1161	1183	1208
1217	1218	1224	1232
1268	1270	1287	1309
1341	1342	1342	2002
2018	2075	2096	2157
2159	2173	2199	2216
2253	2307	2347	2351

J= 51 KT= 4 N(KT)= 22 IKP(KT,NI)= 51 IKC(KT,NI)= 531 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
 J= 51 TIME PERIOD 4 SECTION 3 NO. OF STUDENTS 39 ID OF STUDENTS

1009	1020	1023	1037
1044	1055	1063	1081
1092	1099	1106	1122
1140	1151	1154	1165
1190	1225	1238	1250
1253	1273	1298	1308
1316	1320	1353	2013
2051	2079	2133	2152
2165	2187	2209	2213
2259	2287	2336	

J= 51 KT= 5 N(KT)= 14 IKP(KT,NI)= 51 IKC(KT,NI)= 531 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
 J= 51 TIME PERIOD 5 SECTION 4 NO. OF STUDENTS 40 ID OF STUDENTS

1012	1022	1025	1033
1052	1056	1075	1076
1085	1100	1109	1127
1148	1153	1166	1185
1197	1227	1237	1256
1265	1275	1277	1314
1322	1325	1345	2047
2055	2081	2112	2124
2130	2189	2191	2218
2265	2321	2345	2354

J= 52 COURSE 558 CONFLICT 1 CONFLICTING STUDENT ID 2099 REQUEST LIST 21 309 558 559 571
 J= 52 COURSE 558 CONFLICT 2 CONFLICTING STUDENT ID 2140 REQUEST LIST 24 309 532 558 571

J= 52 COURSE= 558 NO. OF STUDENTS =145 NO. OF SECTIONS = 5 SCHEDULED AT TIME PERIODS 1 3 4 5 6
 J= 52 KT= 1 N(KT)= 17 IKP(KT,NI)= 52 IKC(KT,NI)= 558 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 52 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 23 ID OF STUDENTS
 1092 2002 2006 2009
 2025 2041 2045 2059
 2072 2076 2079 2094
 2122 2155 2159 2175
 2176 2212 2220 2230
 2251 2255 2260 2280
 2316 2318 2345 2351

J= 52 KT= 3 N(KT)= 17 IKP(KT,NI)= 52 IKC(KT,NI)= 558 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 52 TIME PERIOD 3 SECTION 2 NO. OF STUDENTS 27 ID OF STUDENTS
 1057 1332 2007 2015
 2033 2048 2055 2063
 2067 2101 2110 2115
 2148 2163 2165 2178
 2207 2217 2224 2239
 2252 2259 2304 2306
 2321 2331 2360

J= 52 KT= 4 N(KT)= 23 IKP(KT,NI)= 52 IKC(KT,NI)= 558 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
 J= 52 TIME PERIOD 4 SECTION 3 NO. OF STUDENTS 28 ID OF STUDENTS
 1032 1110 1134 1186
 1227 1322 1325 1326
 2001 2023 2092 2121
 2129 2142 2143 2173
 2210 2218 2229 2234
 2253 2271 2278 2291
 2323 2328 2350 2353

J= 52 KT= 5 N(KT)= 15 IKP(KT,NI)= 52 IKC(KT,NI)= 558 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
 J= 52 TIME PERIOD 5 SECTION 4 NO. OF STUDENTS 33 ID OF STUDENTS
 1059 1063 1164 1245
 1250 1254 1311 1333
 2008 2024 2030 2046
 2024 2095 2144 2150
 2152 2174 2181 2187
 2193 2198 2203 2211
 2225 2237 2276 2284
 2287 2289 2308 2355
 2356

J= 52 KT= 6 N(KT)= 14 IKP(KT,NI)= 52 IKC(KT,NI)= 558 IKG(KT,NI)= IKF(KT,NI)=SECTION 5
 J= 52 TIME PERIOD 6 SECTION 5 NO. OF STUDENTS 27 ID OF STUDENTS
 1039 1056 1174 1246
 1263 1320 1328 2004
 2014 2029 2031 2070
 2081 2088 2112 2119
 2124 2125 2130 2192
 2213 2228 2262 2292
 2315 2332 2336

J= 53 SCHEDULING COURSE 556 NO CONFLICTS GENERATED

J= 53 COURSE= 556 NO. OF STUDENTS = 90 NO. OF SECTIONS = 5 SCHEDULED AT TIME PERIODS 1 2 3 4 6
J= 53 K1= 1 N(KT)= 18 IKP(KT,NI)= 53 IKC(KT,NI)= 556 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 53 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 18 ID OF STUDENTS

2003 2028 2051 2058
2093 2126 2137 2138
2149 2169 2195 2227
2256 2261 2283 2294
2317 2327

J= 53 K1= 2 N(KT)= 17 IKP(KT,NI)= 53 IKC(KT,NI)= 556 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 53 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 18 ID OF STUDENTS

2011 2037 2047 2083
2090 2118 2133 2153
2161 2168 2183 2186
2222 2272 2275 2290
2325 2346

J= 53 K1= 3 N(KT)= 18 IKP(KT,NI)= 53 IKC(KT,NI)= 556 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
J= 53 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 18 ID OF STUDENTS

2016 2038 2057 2105
2108 2120 2135 2154
2166 2179 2196 2209
2246 2267 2280 2293
2326 2344

J= 53 K1= 4 N(KT)= 24 IKP(KT,NI)= 53 IKC(KT,NI)= 556 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
J= 53 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 18 ID OF STUDENTS

2005 2039 2060 2085
2111 2127 2139 2146
2164 2182 2201 2219
2250 2273 2279 2309
2334 2357

J= 53 K1= 6 N(KT)= 15 IKP(KT,NI)= 53 IKC(KT,NI)= 556 IKG(KT,NI)= IKF(KT,NI)=SECTION 5
J= 53 TIME PERIOD 6 SECTION 5 NO. OF STUDENTS 18 ID OF STUDENTS

2040 2044 2069 2089
2109 2128 2131 2162
2167 2194 2214 2223
2257 2274 2282 2298
2311 2354

J= 54 SCHEDULING COURSE 535 NO CONFLICTS GENERATED

J= 54 COURSE= 535 NO. OF STUDENTS = 15 NO. OF SECTIONS = 5 SCHEDULED AT TIME PERIODS 1 2 3 4 5

J= 54 KT= 1 N(KT)= 19 IKP(KT,NI)= 54 IKC(KT,NI)= 535 IKG(KT,NI)= IKF(KT,NI)=SECTION 1

J= 54 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 1 ID OF STUDENTS
1011

J= 54 KT= 2 N(KT)= 10 IKP(KT,NI)= 54 IKC(KT,NI)= 535 IKG(KT,NI)= IKF(KT,NI)=SECTION 2

J= 54 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 4 ID OF STUDENTS
1002 1181 1193 1267

J= 54 KT= 3 N(KT)= 19 IKP(KT,NI)= 54 IKC(KT,NI)= 535 IKG(KT,NI)= IKF(KT,NI)=SECTION 3

J= 54 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 0 ID OF STUDENTS
1026 1141 1150 1240
1321 1338 2042 2314

J= 54 KT= 4 N(KT)= 25 IKP(KT,NI)= 54 IKC(KT,NI)= 535 IKG(KT,NI)= IKF(KT,NI)=SECTION 4

J= 54 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 1 ID OF STUDENTS
2118

J= 54 KT= 5 N(KT)= 16 IKP(KT,NI)= 54 IKC(KT,NI)= 535 IKG(KT,NI)= IKF(KT,NI)=SECTION 5

J= 54 TIME PERIOD 5 SECTION 5 NO. OF STUDENTS 1 ID OF STUDENTS
2149

J= 55 COURSE 525 CONFLICT 1 CONFLICTING STUDENT ID 2125 REQUEST LIST 24 309 525 558 571

J= 55 COURSE= 525 NO. OF STUDENTS = 4 NO. OF SECTIONS = 5 SCHEDULED AT TIME PERIODS 1 2 3 4 5
J= 55 KT= 1 N(KT)= 20 IKP(KT,NI)= 55 IKC(KT,NI)= 525 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 55 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 1 ID OF STUDENTS
0

J= 55 KT= 2 N(KT)= 19 IKP(KT,NI)= 55 IKC(KT,NI)= 525 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 55 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 1 ID OF STUDENTS
1020

J= 55 KT= 3 N(KT)= 20 IKP(KT,NI)= 55 IKC(KT,NI)= 525 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
J= 55 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 1 ID OF STUDENTS
2201

J= 55 KT= 4 N(KT)= 26 IKP(KT,NI)= 55 IKC(KT,NI)= 525 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
J= 55 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 1 ID OF STUDENTS
0

J= 55 KT= 5 N(KT)= 17 IKP(KT,NI)= 55 IKC(KT,NI)= 525 IKG(KT,NI)= IKF(KT,NI)=SECTION 5
J= 55 TIME PERIOD 5 SECTION 5 NO. OF STUDENTS 1 ID OF STUDENTS
2143

TOTAL NUMBER OF CONFLICTS = 84

TIME PERIODS

	1	2	3	4	5	6
	*****	*****	*****	*****	*****	*****
LOCKED-IN COURSES						
	578	0	538	540	550	529
	579	0	12	541	0	559
	0	0	0	551	0	561
SEVERAL-HOUR COURSES						
	570	570	570	570	0	0
	571	571	571	0	0	0
	0	0	0	572	572	572
	573	573	573	0	0	0
RESTRICTED SINGLE-SECTION COURSES						
	0	0	0	19	21	0
	0	0	0	320	321	0
	577	0	0	0	0	575
SINGLE-SECTION COURSES						
	0	16	316	555	7	532
	0	547	306	0	0	0
MULTI-SECTION COURSES						
	0	0	0	534 -1	534 -2	534 -3
	0	0	0	581 -1	581 -2	581 -3
	0	0	0	568 -1	568 -2	568 -3
	530 -1	0	0	530 -2	0	0
	0	0	5 -1	0	0	5 -2
	0	562 -1	0	562 -2	0	0
	564 -1	0	0	0	564 -2	0
	0	0	546 -1	0	0	546 -2
	324 -1	0	0	324 -2	0	0
	24 -1	0	0	0	24 -2	0
	0	305 -1	0	0	305 -2	0
	0	0	8 -1	8 -2	0	0
	0	308 -1	0	308 -2	0	0
	0	0	323 -1	323 -2	0	0
	0	23 -1	0	23 -2	0	0
	309 -1	309 -2	309 -3	309 -4	0	0
	9 -1	9 -2	9 -3	9 -4	0	0
	553 -1	553 -2	553 -3	553 -4	0	0
	0	0	549 -1	0	0	549 -2
	548 -1	548 -2	0	0	548 -3	0
	560 -1	560 -2	0	0	0	0
	557 -1	557 -2	0	0	557 -3	557 -4
	0	531 -1	531 -2	531 -3	531 -4	0
	558 -1	0	558 -2	558 -3	558 -4	558 -5
	556 -1	556 -2	556 -3	556 -4	0	556 -5
	535 -1	535 -2	535 -3	535 -4	535 -5	0
	525 -1	525 -2	525 -3	525 -4	525 -5	0

J= 56 SCHEDULING COURSE 569 NO CONFLICTS GENERATED

J= 56 COURSE= 569 K1= 218 KTN(JS)= 1 2 3 4 5 6 6
J= 56 KT= 1 N(KT)= 21 IKP(KT,NI)= 56 IKC(KT,NI)= 569 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 56 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 32 ID OF STUDENTS
2007 2019 2035 2038
2049 2057 2080 2087
2101 2109 2127 2134
2145 2162 2167 2180
2188 2202 2218 2226
2231 2247 2258 2266
2285 2298 2300 2320
2322 2331 2347 2359

J= 56 KT= 2 N(KT)= 20 IKP(KT,NI)= 56 IKC(KT,NI)= 569 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 56 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 32 ID OF STUDENTS
1169 2009 2017 2039
2045 2060 2067 2089
2094 2116 2121 2128
2148 2163 2171 2175
2185 2204 2209 2220
2230 2249 2257 2271
2281 2288 2307 2316
2328 2335 2354 2361

J= 56 KT= 3 N(KT)= 21 IKP(KT,NI)= 56 IKC(KT,NI)= 569 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
J= 56 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 31 ID OF STUDENTS
2011 2012 2025 2040
2051 2065 2076 2097
2107 2118 2123 2137
2155 2161 2169 2183
2189 2210 2212 2227
2240 2250 2260 2273
2275 2294 2309 2323
2327 2352 2357

J= 56 KT= 4 N(KT)= 27 IKP(KT,NI)= 56 IKC(KT,NI)= 569 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
J= 56 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 31 ID OF STUDENTS
2006 2015 2037 2044
2055 2063 2082 2090
2105 2115 2126 2135
2149 2166 2172 2176
2194 2214 2217 2222
2241 2251 2261 2267
2282 2293 2306 2321
2329 2344 2351

J= 56 KT= 5 N(KT)= 18 IKP(KT,NI)= 56 IKC(KT,NI)= 569 IKG(KT,NI)= IKF(KT,NI)=SECTION 5
J= 56 TIME PERIOD 5 SECTION 5 NO. OF STUDENTS 31 ID OF STUDENTS
2001 2013 2028 2048
2058 2072 2085 2096
2108 2117 2132 2138
2154 2164 2173 2182
2195 2207 2215 2227

2243	2253	2259	2272
2283	2299	2310	2318
2334	2346	2360	

J= 56 KT= 6 N(KT)= 16 IKP(KT,NI)= 56 IKC(KT,NI)= 569 IKG(KT,NI)= IKF(KT,NI)=SECTION 6
 J= 56 TIME PERIOD 6 SECTION 6 NO. OF STUDENTS 31 ID OF STUDENTS

2002	2016	2033	2047
2062	2077	2083	2092
2110	2120	2133	2152
2158	2165	2178	2184
2201	2219	2229	2238
2246	2255	2265	2279
2290	2304	2313	2325
2338	2345	2353	

J= 56 KT= 6 N(KT)= 17 IKP(KT,NI)= 56 IKC(KT,NI)= 569 IKG(KT,NI)= IKF(KT,NI)=SECTION 7
 J= 56 TIME PERIOD 6 SECTION 7 NO. OF STUDENTS 30 ID OF STUDENTS

2003	2023	2041	2059
2074	2079	2086	2093
2111	2122	2139	2153
2159	2168	2179	2196
2205	2224	2235	2239
2252	2256	2268	2280
2291	2305	2317	2326
2340	2349		

J= 57 SCHEDULING COURSE 537 NO CONFLICTS GENERATED

J= 57 COURSE= 537 KI= 210 KTN(JS)= 1 2 3 4 5 6 6
 J= 57 KT= 1 N(KT)= 22 IKP(KT,NI)= 57 IKC(KT,NI)= 537 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 57 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 35 ID OF STUDENTS

2055	2056	2062	2065
2074	2082	2085	2086
2110	2115	2117	2132
2139	2158	2163	2165
2178	2182	2189	2207
2214	2238	2243	2259
2267	2268	2272	2299
2304	2310	2313	2328
2338	2346	2353	

J= 57 KT= 2 N(KT)= 21 IKP(KT,NI)= 57 IKC(KT,NI)= 537 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 57 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 34 ID OF STUDENTS

2001	2006	2015	2023
2041	2049	2057	2072
2079	2092	2105	2123
2135	2138	2152	2167
2176	2179	2194	2201
2212	2216	2229	2235
2246	2252	2261	2265
2230	2291	2317	2326
2340	2349		

J= 57 KT= 3 N(KT)= 22 IKP(KT,NI)= 57 IKC(KT,NI)= 537 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
 J= 57 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 35 ID OF STUDENTS

2013	2017	2028	2037
2039	2047	2058	2077

2109	2122	2126	2128
2133	2145	2153	2168
2195	2202	2204	2215
2218	2219	2226	2230
2247	2251	2256	2290
2300	2305	2320	2325
2354	2359	2361	

J= 57 KT= 4 N(KT)= 20 IKP(KT,N1)= 57 IKC(KT,N1)= 537 IKG(KT,N1)= IKF(KT,N1)=SECTION 4
J= 57 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 33 ID OF STUDENTS

2009	2011	2038	2040
2045	2059	2067	2083
2089	2093	2108	2120
2134	2148	2162	2171
2175	2185	2205	2220
2223	2224	2227	2239
2249	2255	2257	2275
2288	2294	2316	2331
2345			

J= 57 KT= 5 N(KT)= 19 IKP(KT,N1)= 57 IKC(KT,N1)= 537 IKG(KT,N1)= IKF(KT,N1)=SECTION 5
J= 57 TIME PERIOD 5 SECTION 5 NO. OF STUDENTS 34 ID OF STUDENTS

2002	2007	2016	2033
2044	2051	2063	2076
2101	2107	2111	2121
2137	2159	2161	2172
2183	2188	2196	2209
2217	2222	2231	2240
2250	2266	2273	2279
2282	2298	2322	2335
2347	2351		

J= 57 KT= 6 N(KT)= 18 IKP(KT,N1)= 57 IKC(KT,N1)= 537 IKG(KT,N1)= IKF(KT,N1)=SECTION 6
J= 57 TIME PERIOD 6 SECTION 6 NO. OF STUDENTS 20 ID OF STUDENTS

2019	2035	2080	2090
2096	2118	2155	2166
2173	2210	2253	2260
2281	2285	2306	2309
2321	2327	2344	2360

J= 57 KT= 6 N(KT)= 19 IKP(KT,N1)= 57 IKC(KT,N1)= 537 IKG(KT,N1)= IKF(KT,N1)=SECTION 7
J= 57 TIME PERIOD 6 SECTION 7 NO. OF STUDENTS 19 ID OF STUDENTS

2075	2048	2087	2094
2116	2154	2164	2169
2180	2241	2258	2271
2283	2293	2307	2318
2323	2334	2357	

J= 58	COURSE 565	CONFLICT 1	CONFLICTING STUDENT 10	2037	REQUEST LIST	21	309	532	537	556	565	569
J= 58	COURSE 565	CONFLICT 2	CONFLICTING STUDENT 10	2039	REQUEST LIST	9	305	537	556	565	569	575
J= 58	COURSE 565	CONFLICT 3	CONFLICTING STUDENT 10	2040	REQUEST LIST	24	309	537	556	565	569	577
J= 58	COURSE 565	CONFLICT 4	CONFLICTING STUDENT 10	2045	REQUEST LIST	9	305	537	558	559	565	569
J= 58	COURSE 565	CONFLICT 5	CONFLICTING STUDENT 10	2067	REQUEST LIST	9	305	537	558	559	565	569
J= 58	COURSE 565	CONFLICT 6	CONFLICTING STUDENT 10	2105	REQUEST LIST	21	309	537	556	565	569	575
J= 58	COURSE 565	CONFLICT 7	CONFLICTING STUDENT 10	2120	REQUEST LIST	7	309	537	548	556	565	569
J= 58	COURSE 565	CONFLICT 8	CONFLICTING STUDENT 10	2127	REQUEST LIST	7	309	533	556	560	565	569
J= 58	COURSE 565	CONFLICT 9	CONFLICTING STUDENT 10	2137	REQUEST LIST	5	323	531	537	556	565	569
J= 58	COURSE 565	CONFLICT 10	CONFLICTING STUDENT 10	2161	REQUEST LIST	5	8	309	323	537	556	565 569
J= 58	COURSE 565	CONFLICT 11	CONFLICTING STUDENT 10	2169	REQUEST LIST	21	308	537	551	556	565	569
J= 58	COURSE 565	CONFLICT 12	CONFLICTING STUDENT 10	2185	REQUEST LIST	7	323	537	557	559	565	569
J= 58	COURSE 565	CONFLICT 13	CONFLICTING STUDENT 10	2209	REQUEST LIST	5	309	531	537	556	565	569

COURSE	CONFLICT	CONFLICTING STUDENT ID	REQUEST LIST	24	309	537	556	560	565	569
J= 58 COURSE 565	CONFLICT 14	CONFLICTING STUDENT ID 2219	REQUEST LIST	24	309	537	556	560	565	569
J= 58 COURSE 565	CONFLICT 15	CONFLICTING STUDENT ID 2239	REQUEST LIST	21	309	537	558	564	565	569
J= 58 COURSE 565	CONFLICT 16	CONFLICTING STUDENT ID 2256	REQUEST LIST	7	309	537	551	556	565	569
J= 58 COURSE 565	CONFLICT 17	CONFLICTING STUDENT ID 2261	REQUEST LIST	24	309	537	556	559	565	569
J= 58 COURSE 565	CONFLICT 18	CONFLICTING STUDENT ID 2327	REQUEST LIST	7	309	537	551	556	565	569
J= 58 COURSE 565	CONFLICT 19	CONFLICTING STUDENT ID 2351	REQUEST LIST	5	309	531	537	558	565	569

J= 58 COURSE= 565 KI= 202 KTN(JS)= 1 2 3 4 5 6 2
J= 58 KT= 1 N(KT)= 23 IKP(KT,NI)= 58 IKC(KT,NI)= 565 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 58 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 16 ID OF STUDENTS

2096	2171	2172	2194
2210	2253	2273	2306
2309	2321	2335	2340
2349	2352	2354	2361

J= 58 KT= 2 N(KT)= 22 IKP(KT,NI)= 58 IKC(KT,NI)= 565 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 58 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 29 ID OF STUDENTS

2002	2013	2028	2044
2055	2059	2066	2085
2108	2111	2132	2164
2173	2182	2109	2196
2217	2223	2243	2255
2266	2282	2298	2310
2320	2331	2334	2347
2360			

J= 58 KT= 3 N(KT)= 23 IKP(KT,NI)= 58 IKC(KT,NI)= 565 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
J= 58 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 27 ID OF STUDENTS

2019	2041	2049	2072
2074	2083	2086	2134
2138	2139	2158	2167
2180	2184	2205	2208
2229	2231	2235	2257
2265	2271	2283	2299
2313	2329	2338	

J= 58 KT= 4 N(KT)= 29 IKP(KT,NI)= 58 IKC(KT,NI)= 565 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
J= 58 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 20 ID OF STUDENTS

2012	2017	2047	2048
2077	2080	2094	2101
2116	2117	2122	2128
2153	2159	2179	2183
2204	2226	2230	2252
2272	2280	2285	2290
2305	2317	2326	2346

J= 58 KT= 5 N(KT)= 20 IKP(KT,NI)= 58 IKC(KT,NI)= 565 IKG(KT,NI)= IKF(KT,NI)=SECTION 5
J= 58 TIME PERIOD 5 SECTION 5 NO. OF STUDENTS 20 ID OF STUDENTS

2003	2009	2023	2038
2079	2087	2090	2109
2123	2135	2148	2163
2168	2176	2202	2220
2227	2241	2247	2251
2267	2275	2280	2293
2307	2323	2328	2357

J= 58 KT= 6 N(KT)= 21 IKP(KT,NI)= 58 IKC(KT,NI)= 565 IKG(KT,NI)= IKF(KT,NI)=SECTION 6
J= 58 TIME PERIOD 6 SECTION 6 NO. OF STUDENTS 27 ID OF STUDENTS

2006	2015	2035	2065
2002	2089	2097	2115
2133	2145	2162	2165

2175	2201	2212	2224
2238	2246	2249	2260
2268	2281	2291	2304
2316	2325	2344	

J= 50 KT= 2 N(KT)= 23 IKP(KT,NI)= 58 IKC(KT,NI)= 565 IKG(KT,NI)= IKF(KT,NI)=SECTION 7
J= 50 TIME PERIOD 2 SECTION 7 NO. OF STUDENTS 28 ID OF STUDENTS

2007	2025	2033	2051
2058	2063	2076	2107
2110	2126	2155	2166
2178	2188	2195	2207
2218	2240	2250	2259
2279	2294	2300	2318
2322	2333	2345	2359

J= 59 COURSE 567 CONFLICT 1 CONFLICTING STUDENT ID 2093 REQUEST LIST 23 24 309 316 537 556 567 569

J= 59 COURSE= 567 KI= 19J KTN(JS)= 1 2 3 4 5 6 6
J= 59 KT= 1 N(KT)= 24 IKP(KT,NI)= 59 IKC(KT,NI)= 567 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 59 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 29 ID OF STUDENTS

1001	1022	1037	1052
1055	1071	1081	1100
1104	1115	1133	1146
1163	1180	1199	1208
1223	1235	1238	1262
1275	1282	1295	1300
1324	1339	1342	2018
2324			

J= 59 KT= 2 N(KT)= 24 IKP(KT,NI)= 59 IKC(KT,NI)= 567 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 59 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 28 ID OF STUDENTS

1009	1031	1039	1053
1067	1074	1090	1098
1107	1124	1135	1147
1170	1184	1201	1209
1228	1236	1243	1268
1280	1288	1297	1302
1330	1337	1347	2073

J= 59 KT= 3 N(KT)= 24 IKP(KT,NI)= 59 IKC(KT,NI)= 567 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
J= 59 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 29 ID OF STUDENTS

1010	1020	1042	1045
1068	1075	1084	1102
1105	1120	1142	1149
1177	1192	1202	1214
1225	1237	1240	1271
1279	1285	1294	1303
1327	1345	2021	2242

J= 59 KT= 4 N(KT)= 30 IKP(KT,NI)= 59 IKC(KT,NI)= 567 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
J= 59 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 28 ID OF STUDENTS

1012	1036	1040	1066
1070	1080	1095	1101
1119	1127	1137	1155
1175	1194	1203	1220
1233	1242	1261	1272
1281	1287	1299	1310
1332	1342	2062	2180

J= 59 KT= 5 N(KT)= 22 IKP(KT,NI)= 59 IKC(KT,NI)= 567 IKG(KT,NI)= IKF(KT,NI)=SECTION 5

J= 59
J= 59 TIME PERIOD 5 SECTION 5 NO. OF STUDENTS 27

ID OF STUDENTS			
1016	1020	1051	1065
1078	1083	1092	1103
1114	1128	1140	1156
1187	1188	1204	1218
1234	1247	1260	1269
1284	1289	1307	1323
1335	1341	2060	

J= 59 KT= 6 N(KT)= 20 IKP(KT,NI)= 59 IKC(KT,NI)= 567 IKG(KT,NI)= IKF(KT,NI)=SECTION 6
J= 59 TIME PERIOD 6 SECTION 6 NO. OF STUDENTS 26

ID OF STUDENTS			
1017	1030	1047	1069
1075	1085	1095	1112
1126	1130	1157	1162
1179	1198	1205	1224
1230	1251	1259	1273
1277	1291	1309	1318
1344	1351		

J= 59 KT= 6 N(KT)= 21 IKP(KT,NI)= 59 IKC(KT,NI)= 567 IKG(KT,NI)= IKF(KT,NI)=SECTION 7
J= 59 TIME PERIOD 6 SECTION 7 NO. OF STUDENTS 26

ID OF STUDENTS			
1018	1038	1054	1073
1082	1087	1097	1113
1129	1132	1159	1166
1183	1200	1207	1225
1232	1252	1264	1276
1283	1292	1317	1329
1349	1353		

J= 60 SCHEDULING COURSE 580 NO CONFLICTS GENERATED

J= 60 COURSE= 580 K1= 191 KTN(JS)= 1 2 3 4 5 6 6 5
J= 60 KT= 1 N(KT)= 25 IKP(KT,NI)= 60 IKC(KT,NI)= 580 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
J= 60 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 28

ID OF STUDENTS			
1016	1030	1047	1053
1065	1074	1087	1090
1096	1101	1127	1132
1184	1188	1194	1202
1207	1220	1224	1237
1277	1292	1302	1303
1335	1344	1345	2073

J= 60 KT= 2 N(KT)= 25 IKP(KT,NI)= 60 IKC(KT,NI)= 580 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
J= 60 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 28

ID OF STUDENTS			
1001	1022	1036	1052
1069	1075	1081	1085
1095	1112	1126	1149
1163	1166	1179	1198
1204	1218	1232	1238
1260	1282	1269	1318
1341	1348	1353	2050

J= 60 KT= 3 N(KT)= 25 IKP(KT,NI)= 60 IKC(KT,NI)= 580 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
J= 60 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 27

ID OF STUDENTS			
1009	1028	1051	1055
1070	1078	1083	1092
1100	1124	1128	1155

1156	1170	1100	1203
1235	1230	1235	1259
1273	1283	1300	1307
1329	1337	1351	

J= 60 KT= 4 N(KT)= 31 IKP(KT,NI)= 60 IKC(KT,NI)= 580 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
J= 60 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 27 ID OF STUDENTS

1010	1038	1039	1045
1071	1084	1097	1107
1113	1120	1135	1162
1177	1183	1208	1209
1243	1248	1252	1264
1275	1278	1324	1329
1342	1347	2360	

J= 60 KT= 5 N(KT)= 23 IKP(KT,NI)= 60 IKC(KT,NI)= 580 IKG(KT,NI)= IKF(KT,NI)=SECTION 5
J= 50 TIME PERIOD 5 SECTION 5 NO. OF STUDENTS 25 ID OF STUDENTS

1017	1037	1066	1068
1075	1082	1104	1115
1130	1147	1159	1200
1214	1225	1233	1261
1268	1279	1281	1288
1294	1309	1317	1339
1349			

J= 60 KT= 6 N(KT)= 22 IKP(KT,NI)= 60 IKC(KT,NI)= 580 IKG(KT,NI)= IKF(KT,NI)=SECTION 6
J= 60 TIME PERIOD 6 SECTION 6 NO. OF STUDENTS 16 ID OF STUDENTS

1012	1031	1042	1103
1119	1137	1142	1175
1192	1234	1242	1269
1272	1287	1297	1323

J= 60 KT= 6 N(KT)= 23 IKP(KT,NI)= 60 IKC(KT,NI)= 580 IKG(KT,NI)= IKF(KT,NI)=SECTION 7
J= 60 TIME PERIOD 6 SECTION 7 NO. OF STUDENTS 16 ID OF STUDENTS

1020	1040	1102	1114
1133	1140	1169	1187
1228	1236	1247	1271
1284	1295	1316	1332

J= 60 KT= 5 N(KT)= 24 IKP(KT,NI)= 60 IKC(KT,NI)= 580 IKG(KT,NI)= IKF(KT,NI)=SECTION 8
J= 60 TIME PERIOD 5 SECTION 8 NO. OF STUDENTS 24 ID OF STUDENTS

1018	1054	1067	1073
1080	1098	1105	1129
1146	1157	1199	1201
1223	1229	1251	1262
1276	1280	1285	1291
1299	1310	1327	1342

J= 61	COURSE 533	CONFLICT 1	CONFLICTING STUDENT ID	1009	REQUEST LIST 531	533	550	553	567	580	
J= 61	COURSE 533	CONFLICT 2	CONFLICTING STUDENT ID	1028	REQUEST LIST 533	548	563	560	567	580	
J= 61	COURSE 533	CONFLICT 3	CONFLICTING STUDENT ID	1039	REQUEST LIST 5	321	533	558	564	567	580
J= 61	COURSE 533	CONFLICT 4	CONFLICTING STUDENT ID	1045	REQUEST LIST 24	305	533	546	557	567	580
J= 61	COURSE 533	CONFLICT 5	CONFLICTING STUDENT ID	1066	REQUEST LIST 531	533	553	567	577	580	
J= 61	COURSE 533	CONFLICT 6	CONFLICTING STUDENT ID	1090	REQUEST LIST 5	21	321	323	533	553	567 580
J= 61	COURSE 533	CONFLICT 7	CONFLICTING STUDENT ID	1156	REQUEST LIST 5	324	530	533	553	567	580
J= 61	COURSE 533	CONFLICT 8	CONFLICTING STUDENT ID	1166	REQUEST LIST 24	308	531	533	553	567	580
J= 61	COURSE 533	CONFLICT 9	CONFLICTING STUDENT ID	1224	REQUEST LIST 8	305	531	533	553	567	580
J= 61	COURSE 533	CONFLICT 10	CONFLICTING STUDENT ID	1277	REQUEST LIST 8	305	531	533	553	567	580
J= 61	COURSE 533	CONFLICT 11	CONFLICTING STUDENT ID	1282	REQUEST LIST 5	324	533	557	561	567	580
J= 61	COURSE 533	CONFLICT 12	CONFLICTING STUDENT ID	1284	REQUEST LIST 5	324	533	546	553	567	580

1020	1040	1102	1119
1142	1152	1242	1295
1316	1329	1342	

J= 61 KT= 2 N(KT)= 27 IKP(KT,NI)= 61 IKC(KT,NI)= 533 IKG(KT,NI)= IKF(KT,NI)=SECTION 7
 J= 61 TIME PERIOD 2 SECTION 7 NO. OF STUDENTS 28

ID OF STUDENTS
 1017 1027 1054 1065
 1078 1092 1100 1103
 1115 1157 1175 1180
 1188 1200 1203 1208
 1225 1235 1248 1264
 1272 1275 1283 1303
 1310 1335 1344 2352

J= 62 COURSE 552 CONFLICT 1 CONFLICTING STUDENT ID 1040 REQUEST LIST 24 308 531 533 552 567 580
 J= 62 COURSE 552 CONFLICT 2 CONFLICTING STUDENT ID 1075 REQUEST LIST 8 324 531 533 552 567 580
 J= 62 COURSE 552 CONFLICT 3 CONFLICTING STUDENT ID 1100 REQUEST LIST 5 308 531 533 552 567 580
 J= 62 COURSE 552 CONFLICT 4 CONFLICTING STUDENT ID 1107 REQUEST LIST 5 323 533 552 560 567 580
 J= 62 COURSE 552 CONFLICT 5 CONFLICTING STUDENT ID 1112 REQUEST LIST 21 324 533 546 552 567 580
 J= 62 COURSE 552 CONFLICT 6 CONFLICTING STUDENT ID 1169 REQUEST LIST 8 324 533 538 552 569 580
 J= 62 COURSE 552 CONFLICT 7 CONFLICTING STUDENT ID 1197 REQUEST LIST 5 308 531 552 571
 J= 62 COURSE 552 CONFLICT 8 CONFLICTING STUDENT ID 1244 REQUEST LIST 7 323 552 559 573
 J= 62 COURSE 552 CONFLICT 9 CONFLICTING STUDENT ID 1279 REQUEST LIST 5 324 531 533 552 567 580
 J= 62 COURSE 552 CONFLICT 10 CONFLICTING STUDENT ID 1286 REQUEST LIST 5 324 550 552 573
 J= 62 COURSE 552 CONFLICT 11 CONFLICTING STUDENT ID 1318 REQUEST LIST 19 321 533 552 567 577 580
 J= 62 COURSE 552 CONFLICT 12 CONFLICTING STUDENT ID 2226 REQUEST LIST 7 309 537 552 559 565 569

J= 62 COURSE= 552 K1= 172 KTN(JS)= 1 2 3 4 5 6
 J= 62 KT= 1 N(KT)= 27 IKP(KT,NI)= 62 IKC(KT,NI)= 552 IKG(KT,NI)= IKF(KT,NI)=SECTION 1
 J= 62 TIME PERIOD 1 SECTION 1 NO. OF STUDENTS 15

ID OF STUDENTS
 1066 1075 1084 1102
 1130 1157 1175 1201
 1228 1229 1236 1291
 1297 1309 1310

J= 62 KT= 2 N(KT)= 28 IKP(KT,NI)= 62 IKC(KT,NI)= 552 IKG(KT,NI)= IKF(KT,NI)=SECTION 2
 J= 62 TIME PERIOD 2 SECTION 2 NO. OF STUDENTS 26

ID OF STUDENTS
 1070 1093 1119 1133
 1142 1161 1192 1199
 1295 1324 2019 2035
 2062 2065 2096 2097
 2117 2145 2100 2202
 2215 2231 2247 2258
 2268 2338

J= 62 KT= 3 N(KT)= 27 IKP(KT,NI)= 62 IKC(KT,NI)= 552 IKG(KT,NI)= IKF(KT,NI)=SECTION 3
 J= 62 TIME PERIOD 3 SECTION 3 NO. OF STUDENTS 35

ID OF STUDENTS
 1012 1031 1037 1053
 1065 1071 1113 1114
 1146 1162 1163 1194
 1200 1209 1220 1223
 1225 1242 1247 1251
 1264 1272 1292 1299
 1329 1344 1349 1353
 2002 2172 2241 2243
 2249 2266 2335

J= 62 KT= 4 N(KT)= 33 IKP(KT,NI)= 62 IKC(KT,NI)= 552 IKG(KT,NI)= IKF(KT,NI)=SECTION 4
 J= 62 TIME PERIOD 4 SECTION 4 NO. OF STUDENTS 35

ID OF STUDENTS
 1017 1022 1025 1042
 1051 1052 1060 1067
 1078 1109 1176 1189
 1258 1314 1315 1335
 2074 2086 2123 2132
 2170 2186 2188 2189

2235	2238	2240	2265
2281	2299	2307	2310
2313	2322	2347	

J= 62 KT= 5 N(KT)= 27 IKP(KT,NI)= 02 IKC(KT,NI)= 552 IKG(KT,NI)= IKF(KT,NI)=SECTION 5
J= 62 TIME PERIOD 5 SECTION 5 NO. OF STUDENTS 29 ID OF STUDENTS

1001	1023	1029	1030
1044	1099	1106	1123
1143	1154	1168	1178
1190	1222	1248	1308
2017	2049	2077	2116
2134	2158	2171	2205
2265	2300	2305	2320
2340			

J= 62 KT= 6 N(KT)= 24 IKP(KT,NI)= 02 IKC(KT,NI)= 552 IKG(KT,NI)= IKF(KT,NI)=SECTION 6
J= 62 TIME PERIOD 6 SECTION 6 NO. OF STUDENTS 20 ID OF STUDENTS

1004	1013	1034	1043
1054	1076	1122	1151
1153	1165	1185	1206
1239	1253	1256	1304
2013	2157	2329	2358

TOTAL NUMBER OF CONFLICTS = 48

	TIME PERIODS					
	1	2	3	4	5	6
LOCKED-IN COURSES	*****	*****	*****	*****	*****	*****
	578	0	548	540	550	529
	579	0	12	541	0	559
	0	0	0	551	0	561
SEVERAL-MCUM COURSES						
	570	570	570	570	0	0
	571	571	571	0	0	0
	0	0	0	572	572	572
	573	573	573	0	0	0
RESTRICTED SINGLE-SECTION COURSES						
	0	0	0	19	21	0
	0	0	0	320	321	0
	577	0	0	0	0	575
SINGLE-SECTION COURSES						
	0	10	310	555	7	532
	0	547	306	0	0	0
MULTI-SECTION COURSES						
	0	0	0	534 -1	534 -2	534 -3
	0	0	0	581 -1	581 -2	581 -3
	0	0	0	568 -1	568 -2	568 -3
	530 -1	0	0	530 -2	0	0
	0	0	5 -1	0	0	5 -2
	0	562 -1	0	562 -2	0	0
	564 -1	0	0	0	564 -2	0
	0	0	546 -1	0	0	546 -2
	324 -1	0	0	324 -2	0	0
	24 -1	0	0	0	24 -2	0
	0	305 -1	0	0	305 -2	0
	0	0	8 -1	8 -2	0	0
	0	308 -1	0	308 -2	0	0
	0	0	523 -1	523 -2	0	0
	0	23 -1	0	23 -2	0	0
	309 -1	309 -2	309 -3	309 -4	0	0
	9 -1	9 -2	9 -3	9 -4	0	0
	553 -1	553 -2	553 -3	553 -4	0	0
	0	0	549 -1	0	0	549 -2
	548 -1	548 -2	0	0	548 -3	0
	560 -1	560 -2	0	0	0	0
	557 -1	557 -2	0	0	557 -3	557 -4
	0	531 -1	531 -2	531 -3	531 -4	0
	558 -1	0	556 -2	556 -2	556 -4	556 -5
	556 -1	556 -2	556 -3	556 -4	0	556 -5
	525 -1	525 -2	525 -3	525 -4	525 -5	0
	566 -1	566 -2	566 -3	566 -4	566 -5	566 -6
	0	0	0	0	0	566 -7
	537 -1	537 -2	537 -3	537 -4	537 -5	537 -6
	0	0	0	0	0	537 -7
	565 -1	565 -2	565 -3	565 -4	565 -5	0
	0	0	0	0	565 -6	0
	0	565 -7	0	0	0	0
	567 -1	567 -2	567 -3	567 -4	567 -5	567 -6
	0	0	0	0	0	567 -7
	580 -1	580 -2	580 -3	580 -4	580 -5	580 -6
	0	0	0	0	0	580 -7
	0	0	0	0	580 -8	0
	533 -1	533 -2	533 -3	533 -4	533 -5	0
	0	0	0	0	533 -6	0
	0	533 -7	0	0	0	0
	552 -1	552 -2	552 -3	552 -4	552 -5	552 -6