

REVERSIBLE SEISMIC DATA COMPRESSION

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CHAPTER I

OVERVIEW OF DATA COMPRESSION

Introduction

At the present time, analysis of seismic data is based upon a bulk of data collected from the field, recorded onto tapes, and transported to the research center where it is analyzed for possible use. This procedure introduces a delay in the analysis and usually will not allow for on-line decision. With recent developments in telecommunications, it has been found to be desirable to transmit the data via satellite in order to reduce delay time. However, the bit rate for telecommunications is limited and the cost for transmitting a large amount of seismic data may be prohibitive. For these reasons, the seismic data need to be compressed before transmission over communication channels. This thesis is devoted to the development of a seismic data compression method for satellite transmission.

Data Compression

Data compression can be viewed as any method of representing source data in an efficient manner while maintaining the required information content. An efficient representation of the source data can be obtained first, by reducing the redundant information; second, by reducing the bandwidth; and third, by efficient coding techniques [1]. The usage of the term "data" in this thesis is to indicate any digitized signal to be

transmitted or stored. In order to avoid confusion, "information" is defined as a measure of the full range of important features contained in a set of data. This definition differs from the information theoretic definition, where it is used as a quantitative measure of information content. Both definitions are used in this work and the difference will be apparent from context. Such data compression is required due to storage constraints, digital computer memory size, limited bandwidth in communication links, limited capacity in channels, or by the desire of extracting important attributes from the source data.

Data compression can be defined in mathematical terms using the concept of entropy which, in simple terms, is referred to as the average information content measured as a number of information units (bits). The mathematical derivation of data compression, alternatively entropy compression, is given below.

The entropy can be written as [2]

$$H = - \sum_{i=1}^D p_i \log_2 p_i \quad (\text{bits/configuration}) \quad D = N^K \quad (1.1)$$

where K is the number of data, N is the number of quantization levels of data, and p_i is the probability of the i th configuration. Also, the information content of the i th configuration is given by

$$I_i = -\log_2 p_i \quad (\text{bits}). \quad (1.2)$$

When there is no redundancy in the data or every configuration has the same probability, the entropy is maximum and it is expressed as

$$H_{\max} = K \log_2 N \quad (\text{bits/configuration}). \quad (1.3)$$

The redundancy ratio (R) is defined by

$$R = \frac{H_{\max} - H}{H_{\max}} . \quad (1.4)$$

The entropy compression ratio (C) can be defined by

$$C = \frac{H_{\max}}{H} = \frac{K \log_2 N}{\sum_{i=1}^D p_i \log_2 p_i} = \frac{1}{1-R} . \quad (1.5)$$

Equation (1.5) shows that the maximum entropy of transformed data can approach their own entropy by reducing the redundancy of the data (K), and/or by reducing the redundancy from the quantization levels (N), both of which result in reducing the number of possible configurations (D).

Most of the present data compression methods are combinations of the following methods; these include: entropy reducing transform [2], prediction [3], interpolation [4], orthogonal transforms [5], and digital code representations [6]. The five general data compression techniques mentioned above are examined in relation to entropy compression below.

Some entropy reducing transforms achieve data compression by deleting a portion of information in source data. For example, filtering (low-pass, high-pass, or band-pass) is one form of an entropy reducing transformation. Usually this method is applied when intelligibility is the main objective, and a portion of the data in terms of either time or frequency is used in extracting the desired information.

The prediction method can be used to compress data when a system model is available. Assuming that the system model is incorporated in a prediction algorithm at the receiver and the transmitter has a way of acknowledging the predicted value with no channel errors, it is acceptable

to omit data which can be predicted at the receiver [3]. It is clear that the number of data to be transmitted can be reduced, and thus maximum entropy is compressed. Another advantage of this method, for entropy compression, can be realized by transmitting the difference between the source data and the predicted value (prediction error). This advantage can be seen because the prediction error variance is less than the variance of the source data. The minimization of the prediction error variance can be achieved by various techniques. These will be discussed in detail later.

The interpolation method [4] estimates the values between a given transmitted value and the most distant possible point, such that the maximum interpolation error is below the preset threshold. One simple example of interpolation techniques in data compression would be approximation of the source data by polynomial segments.

Orthogonal transforms are used in data compression because of their properties in representing data with linearly independent eigenvectors, thus reducing the redundancies. Data compression can be achieved if fewer numbers of transformed vector coefficients can represent the source data. This depends on the statistical characteristics of the source data.

Finally, the digital code representation achieves data compression by reducing the number of quantization levels while maintaining tolerable quantization error. When sampling an analog signal, the original bit rate is determined by the number of quantization levels employing the maximum entropy (Equation (1.3)). Assuming that there is a high rate of redundancy, the original bit rate can be reduced with an efficient code representation, thus achieving entropy compression.

Common application of data compression are found in systems for communications, speech and image processing, and pattern recognition [1].

A Thesis Review

Chapter II examines seismic data acquisition and reviews previous efforts in seismic data compression. In this thesis vibroseis data and impulsive seismic data are used for the proposed data compression. In order to visualize the characteristics of seismic data, statistics are measured and displayed. The proposed approach for "reversible seismic data compression" is discussed in relation to the statistical characteristics of seismic data.

In Chapter III, the proposed techniques are compared with other candidates and their performances are evaluated. First, orthogonal transforms, such as the Karhunen-Loéve transform, the Walsh-Hadamard transform, the discrete Fourier transform, and the discrete cosine transform are briefly examined and their compression ratios on seismic data are compared. Second, digital coding techniques are discussed in two steps, quantization methods and cosine representations. Brief derivations of signal-to-noise ratios are discussed for each technique, and their performance on seismic data compression is compared.

Chapter IV discusses compression methods for vibroseis data and impulsive seismic data. The vibroseis data compression is examined in the following order. First, the hybrid technique is introduced with respect to its compression ratio and signal-to-noise ratio relations. Second, the selection method is examined in terms of threshold. Third, μ -law quantization is studied in detail. Fourth, the implementation considerations are discussed. The impulsive seismic data compression is examined

in three stages: data slicing, predictive coding using an optimum linear predictor, and its implementation. Finally, the results are evaluated for both types of seismic data.

In Chapter V, the general applications of satellite communications are examined. Compressed seismic data transmission via satellite is studied with respect to its time and space configurations.

Finally, Chapter VI suggests possible future research areas for the "reversible seismic data compression."

Summary

The main interests in data compression techniques are the compression ratio (CPR) and the distortion rate, often called the "signal-to-noise ratio" (SNR) [1]. This study defines the CPR as the ratio between the bit rate of the original signal and the compressed signal. The SNR is defined as the ratio between energy of the original signal and the noise, where the noise indicates the difference between the original signal and the reconstructed signal.

Data compression techniques can be divided into two categories, irreversible and reversible compression. In order to distinguish one from the other, reversibility should be defined. In general, techniques which can reconstruct the original data with an adequate error fidelity criterion are defined as reversible data compression techniques. When the original data cannot be reconstructed due to compression techniques, they are called irreversible techniques [2].

The entropy reducing transforms are irreversible data compression techniques, since some parts of the original signal have been discarded. For irreversible techniques, SNR cannot be considered. Rather, subjective

measurements, such as human intelligibility, are considered as a measure of their performances.

The prediction, interpolation, orthogonal transforms, and digital coding representation methods are all reversible. However, the reconstructed data include the process error. The process error is defined as all the errors involved in techniques of compression and decompression. They include prediction error, interpolation error, transformation error, and quantization error. For reversible techniques, SNR is the measure of their performance. However, it is hard to determine the threshold of error fidelity when the original data have involved large noise. For this reason, the original data are assumed to be ideal in this thesis.

The development of a "reversible seismic data compression" technique is pursued, where "reversible" indicates that the original seismic data can be reconstructed with approximately a 30 dB signal-to-noise ratio. Figure 1 gives the block diagram identifying the noise (E_1) introduced at various locations. In this study the processed noise, denoted by E_2 , is of interest.

The next chapter deals with seismic data acquisition methods, previous work, the proposed approach for seismic data compression, and statistical characteristics of seismic data.

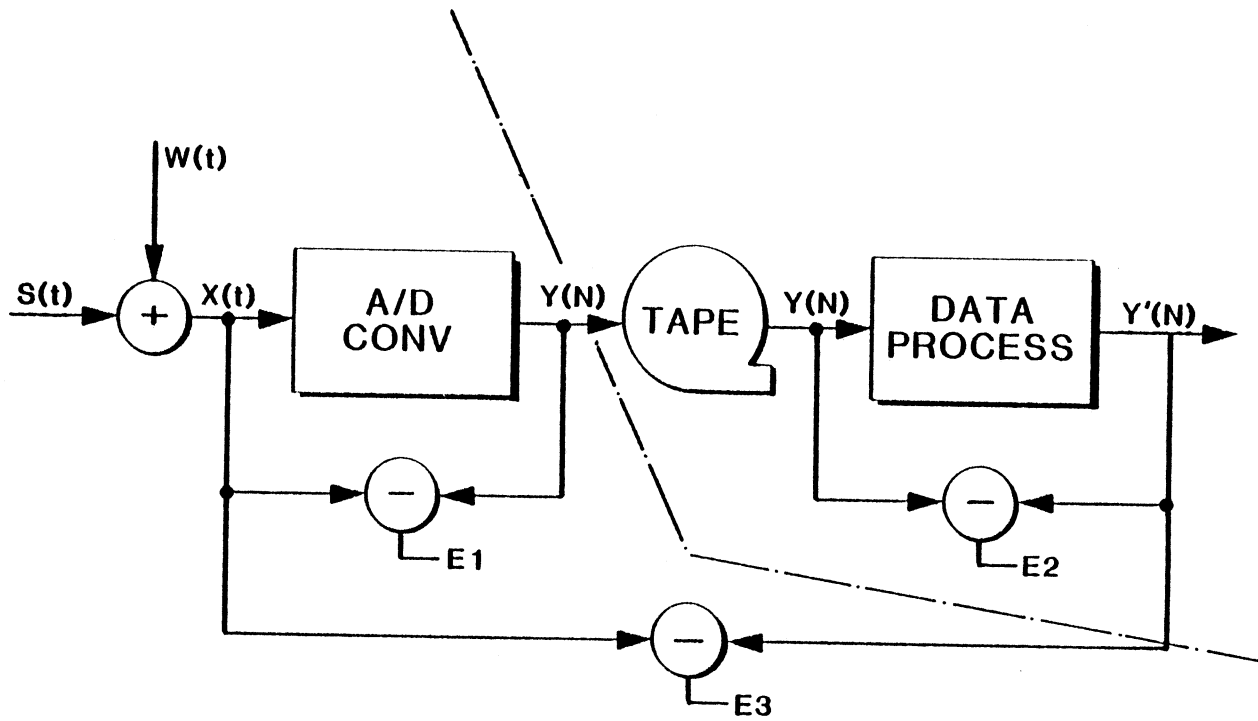


Figure 1. Simulation of Seismic Data Compression Techniques

CHAPTER II

GENERAL CONSIDERATIONS ON SEISMIC

DATA COMPRESSION

Introduction

In Chapter I, a brief overview of data compression was introduced, with respect to general compression techniques in relation to entropy reduction.

This chapter focuses on several seismic data compression considerations. First, the basic concepts of seismic data acquisition are examined, with example cases of acquisition corresponding to vibroseis and impulsive inputs. Second, previous work on seismic data compression is reviewed and its performance is evaluated with respect to reversible considerations. Third, a reversible seismic data compression scheme is viewed and directions for this proposed approach are discussed. Finally, various statistical characteristics of seismic data are examined and some important results are displayed.

Seismic Data Acquisition

As mentioned earlier, seismic data are collected from the field and recorded onto digital tapes. The essence of such a data collection system is illustrated in Figure 2 [7]. Robinson and Treitel describe the seismic data acquisition procedures as follows:

Disturbances created by seismic energy sources propagate through the earth, where interfaces between geophysical strata reflect spreading wave fronts. The receivers shown in Figure 1-2 [Figure 2 in this thesis] actually represent a composite group of transducers (seismometers). These groups may consist of up to 100 individual geophones laid out in various linear and spatial patterns, with group intervals (between group distance) ranging from 50 to 900 ft. Each time a source is activated it is common practice to record either 24, 48, or 96 group traces on digital tape simultaneously as a single recording (p. 74).

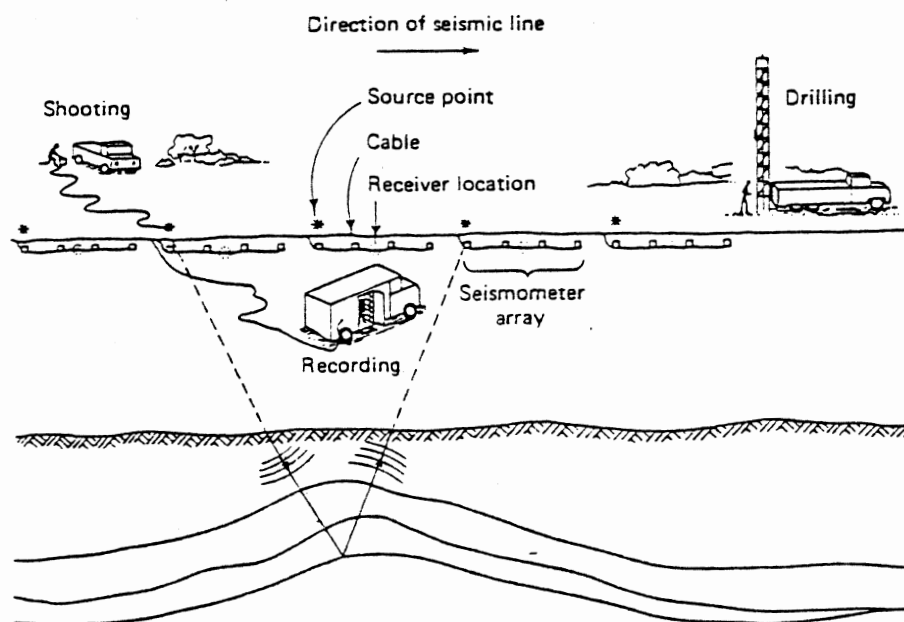


Figure 2. Seismic Data Acquisition (After Robinson and Treitel [7])

Robinson and Treitel state that many different types of energy sources are used to generate seismic waves. Dynamite and other high-energy explosive sources provide the simplest and most efficient means of releasing energy, but environmental considerations have led to the development of many alternative sources, such as explosive air guns,

electrical sparks, vibrating chirp systems, and so on. The different types of energy sources require different compression techniques.

Seismic data compression will be investigated with vibroseis data and impulsive seismic data in this thesis, where vibroseis data refer to seismic data that are collected from the reflected wave generated by a vibrating chirp disturbance propagating through the earth, and impulsive seismic data refer to seismic data that are collected from the reflected wave generated by a high-energy explosive source disturbance propagated through the earth. Examples of both types of seismic data acquisition environments are illustrated below.

In Appendix A, a set of vibroseis data is given. The data were sampled at 250 Hz for 19 seconds. Each time a vibrating chirp system was initiated, a group of 48 traces was simultaneously recorded onto a digital tape as a single recording, and is defined as a record. A 20 bit code consisting of 16 bits of mantissa and 4 bits of gain was used in this representation. The vibrating chirp system propagated a chirp signal ranging from 10 to 55 Hz for 14 seconds over the test interval.

Also, in Appendix A, a set of impulsive seismic data is given. The data were sampled at 1000 Hz for 1 second. Each time an explosion was initiated, a group of 198 traces was simultaneously recorded onto digital tape in a single recording. The same digital coding representation technique was used as in the vibroseis data case.

Due to the constraints in the bandwidth of communication links, and in storage, the bit rate (bits/second) and the block size (bits/block) are two important parameters for data compression considerations. A block is a number of seismic source initiations. In the following, the values of these parameters for the seismic data shown in Appendix A are given.

The bit rate (B) is

$$B = N_s \cdot F_s \cdot b \text{ (bits/second)} \quad (2.1)$$

where N_s is the number of traces in a record, F_s is the sampling rate, and b is the number of bits per sample. The block size (K) of the seismic data is

$$K = B \cdot N_t \cdot T \text{ (bits/block)} \quad (2.2)$$

where N_t is the number of tests performed and T is the sampling duration.

Using Equation (2.1), the bit rate of the vibroseis data (B_v) in Appendix A is

$$B_v = 48 \cdot 250 \cdot 20 = 192 \text{ k (bits/second)}$$

where $N_s = 48$, $F_s = 250 \text{ Hz}$ and $b = 20$ bits/sample. The bit rate of the impulsive seismic data (B_i) in Appendix A is

$$B_i = 198 \cdot 1000 \cdot 20 = 3920 \text{ k (bits/second)}$$

where $N_s = 198$, $F_s = 1000 \text{ Hz}$ and $b = 20$ bits/sample. Thus, the block size of the vibroseis data (K_v) is

$$K_v = 192\text{k} \cdot 16 \cdot 19 = 73 \text{ M (bits/block)}$$

where $B_v = 192\text{k}$, $N_t = 16$, and $T = 19$ seconds. The block size of the impulsive seismic data (K_i) is

$$K_i = 3960\text{k} \cdot 9 \cdot 1 = 1.8 \text{ M (bits/block)}$$

where $B_i = 3960\text{k}$, $N_t = 9$, and $T = 1$ second.

In general, 2400, 4800, 9600, and/or 56k baud rate (bits/second) are available for most communication links. Also most small computers have less than 256 k byte (2048 k bits) of memory size. Considering

the above figures, the values of bit rate and block size for seismic data indicate that data compression is necessary in order to utilize a modern telecommunication system for transmission and for processing with mini- or microcomputers.

Data compression techniques have been investigated in many areas, such as speech and image processing. However, seismic data compression has not been investigated as thoroughly as the others. A review of one previous effort on seismic data compression is discussed in the next section.

Previous Work

A survey of previous seismic data compression yields one notable work by Wood [8], in which he used bandwidth limiting and efficient digital coding techniques. The results were obtained using two methods on a set of vibroseis data sampled at 500 Hz. These are data resampling and interpolation in the time domain and the Walsh transform, and sequency limiting and the Walsh inverse transform in the Walsh domain. These techniques are examined and evaluated below. In his paper, the resampling technique is a time domain compression technique described as follows.

From sampling theory [9], if a signal is bandlimited and sampled with anti-aliasing considerations, no information is lost through proper resampling and interpolation. Based on this statement, resampling (with averaging) is applied at $1/2$, $1/4$, and $1/8$ of the original sampling rate at the transmitter. Then interpolation is performed at the receiver.

In the second method, the seismic data are transformed using the Walsh transform. Then sequency limiting is performed, by windowing the first major sequency range that contains 80 to 85 percent of the total

energy. At the receiver, the data are buffered to their original size and then the inverse Walsh transform is applied. It has been found that resampling in the time domain amounts to sequencing limiting in the Walsh domain and interpolation in the time domain amounts to the inverse Walsh transform. Flow charts corresponding to the above two methods are shown in Figure 3.

In the second stage, nonuniform quantization was applied with variable code word lengths for both methods. As illustrated in Figure 4, the step sizes are computed according to the probability densities of the data. The code word lengths are determined by the information rate of the quantization levels. For example, the first step size was illustrated to be 25.1, where 50 percent of the data could be represented. As shown in Equation (1.2), the information rate was computed as 1 bit; thus the code length was determined to be 1 bit.

In Wood's work [8], visible degradation in the final plotted seismic data was the measure of performance. Reversibility was not considered, since teleprocessing of the plot data was the major interest in his compression. In order to evaluate the applicability on reversible seismic data compression, the signal-to-noise ratios of the compression method given by Wood was estimated according to the percentage of the total energy he claimed to maintain in the Walsh domain. It was found to be approximately 7 dB without considering the processing error.

Wood pursued further compression by applying nonuniform quantization using an average of 3 bits per sample. This quantization introduces significant error. Though he claimed to achieve a 16:1 compression ratio, the signal-to-noise ratio (<7 dB) is too small to be of use for reversible seismic data compression. Also, it should be pointed out that the

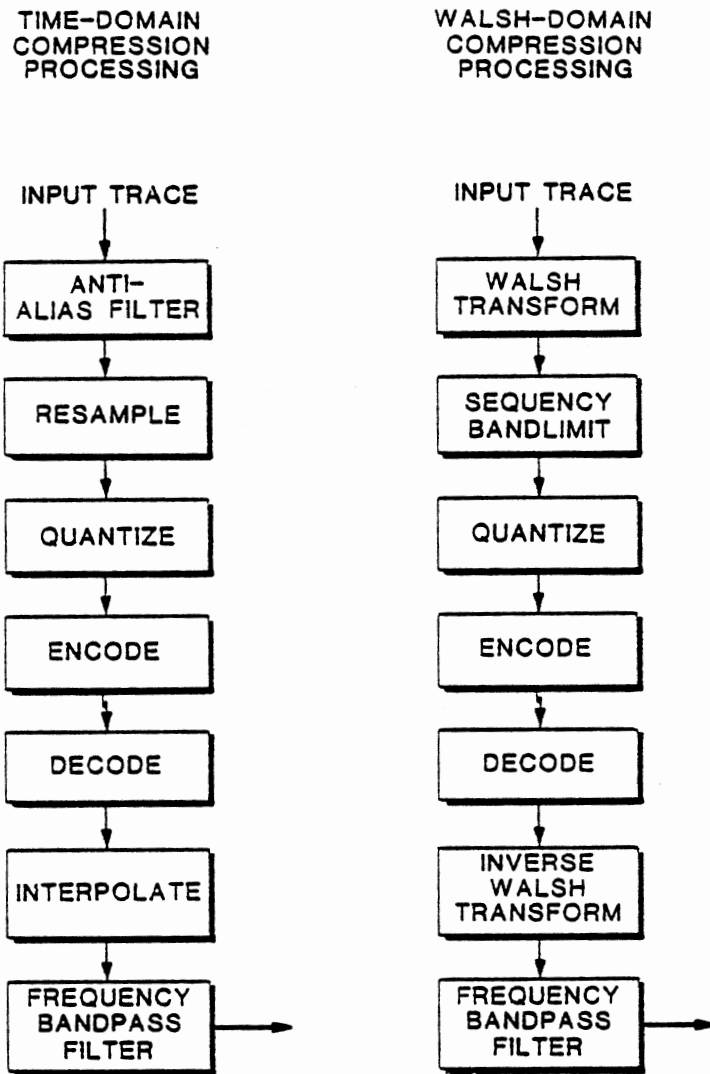


Figure 3. Data Compression Procedures
(After Wood [8])

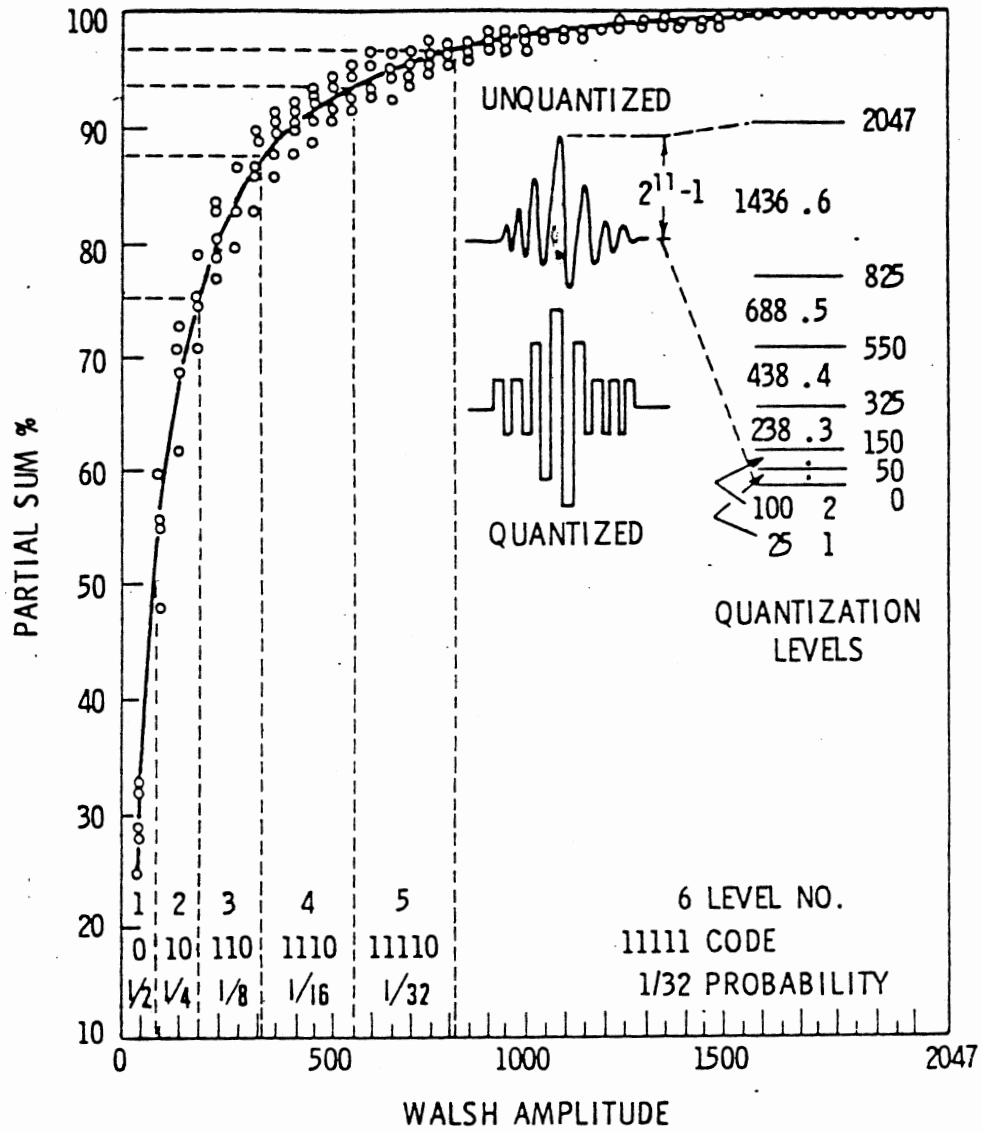


Figure 4. Non-Uniform Quantization (After Wood [8])

compression ratio is directly related to the sampling rate. This indicated that an 8:1 compression ratio is expected for the vibroseis data, because the sampling rate Wood used was 500 Hz instead of 250 Hz for the sample given in Appendix A. The proposed study examines alternative approaches for achieving seismic data compression with significant improvement in signal-to-noise ratio.

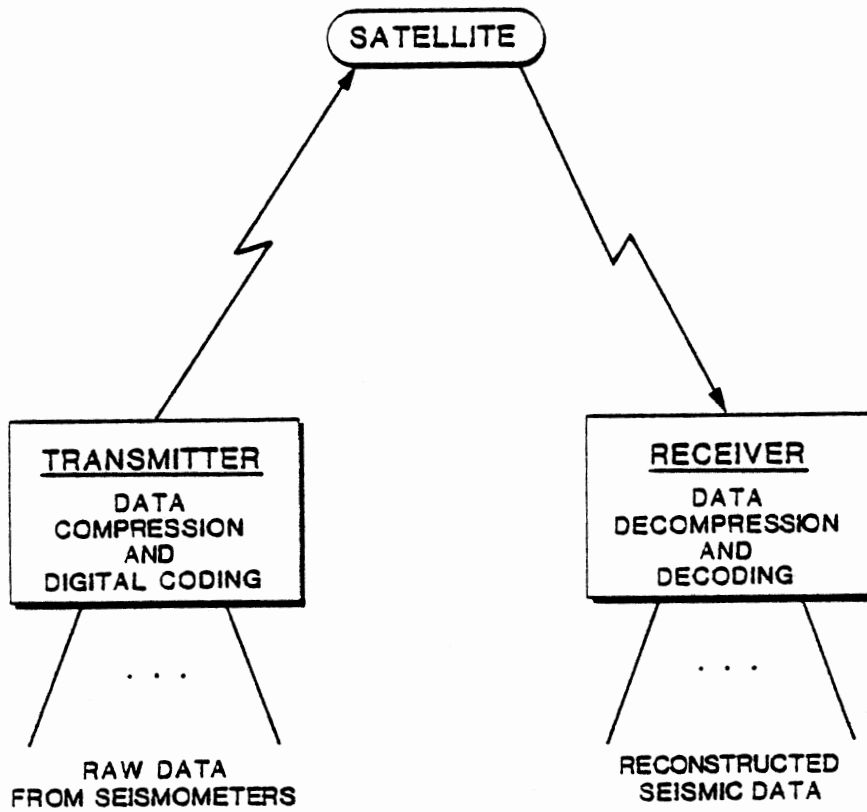
Proposed Approach

The proposed method of compression and decompression and satellite transmission is shown in Figure 5. The seismic data collected from the field are compressed and coded at the transmitter, and transmitted via satellite. The receiver performs decoding and decompression for reconstructing the original seismic data. The entire procedure involves data acquisition, data compression, satellite communication, and data reconstruction. This work will concentrate on the data compression and reconstruction for two types of seismic data, namely the vibroseis data and the impulsive seismic data. In Chapter V, the mechanics associated with satellite transmission and real-time implementation, such as time and space configurations and throughput considerations, will be discussed.

Before the compression methods are discussed, the statistical aspects of vibroseis data and impulsive seismic data need to be investigated. These are presented in the next section. The compression methods will be discussed in Chapters III and IV.

Statistical Analysis of the Seismic Data

It can be observed that signal processing techniques have been often developed based on a statistical model of signals [6]. For example, some



Available Baud Rate - 2400, 4800, 9600 and/or 56k bps
Reversible Data Compression - above 30 dB of SNR

Figure 5. Seismic Data Transmission via Satellite

digital coding methods for speech signals have been developed and evaluated under the assumption that the probability density function of speech signals can be modeled statistically as Laplace or Gamma probability densities. Similarly, it is necessary to obtain statistical models of seismic signals or to describe their statistical behaviors, so that an efficient compression technique can be chosen accordingly [10].

Statistical characteristics of seismic signals are examined with respect to several aspects. Such aspects would be statistical parameters, time-varying characteristics, the probability density functions, and correlation considerations in adjacent samples as well as in traces and tests. The term test indicates a source initiation in this thesis.

The statistical parameters include mean, variance, and peak-to-peak range. The examination of such parameters shows that the mean values are approximately zero for all traces and that the seismic data have nonuniform energy distribution among sensors. Amplitude distributions of several traces shown in Figure 6 illustrate symmetry of seismic data. Figures 7a and 7b give a comparison of the peak values among traces and among tests. Figures 8a and 8b similarly compare the ratios between the peak value and the standard deviation.

For testing the stationarity of the seismic data, the time-varying characteristics are analyzed. This analysis has been performed on the seismic data trace by trace by dividing the sampling duration into time unit segments. For vibroseis data, one second is used as a time unit while 100 milliseconds are used for impulsive seismic data. The results are obtained by examining the statistical parameters mentioned above for each time unit interval. Figures 9a and 9b illustrate the time-varying characteristics for the peak values and the ratios between the peak value

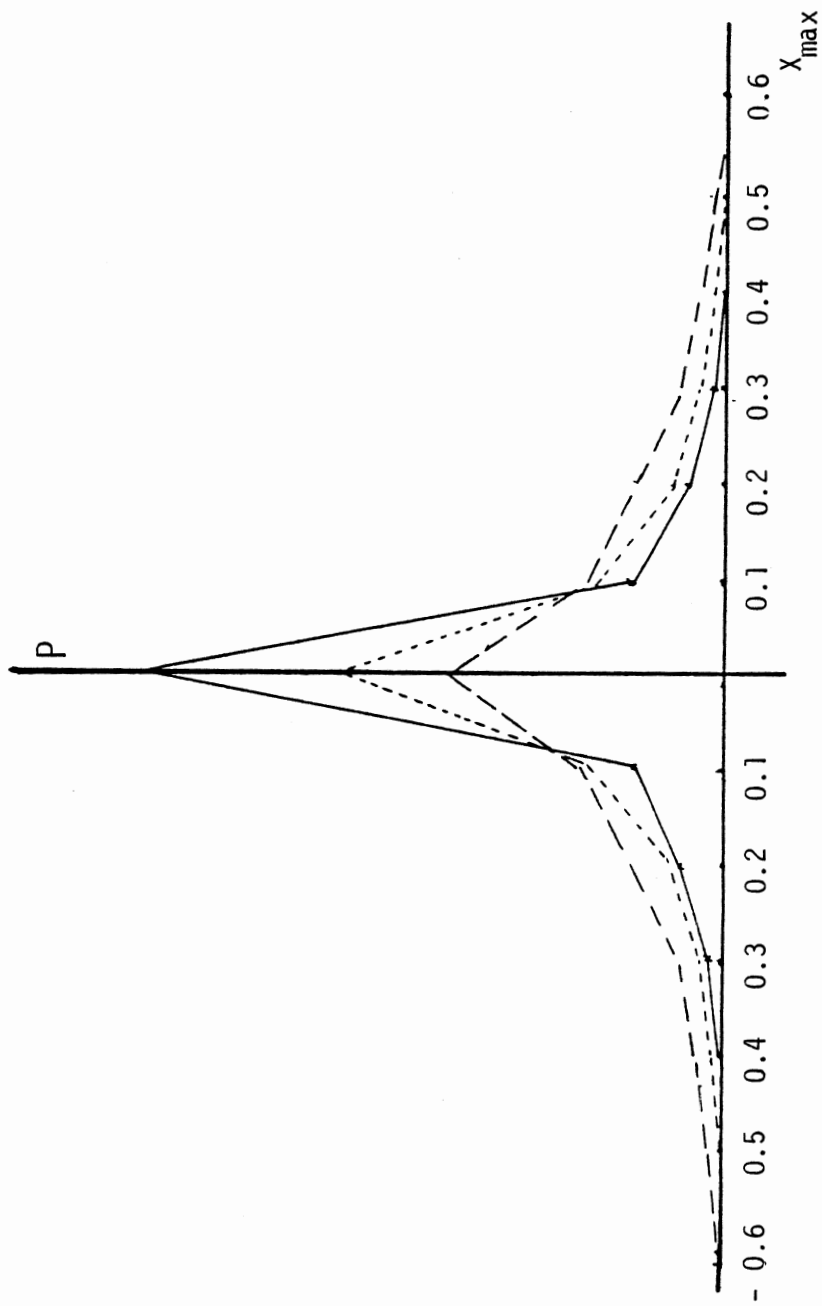
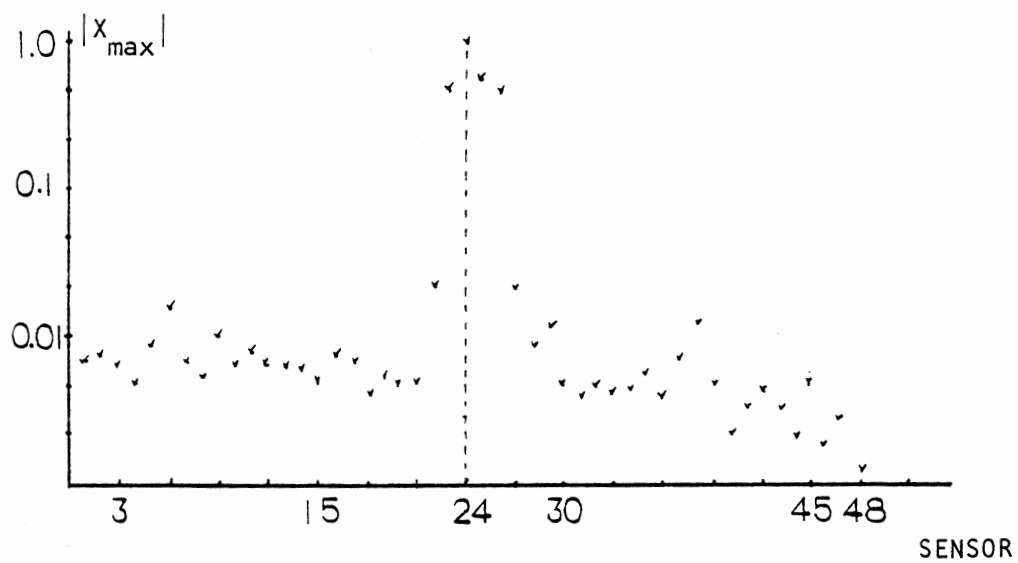
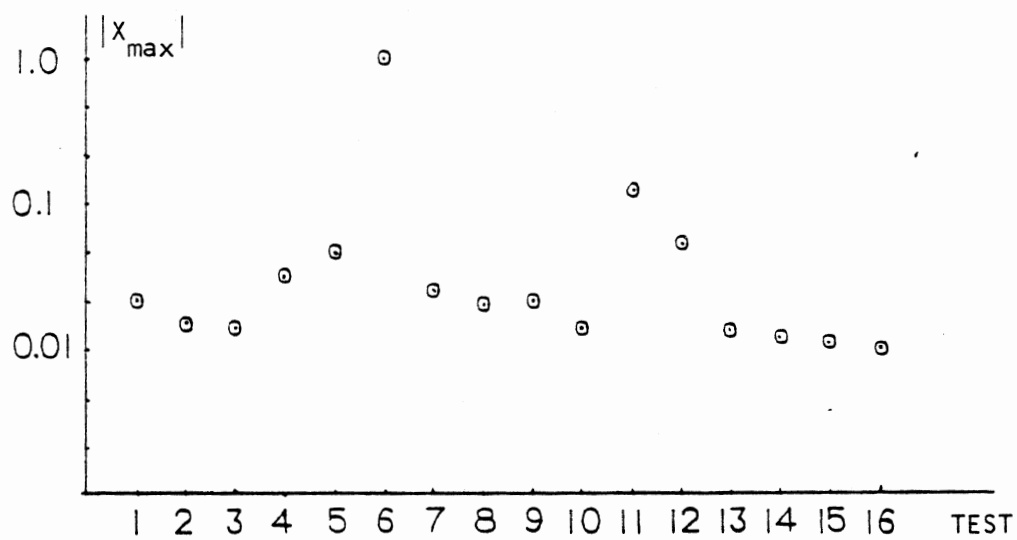


Figure 6. Amplitude Distribution of Vibroseis Data

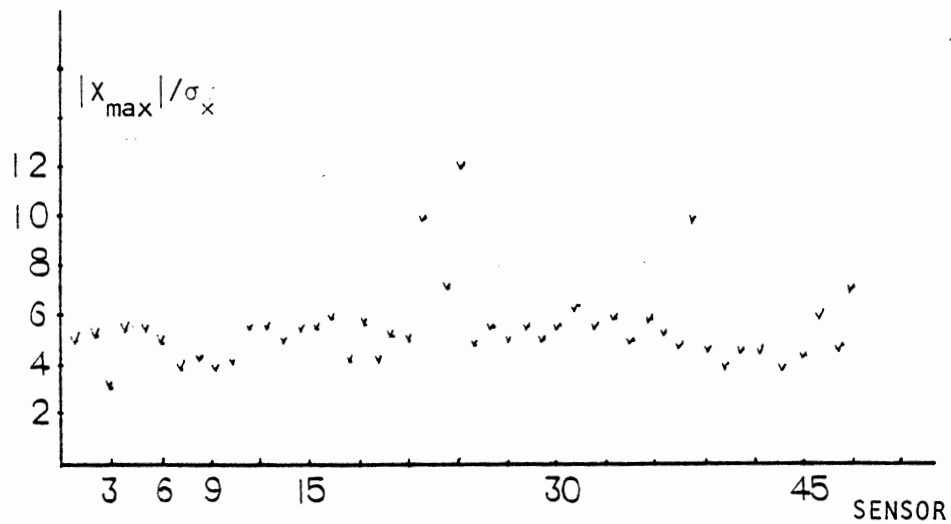


(a) Characteristics of $|X_{\max}|$ among sensors

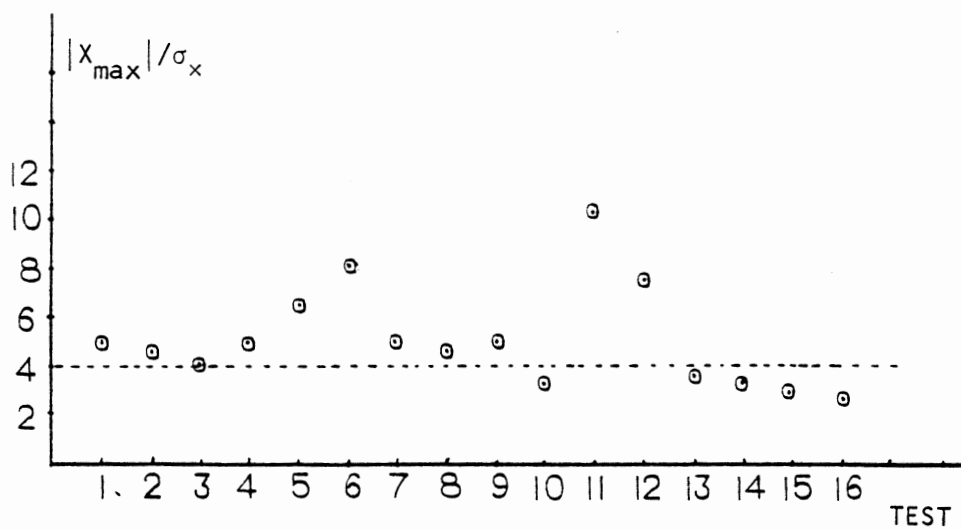


(b) Characteristics of $|X_{\max}|$ among tests for one sensor

Figure 7. Illustrations of $|X_{\max}|$ Characteristics of Vibroseis Data

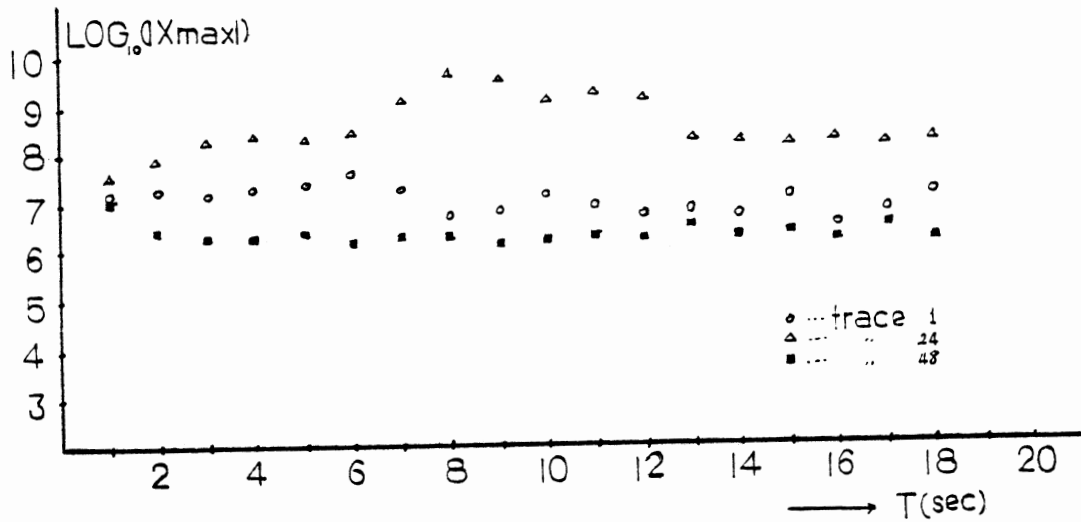


(a) Characteristics of $|X_{\max}|/\sigma_x$ among sensors

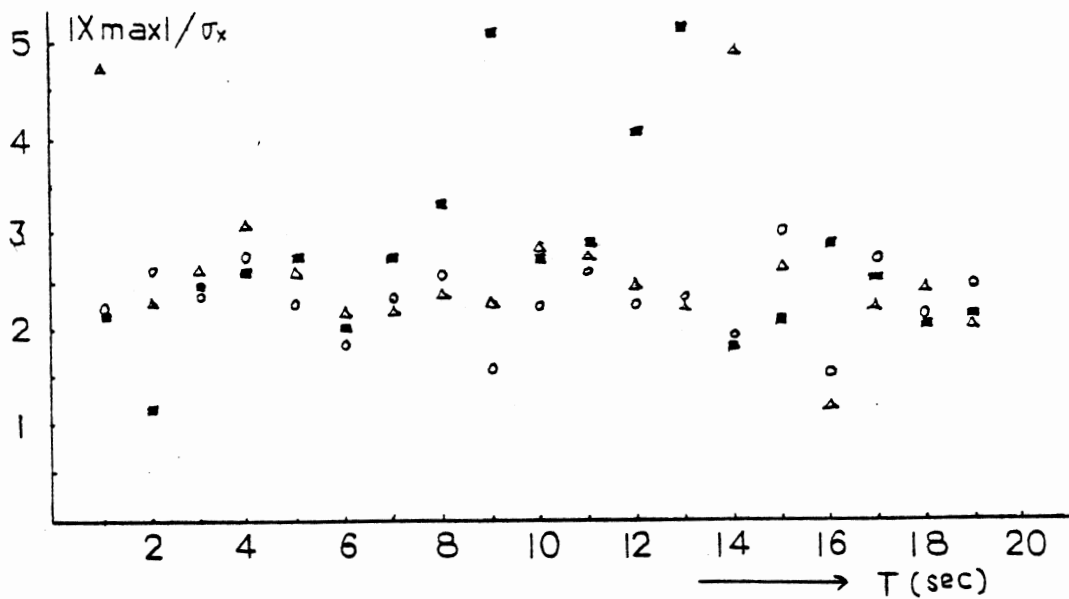


(b) Characteristics of $|X_{\max}|/\sigma_x$ among tests for one sensor

Figure 8. Illustrations of $|X_{\max}|/\sigma_x$ Characteristics of Vibroseis Data

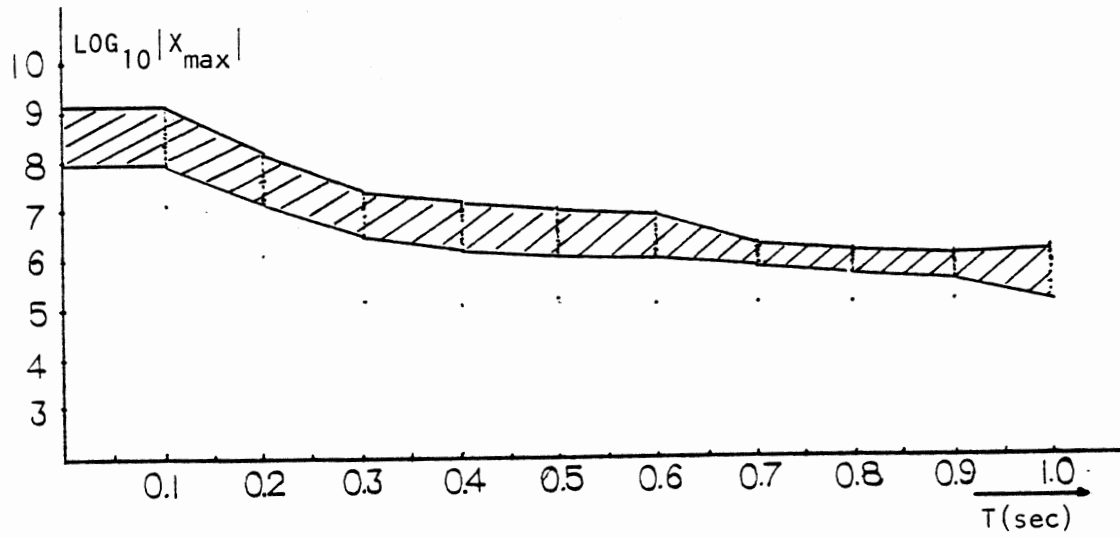


(a) Time Varying Characteristics of $|X_{\max}|$ of Vibroseis Data

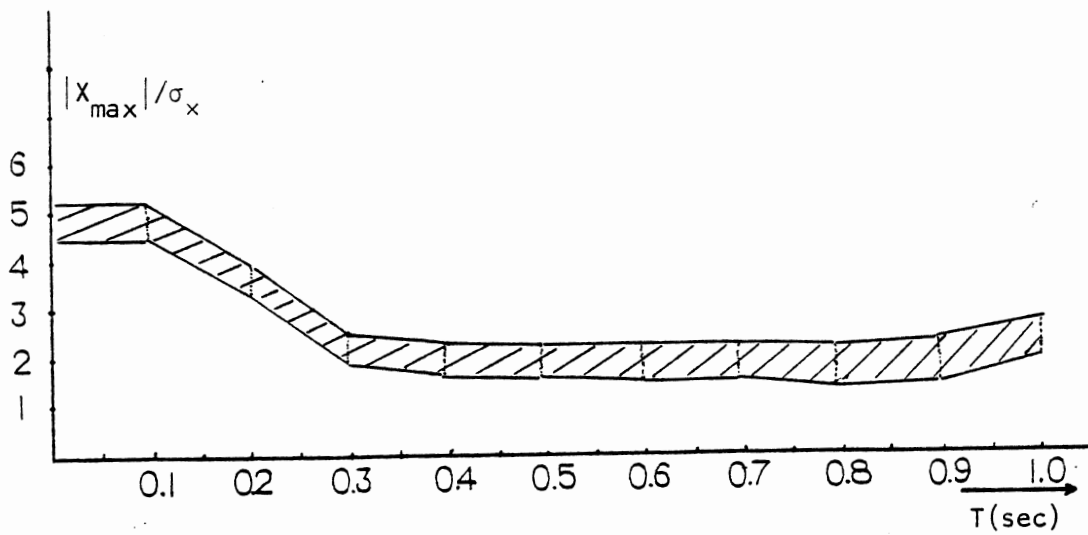


(b) Time Varying Characteristics of $|X_{\max}|/\sigma_x$ of Vibroseis Data

Figure 9. Illustrations of Time Varying Characteristics of Seismic Data



(c) Time Varying Characteristics of $|X_{\max}|$ of Impulsive Seismic Data



(d) Time Varying Characteristics of $|X_{\max}|/\sigma_x$ of Impulsive Seismic Data

Figure 9. (Continued)

and the standard deviation of vibroseis data and of impulsive seismic data, respectively. It should be pointed out that three traces are displayed individually for vibroseis data and the ranges of 100 traces are displayed for impulsive seismic data. This is based on the fact that vibroseis data vary nonuniformly among traces, while impulsive seismic data vary uniformly among traces.

The probability density functions are obtained using the frequency histogram method. The frequency indicates the number of occurrences of data in a referenced range. The comparison of seismic data with well-known probability density functions, Laplace, Gamma, and speech signals, was performed in order to obtain a statistical model for seismic data. The results are shown in Figure 10. It can be observed that vibroseis data have a similar probability density function as speech signals. This indicates that it may be possible to use the results from speech processing. These results may include some performance measures of various digital speech coding methods. It can also be observed that the Laplace or Gamma density functions generally define the seismic data in a statistical sense.

The correlation between adjacent samples can be examined from the value of the autocorrelation function at the first lag. For example, the autocorrelation function of a partial trace of vibroseis data is shown in Figure 11. It can be seen that there is a relatively high correlation among sample points.

Also, correlation characteristics among traces and among tests are examined via computing the correlation coefficients. These examinations are used for selecting traces or tests to be implemented for data compression techniques, so that statistically similar traces and tests can be

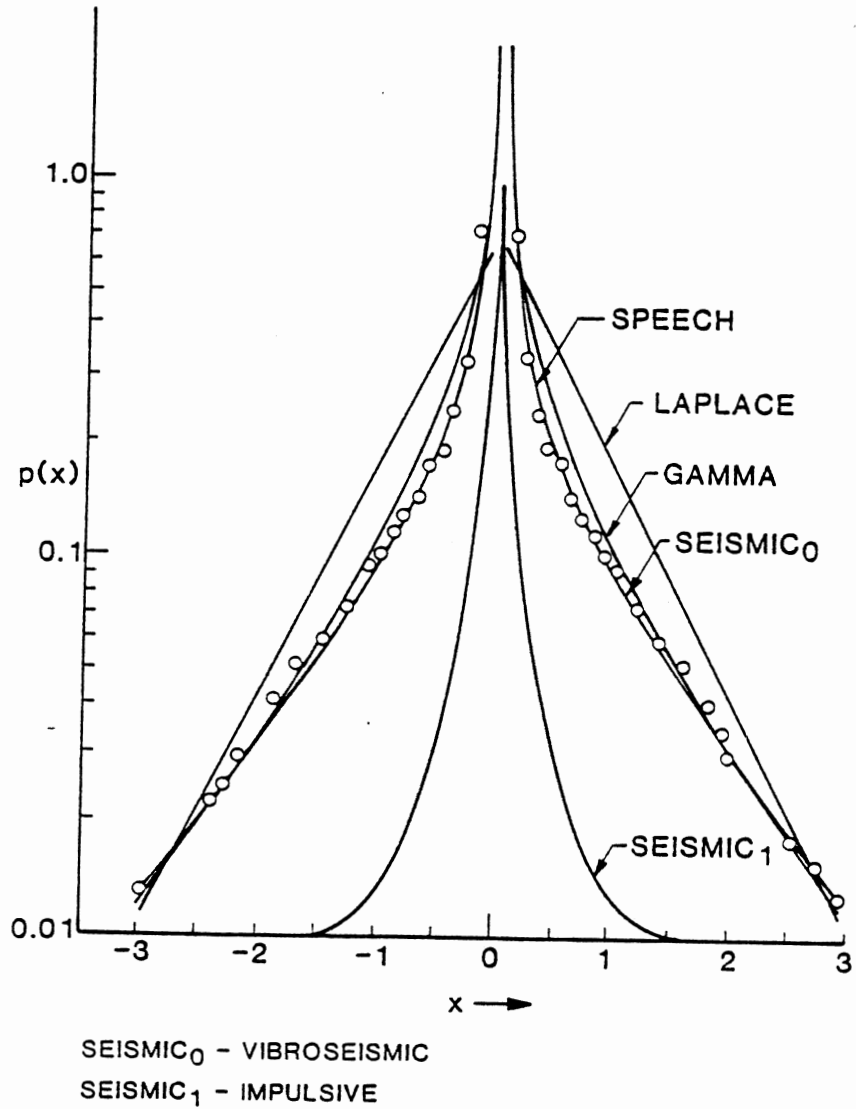


Figure 10. Probability Density Function of Seismic Data in Comparison of Laplace, Gamma and Speech (After Max and Paez [10])

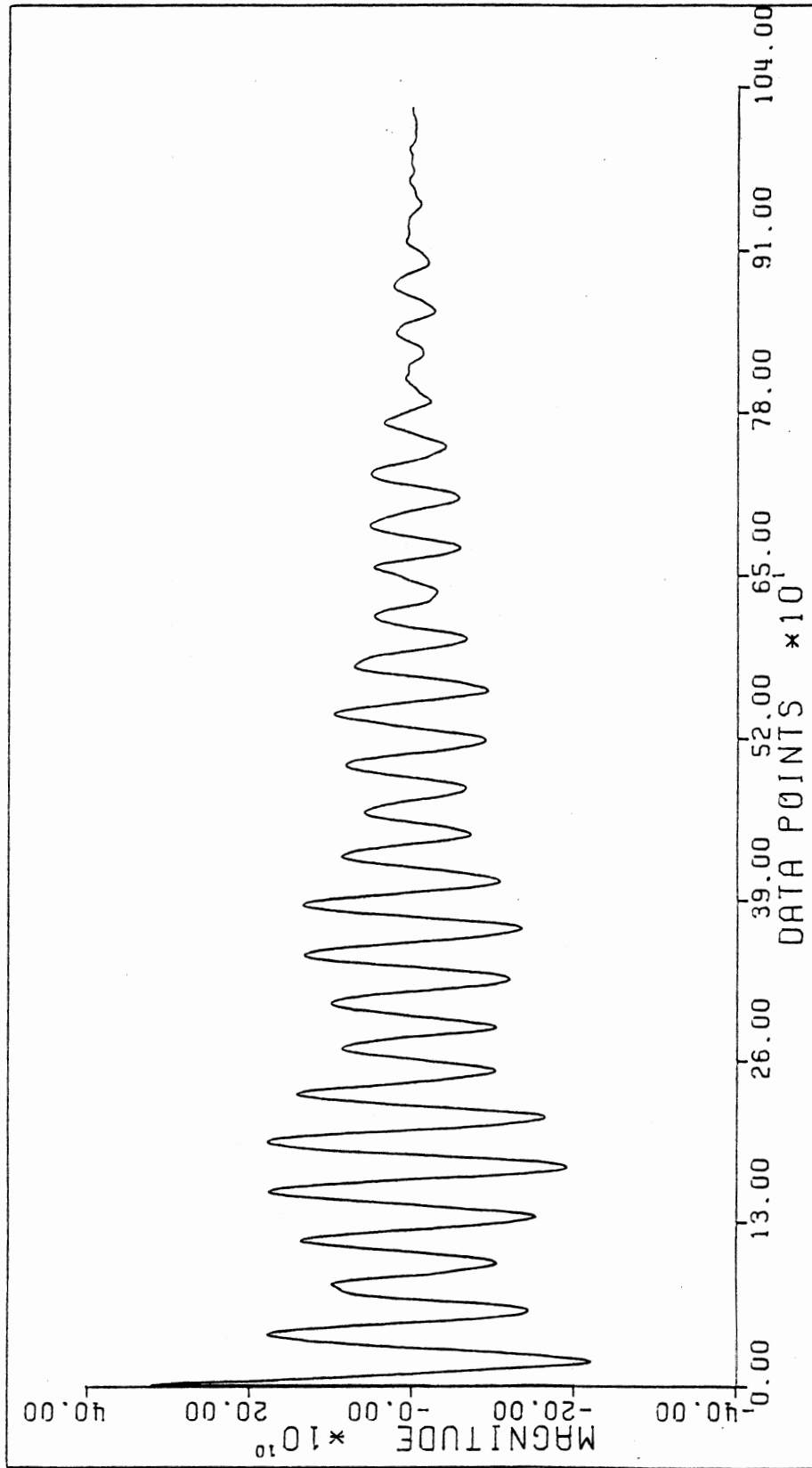


Figure 11. Auto-correlation Function of Vibroseis Data

grouped together. The correlation coefficients are ranged between +1 and -1, where +1 implies a complete linear relation and -1 implies a complete inversely linear relation [10]. Appendix B illustrates these coefficients of 16 traces of one sensor.

Further statistical analysis has been performed by applying the pre-emphasis method with low frequency and high frequency. The low-frequency pre-emphasized signal (D_n) can be computed from

$$D_n = X_n - a D_{n-1} \quad (2.3)$$

where $D_0 = X_0$ for $n = 1, 2, \dots, N-1$, X_n is the original signal, and a is a gain factor. The high-frequency pre-emphasized signal (H_n) can be computed from

$$H_n = X_n - b X_{n-1} \quad (2.4)$$

where $H_0 = X_0$ for $n = 1, 2, \dots, N-1$, and b is a gain factor.

The above pre-emphasis methods will be referred to as the first-order pre-emphasis. Also, the second-order pre-emphasis methods are investigated by applying the pre-emphasis on previously emphasized signals. For these first- and second-order pre-emphasized signals, the statistical characteristics mentioned above were examined. The important results obtained from these investigations are that the correlation coefficients increase (decrease) for low-frequency (high-frequency) pre-emphasized signals. In particular, for the second-order low-frequency pre-emphasized signals, the correlation coefficients are almost 1.0, as shown in Appendix B. The correlation considerations are not discussed for impulsive seismic data, since it is clearly seen from the plot in Appendix A that all of the traces are highly correlated.

The seismic data have been statistically analyzed in this section. The analysis results will be utilized for evaluating the utility of various data compression techniques for seismic data in Chapter III.

Summary

In this chapter, general considerations involved in seismic data compression, such as seismic data rate and size, reversibility of compression techniques, and statistical characteristics of seismic data, have been discussed. Data rate and size of vibroseis data were found to be 192k bps and 9M byte per block. Impulsive seismic data were shown to have 3960k bps of data rate and 230k byte of data size per block.

Due to bandwidth constraints of telecommunication channels and computer storage limits, data compression techniques are found to be necessary for seismic data transmission. One previous work by Wood [8], which is basically sequency limiting in the Walsh domain, was examined and it was determined to be improper since its signal-to-noise ratio is too small (less than 7 dB). For maintaining the desired quality of signal after reconstruction at the receiver, a SNR of 30 dB was considered essential for a compression method to be useful.

Analysis of statistical characteristics of vibroseis data showed that vibroseis data have a similar probability distributional characteristic as speech signals. Also, it was observed that vibroseis data traces are slightly correlated among each other, while impulsive seismic data traces are highly correlated. These characteristics will be used for evaluating performances of various compression techniques in the next chapter.

CHAPTER III

EVALUATION OF DATA COMPRESSION TECHNIQUES

Introduction

The statistical characteristics of vibroseis and impulsive seismic data were examined in Chapter II. Based on these characteristics, this chapter is devoted to evaluating techniques for reversible seismic data compression. The reversibility constraint is defined as 30 dB of SNR in this thesis, assuming that this constraint may provide the required quality for the signal to be reconstructed at the receiver.

In this chapter, five basic data compression techniques mentioned in Chapter I are discussed with respect to their use for reversible seismic data compression. First, it can be argued that the entropy reducing transforms can be neglected, since these methods extract the desired information from a portion of the data in terms of time and frequency, and thus lose the reversibility.

Second, prediction techniques are considered with respect to identifying an optimal predictor where the "linear prediction" algorithm is used for obtaining parameters for the predictor. For vibroseis data, these methods are ignored due to the computational complexities involved in obtaining the predictor parameters. However, these methods are implementable for impulsive seismic data since the predictor parameters can be obtained in a simple manner; these parameters can be used for designing a predictor for other traces. This is based on the characteristics of

waveforms of the impulsive seismic data, which will be discussed later.

Third, interpolation techniques are discussed in relation to resampling techniques where some sample values can be reconstructed using a polynomial interpolator. The parameters associated with a polynomial interpolator should be generated at the transmitter and transmitted in place of the original data. This technique involves basically similar computational complexities as prediction techniques. For a seismic signal, the number of parameters associated with a polynomial interpolator is excessive and for these reasons, interpolation techniques are ignored.

Orthogonal transforms can be chosen for vibroseis data and for a nonimpulsive section of impulsive seismic data. These transforms remove redundancies and also they are simple implementation-wise. Energy conservation and the invertibility property of the orthogonal transforms allow for using other compression techniques in the transformed domain. Similarly, other compression techniques can be used prior to orthogonal transforms. These aspects will be examined in more detail in the next section.

Digital coding methods can be proposed for two reasons. First, digital coding is required for transmission over a digital communication channel. Second, data compression can be achieved by applying an efficient digital coding technique. Data compression consideration via digital coding is based on reducing the average number of bits (or bit rate) per symbol rather than reducing the number of messages. An efficient digital coding technique can be approached by various quantization techniques and encoding techniques. These various techniques will be discussed in a later part of this chapter.

Finally, the compression results of various compression techniques are compared and combinations of these techniques are proposed for vibroseis data and impulsive seismic data.

Orthogonal Transforms in Data Compression

Orthogonal transforms can be used for redundancy removal. This property is one of the key elements of data compression. The redundancy removal property is discussed below [5].

Let a sequence of N data points be represented by an N -dimensional vector \underline{X} . A transformed vector \underline{Y} can be formed from \underline{X} by

$$\underline{Y} = A\underline{X} , \quad (3.1)$$

where A is a unitary matrix. That is,

$$A^* A = I . \quad (3.2)$$

where $*$ indicates a complex conjugate transpose. The objective is to select a subset of M components of \underline{Y} , where M is substantially less than N . The remaining $(N - M)$ components can then be discarded without introducing objectionable error, where the error is, of course, the difference between the signal and the reconstructed signal using the retained M components of \underline{Y} . The error criterion often used for orthogonal transforms is the mean-square error criterion, and is discussed below.

From Equations (3.1) and (3.2),

$$\underline{X} = A^* \underline{Y} = \sum_{i=1}^N y_i \phi_i . \quad (3.3)$$

where

$$A^* = [\phi_1, \phi_2, \dots, \phi_N] .$$

The reconstructed vector \underline{X}' from the retained M components of \underline{Y} can be given by

$$\underline{X}' = A^* \underline{Y}' , \quad (3.4)$$

where \underline{Y}' contains M components of \underline{Y} and a constant replaced for the discarded N-M components. It should be noted that M largest eigenvalues should be selected for the subset of \underline{Y}' . The mean square error is usually defined as [5]

$$e = E[(X - X')^* (X - X')] = \sum_{i=M+1}^N \phi_i^* \Sigma_x \phi_i \quad (3.5)$$

where

$$\Sigma_x = E[(X - \bar{X})(X - \bar{X})^*]$$

with \bar{X} being the mean value of \underline{X} . From the above relations, the minimum error is given by

$$e_{\min} = \sum_{i=M+1}^N \lambda_i$$

with ϕ_i and λ_i being the eigenvector and the corresponding eigenvalue of the covariance matrix Σ_x .

The unitary transform A, composed of the eigenvectors of the covariance matrix of the given data, is called the Karhunen-Loève transform. Also, it can be seen that the covariance matrix of the transformed vector \underline{Y} is uncorrelated and it is expressed as [5]

$$\Sigma_y = \text{diag} (\lambda_1, \lambda_2, \dots, \lambda_N) . \quad (3.7)$$

Equation (3.7) indicates that \underline{Y} has no redundancy; thus \underline{Y} is the most efficient representation of \underline{X} . However, there exists no general fast algorithm to compute the KLT, since the KLT depends upon the data

covariance matrix. The computation involves $2N^2$ multiplications, and as N increases, the task of computing the transformed vector becomes a formidable one [11]. For this reason, several suboptimal orthogonal transforms are investigated, which include the Walsh-Hadamard, the discrete Fourier transform, and the discrete cosine transform. For illustrative purposes, vibroseis data shown in Figure 12 are used to compare these transforms.

Walsh-Hadamard Transform (WHT)

The matrix A in Equation (3.1), in terms of the Hadamard matrices, is defined by [11]

$$A = A_{\text{WHT}} = H(\nu) \quad (3.8a)$$

$$H(\nu) = \begin{bmatrix} H(\nu-1) & H(\nu-1) \\ H(\nu-1) & -H(\nu-1) \end{bmatrix} \quad (3.8b)$$

where

$$H(0) = 1, \quad N = 2^\nu$$

and

$$A_{\text{WHT}}^{-1} = \frac{1}{N} H_\nu.$$

Noting that $H(\nu)$ has only ± 1 's, the WHT algorithm requires only $N \log_2 N$ summations [11]. This computational simplicity is the main reason for its wide usage.

Figure 13 illustrates the Walsh-Hadamard transformed vector of the vibroseis data shown in Figure 12. As will be shown later, the WHT compression ratio is not as good as some of the other suboptimal transforms and, therefore, WHT is not used in the proposed compression methods.

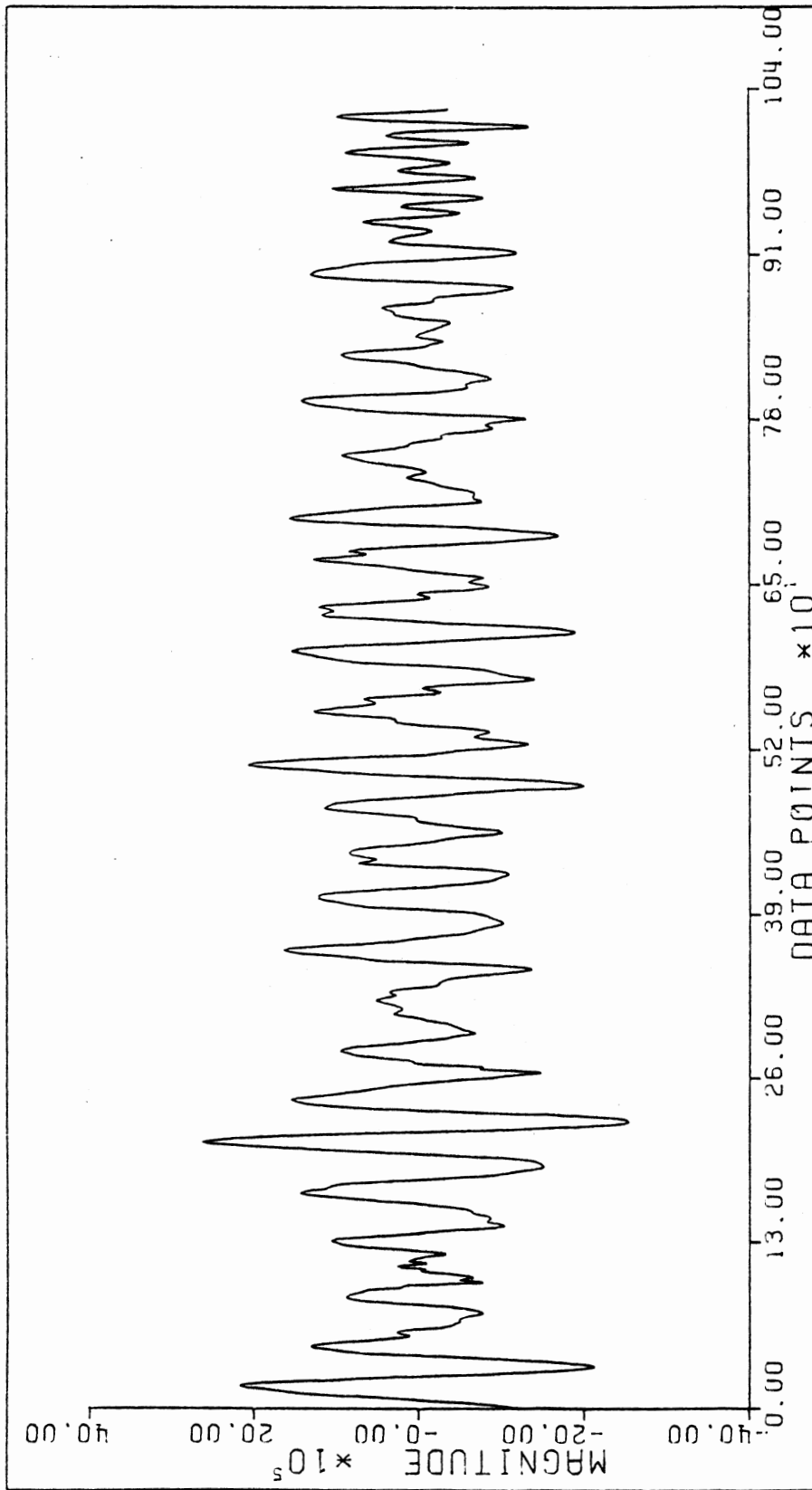


Figure 12. Partial Trace of Vibroseis Data

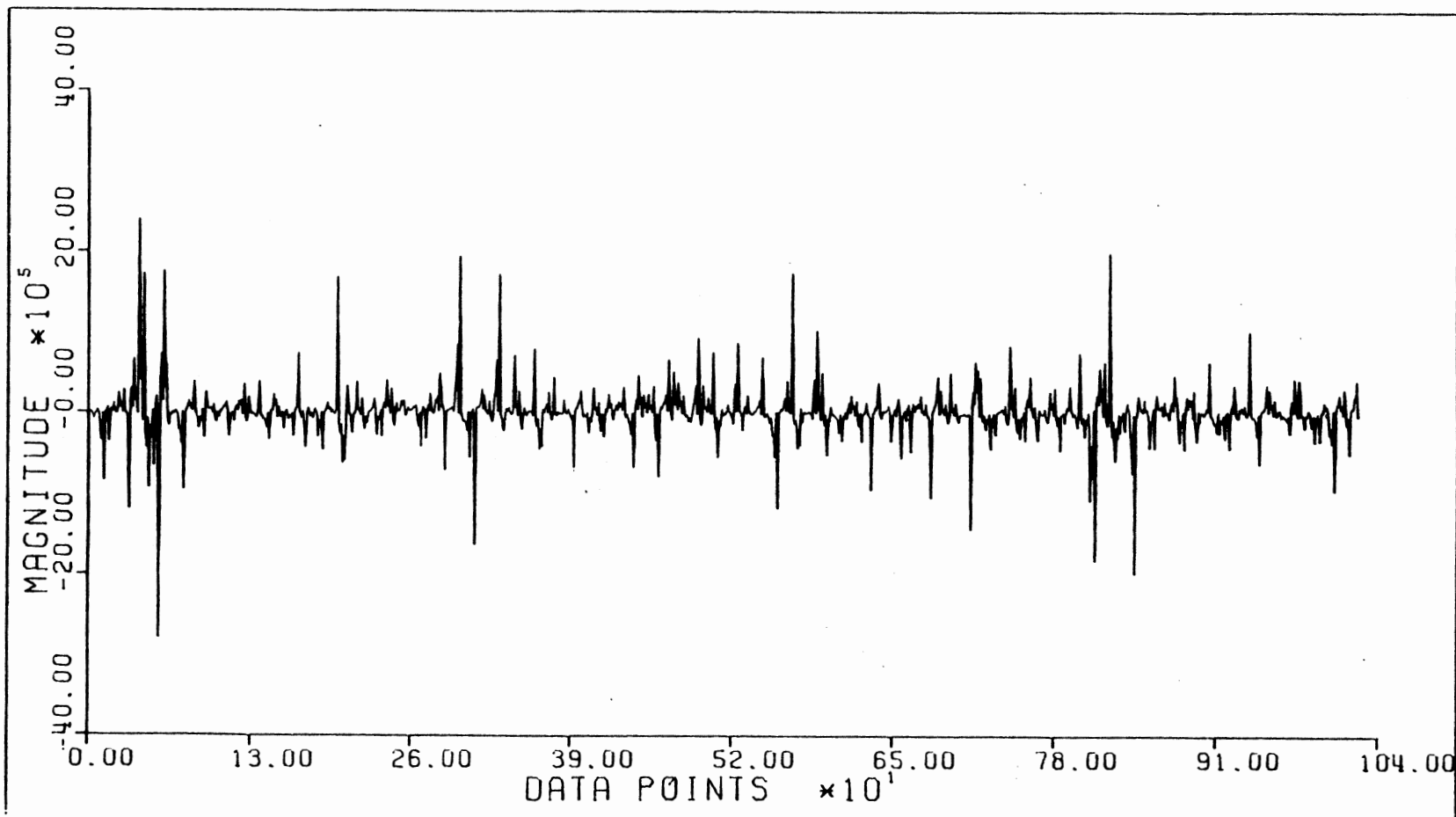


Figure 13. Walsh-Hadamard Transformed Vector Coefficients

Discrete Fourier Transform (DFT)

The matrix A in Equation (3.1), corresponding to the DFT, is defined by [5]

$$A = A_{\text{DFT}} = (a_{kj}) \quad (3.9)$$

where

$$a_{kj} = \exp[-i2\pi \frac{kj}{N}], \quad 0 \leq k, j \leq N - 1$$

and

$$A_{\text{DFT}}^{-1} = \frac{1}{N} A^*$$

where A^* is the complex conjugate transpose of A . Note that $(1/\sqrt{N})$ can be incorporated into A to make it unitary and, of course, it has no important effect on the nature of the representation.

The DFT is of interest primarily because it approximates the continuous Fourier transform and fast algorithms are available. The basic FFT algorithm requires $2N \log_2 N$ multiplications for $N = 2^v$ to compute the transformed vector, and it is significantly simpler than the KLT computation [5].

Figure 14 shows the DFT transformed vector of the vibroseis data shown in Figure 12. The first half of the plot is the real part of the transformed vector components and the second half is the imaginary part of the transformed vector components. The DFT requires a complex components array and two threshold values are necessary for selecting the significant coefficients. Compared to the discrete cosine transform below, the DFT is not attractive.

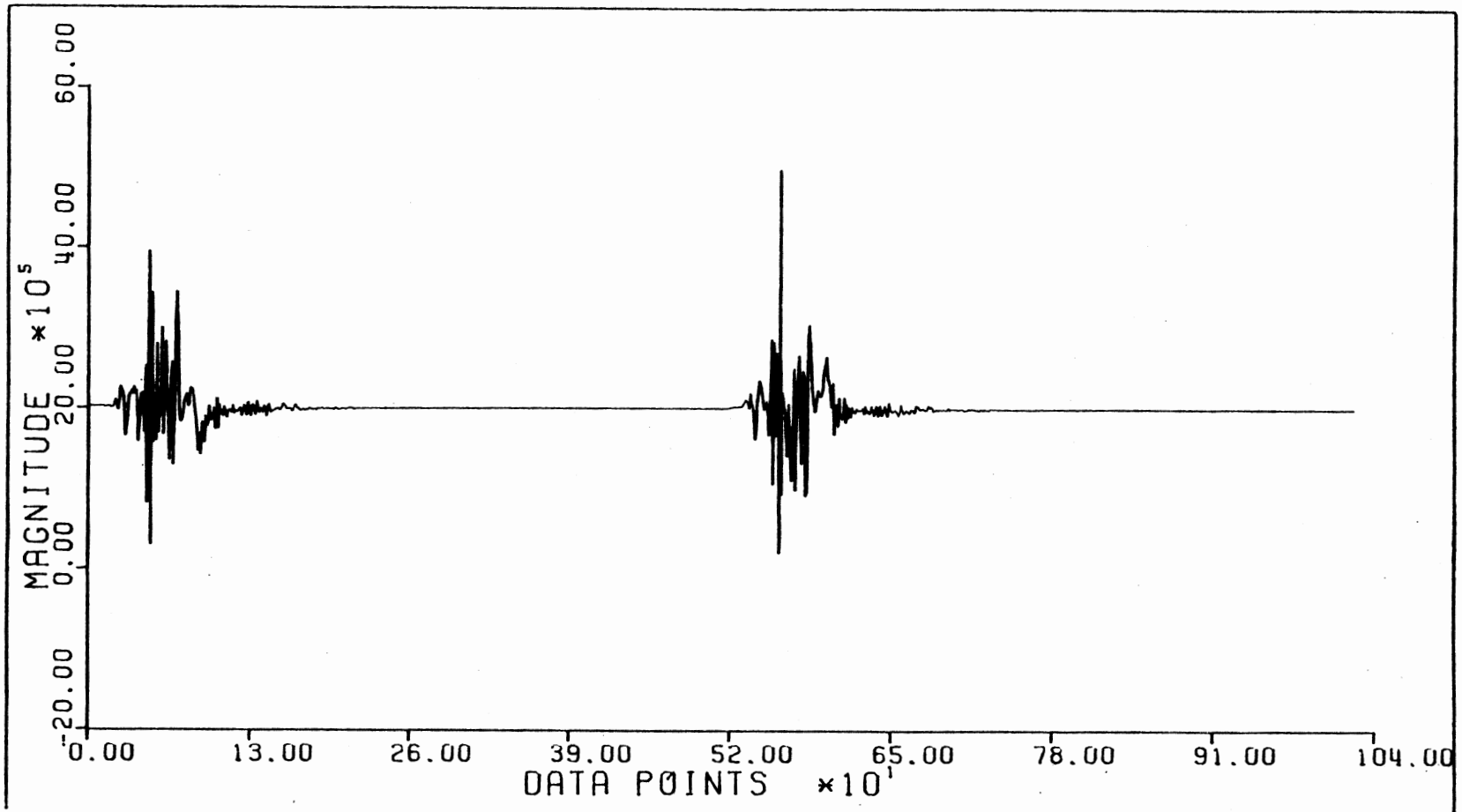


Figure 14. Discrete Fourier Transformed Vector Coefficients

Discrete Cosine Transform (DCT)

The DCT is similar to the DFT in that it uses sinusoidal waveforms as the basis of its orthogonal transform matrix. The matrix A in Equation (3.1) corresponding to the DCT is defined by [5]

$$A = A_{\text{DCT}} = (a_{kj}) \quad (3.10)$$

where

$$a_{kj} = \frac{2}{N} \cos \frac{2j+1}{2N} k\pi \quad 0 \leq j \leq N-1; 1 \leq k \leq N-1$$

and

$$a_{0j} = \frac{1}{N} \quad 0 \leq j \leq N-1$$

The inverse of A_{DCT} is given by

$$A_{\text{DCT}}^{-1} = (b_{jk})$$

where

$$b_{jk} = \frac{2}{N} \cos \left(\frac{2j+1}{2N} k\pi \right) + \frac{1}{N} \quad 1 \leq k \leq N-1; 0 \leq j \leq N-1$$

and

$$b_{k0} = \frac{1}{N} \quad 0 \leq k \leq N-1$$

The DCT can be computed using the FFT as follows [13]:

$$a_{kj} = \frac{2}{N} \operatorname{Re} \left[\exp \frac{2j+1}{2N} k\pi \right]$$

Similarly, it can be shown that the algorithm can be used to compute the inverse DCT coefficients. Also, it has been shown that the basis vectors

of the DCT closely approximate the eigenvectors of a class of Toeplitz matrices and that the DCT approaches the KLT as far as optimality is concerned [11]. For these reasons, the DCT has been chosen as the orthogonal transformation technique for the vibroseis data compression.

Figure 15 shows the DCT transformed vector of the vibroseis data shown in Figure 13. For the three suboptimal transforms discussed above, the compression ratios (CPR) are tabulated in Table I corresponding to the vibroseis data in Figure 12. From these results it can be seen that the DCT gives the best compression among the three suboptimal transforms.

Digital Coding Techniques

It was pointed out earlier that efficient step-size to reduce the number of quantization levels and efficient bit allocation techniques to minimize the number of bits play an important role in digital coding for data compression. These ideas are implemented in various forms of quantization and encoding; some of these ideas are discussed below.

At the transmitter, the source signal is coded using the digital coding method, and the source signal is reconstructed at the receiver from the coded signal, subject to some error fidelity criteria. This is shown in Figure 16, where the input signal is denoted by $X(n)$, the quantized signal is denoted by $\hat{X}(n)$, and $C(n)$ is the code word for $\hat{X}(n)$. The received code $C'(n)$ is decoded and the decoded signal is denoted by $X'(n)$. Transmission error (or channel error) is not considered in this thesis for evaluating the compressed results of various techniques. Therefore, it is assumed that $C(n) = C'(n)$ and thus $\hat{X} = \hat{X}'(n)$. This implies the error $(X - X')$ is mainly quantization error.

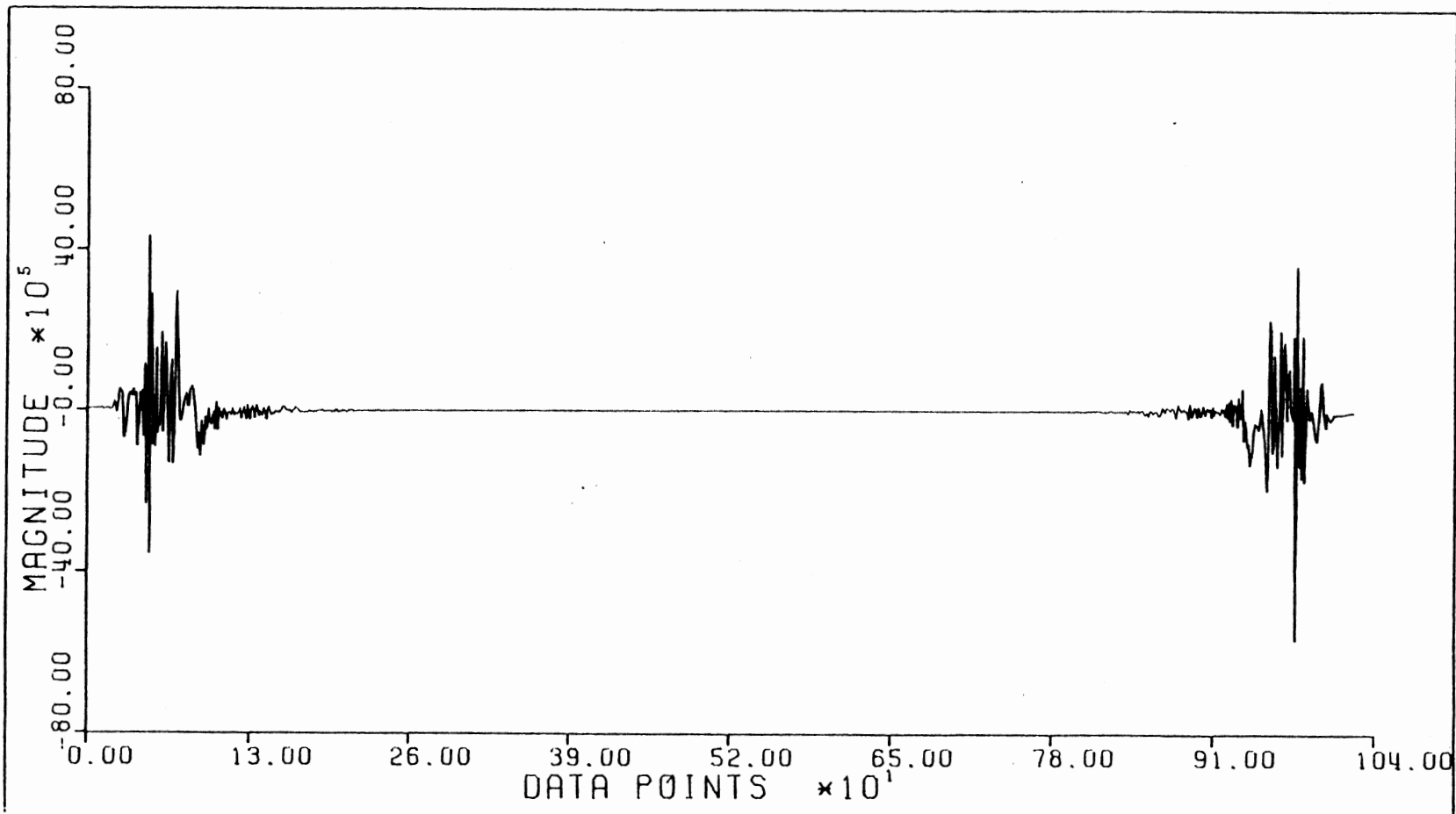
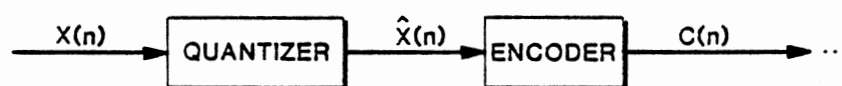


Figure 15. Discrete Cosine Transformed Vector Coefficients

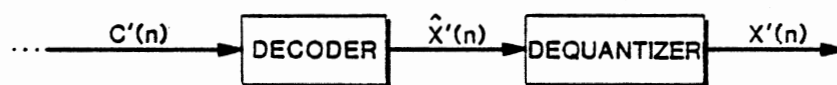
TABLE I
 CPR* OF ORTHOGONAL TRANSFORMS
 ON VIBROSEIS DATA

TEST	WHT	DFT	DCT
1	1.19	1.44	2.41
2	1.17	1.24	1.70
3	1.21	2.74	3.02
4	1.21	1.17	1.40
5	1.21	2.59	2.96
6	1.19	1.37	2.17
7	1.19	2.99	3.32
8	1.15	2.68	2.88
9	1.10	1.82	2.73
10	1.16	1.20	1.35
11	1.16	1.30	1.66
12	1.15	1.22	1.62
13	1.14	2.24	2.50
14	1.12	2.73	2.88
15	1.13	1.27	1.60
16	1.21	2.31	2.52
AVG	1.17	1.89	2.29

*Threshold = $0.05\sigma_x$
 SNR > 35 dB



(a) TRANSMITTER



(b) RECEIVER

Figure 16. Digital Coding Techniques

For a given source and a given error fidelity criterion, the minimum transmission rate can be computed from the rate distortion function of Shannon [14], which will be discussed later. This optimum rate cannot readily be achieved because the coding technique is usually extremely complex or theoretically intractable. Most data compression techniques are suboptimal in the sense that they exceed the minimum possible transmission rate. It is not possible to choose the "best" way of coding for a given application, as the computational complexities and hardware design play important roles. The decision must be based on some vague factors, such as generality of the method relative to the source information and relative equipment complexities.

The following five sections are devoted to examining various coding techniques, which involve quantization and encoding for seismic data compression. Also, comparative merits and demerits are discussed in terms of quantization noise and compression ratio. The quantization noise or signal-to-noise ratio considerations are examined with four basic step-size algorithms: optimum, uniform, logarithmic, and adaptive algorithms. The compression ratio considerations are examined with two algorithms, fixed code word length and optimum code word length. The differential coding techniques are discussed with respect to their contributions on signal-to-noise ratio. First, the optimum quantizer is examined in the next section.

Optimum Step-Size Quantization

The optimum quantization technique is discussed to indicate the computational complexities and to obtain a standard result for comparing other quantization techniques.

The variance of the quantization error is given by

$$\sigma_e^2 = \int e^2 p_e(e) de \quad (3.11)$$

where e is the quantization error $(x - \hat{x})$, and $P_e(e)$ is the probability density function of the error signal and can be expressed in terms of the probability density function of $X(n)$, $P_x(x)$ [10]. Equation (3.11) can thus be expressed as

$$\sigma_e^2 = 2 \sum_{i=1}^{M/2} \int_{x_{i-1}}^{x_i} (\hat{x}_i - x)^2 P_x(x) dx \quad (3.12)$$

where M is the number of quantization levels, and $P_x(x)$ is assumed to be equal to $P_x(-x)$. Equation (3.12) indicates that it is possible to choose the quantization levels so as to minimize the quantization error variance, and thus maximize the SNR, when $P_x(x)$ is known. A brief discussion of selecting $\{x_i\}$ and $\{\hat{x}_i\}$ which minimize σ_e^2 is given below.

By using the minimization process, the optimum location of the quantization level \hat{x}_i can be shown to be the centroid of the probability density interval x_{i-1} to x_i . Also, it has been found that the optimum boundary points lie halfway between the $M/2$ quantizer levels, \hat{x}_i . These nonlinearly related conditions must be met simultaneously, and iterative procedures are generally used to solve this problem. The parameter here is the step size.

As mentioned in Chapter 11, the vibroseis data have a probability density function close to Laplace and Gamma density functions. For this reason, the optimal quantizer for signals with Laplace density and Gamma density developed by Max, Paez and Glisson [10] are examined and the optimum quantization step size for Laplace density is shown in Figure 17.

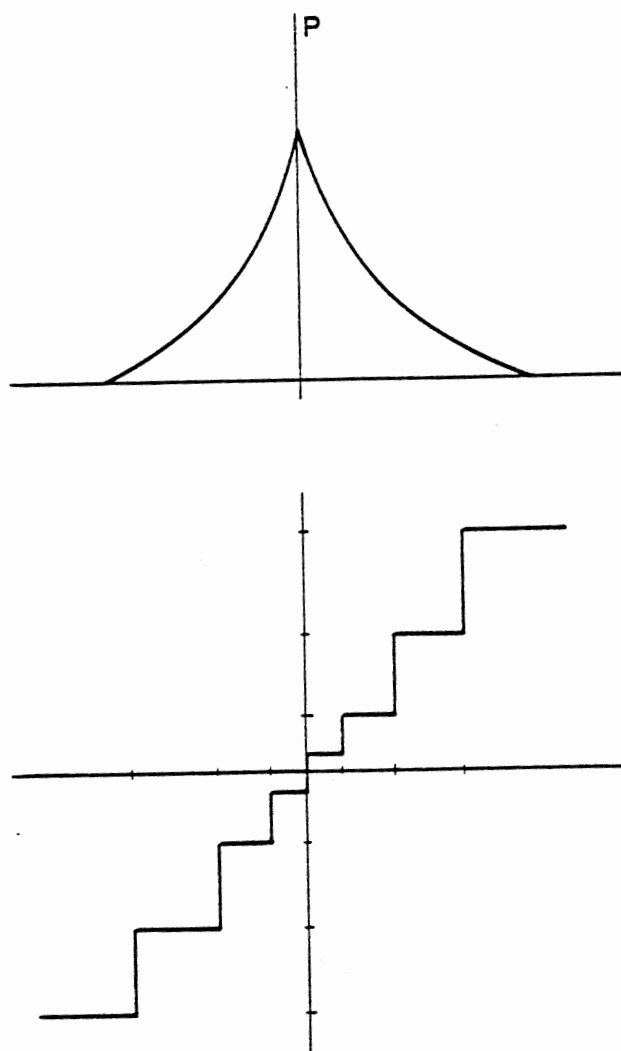


Figure 17. 3-bit Optimum Quantization
for Laplace Density
(After Max and Paez [10])

It can be easily seen from this figure that the quantization levels get farther apart as the probability density decreases. This indicates that large quantization errors should be reserved for the least frequently occurring samples.

This technique will not be used due to the computational complexity and the lack of exact information on the $P_x(x)$ for the vibroseis data. For these reasons, several suboptimum quantization techniques are discussed, including the uniform step-size quantization, non-uniform step-size, and adaptive step-size quantization.

Uniform Step-Size Quantization

The uniform step-size quantizer is the simplest kind of all the quantizers. This quantizer involves only two parameters, the number of levels and the quantization step-size, denoted as Δ . For a b -bit uniform quantizer, there are 2^b levels. This is illustrated in Figure 13 for a 3-bit quantizer. It has been found that the quantization error of this technique approaches the optimum quantizer when the signal is described by Gaussian distribution [15].

For future use, the parameter Δ and the SNR for this method is given below. The step-size is

$$\Delta = \frac{\text{Peak-to-peak range}}{\text{Number of levels}} .$$

If a symmetrical probability density function can be assumed for $X(n)$, the Δ can be expressed as

$$\Delta = \frac{2 |X_{\max}|}{2^b}$$

where $|X_{\max}|$ is the absolute maximum of $X(n)$.

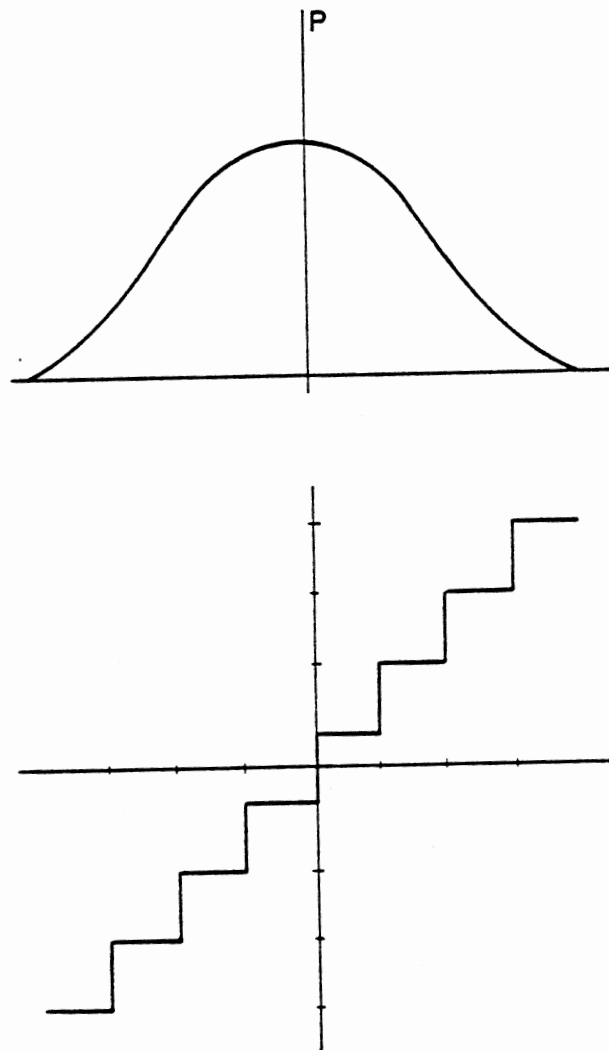


Figure 18. 3-bit Uniform Quantization
for Gauss Density

The quantization error of the uniform step-size quantizer [16]

$$\text{SNR(dB)} = 6b + 4.77 - 20 \log (|X_{\max}|/\sigma_x) \quad (3.13)$$

which points out that each additional bit contributes a 6 dB to the signal-to-noise ratio. This derivation was based on the following assumptions. First, the quantization error is a stationary white noise process; second, the quantization error is uncorrelated with the input signal; and third, the distribution of quantization error is uniform over each quantization interval [17].

The first assumption is true when the input signal fluctuates in a complicated manner. The second assumption can be met if there are enough quantization levels available so that the step-size is properly small. The step-size can be determined empirically by applying various step sizes and examining the correlation between the quantization error and the input signal. The third assumption can be true if the range of the quantizer is set so as to match the peak-to-peak range of the signal, which is difficult to meet due to time varying characteristics of seismic signals. The time varying characteristics of seismic data are shown in Figure 9a and 9b indicating that the peak-to-peak range varies significantly from one time frame to another. Thus the number of quantization levels are not fully used in each frame, and the SNR in Equation (3.13) may not be achieved. Also, the uniform step-size quantization suffers from the dependence upon the signal variance. For example, the $(|X_{\max}|/\sigma_x)$ for the vibroseis data may vary from four to twelve, which can be seen in Figures 8a and 8b. This indicates a significant reduction in SNR. For this reason, quantization techniques which are less sensitive to the signal variance are investigated in the following.

Logarithmic Step-Size Quantization

The logarithmic quantizer, often called the instantaneous compressor/expander or simply compander, resolves the dependence of the SNR upon the signal variance by logarithmically spaced quantization levels. This can be alternatively achieved by quantizing the logarithm of the input rather than the input itself. Also, the companding may be used to improve the signal-to-noise ratio by producing effectively non-uniform quantization so that the largest quantization errors should be reserved for the least probable samples.

The signal-to-noise of the logarithmic quantizer is [6]

$$\text{SNR} = \frac{1}{\sigma_e^2} \quad (3.14)$$

This equation shows that the SNR depends upon only the step-size. Since the logarithm of very small numbers can be very large, this type of quantizer in general needs infinite number of quantization levels and therefore is impractical.

For this reason, Smith [18] has developed the alternative compression characteristics called μ -law. The μ -law is expressed as

$$\begin{aligned} Y(n) &= F[X(n)] \\ &= X_{\max} \frac{\log 1 + \mu \frac{|X_n|}{X_{\max}}}{\log(1 + \mu)} \text{sign}(X(n)) \end{aligned} \quad (3.15)$$

The parameter μ controls the degree of compression and may be chosen that large changes in the input produce relatively small changes in the output. When μ is zero, it corresponds to uniform step-size quantization. The μ -law step-size with μ -value of 500 and several μ -curves

with associated μ -values are illustrated in Figure 19. The signal-to-noise ratio of this quantizer for $\mu > 0$ is [18]

$$\begin{aligned} \text{SNR}(\text{dB}) = & 6b + 4.77 - 20 \log (\ln(1 + \mu)) \\ & - 10 \log \left(1 + \frac{\sigma_x^2}{\mu \sigma_x} + 2 \frac{\sigma_x^2}{\mu \sigma_x} \right). \end{aligned} \quad (3.16)$$

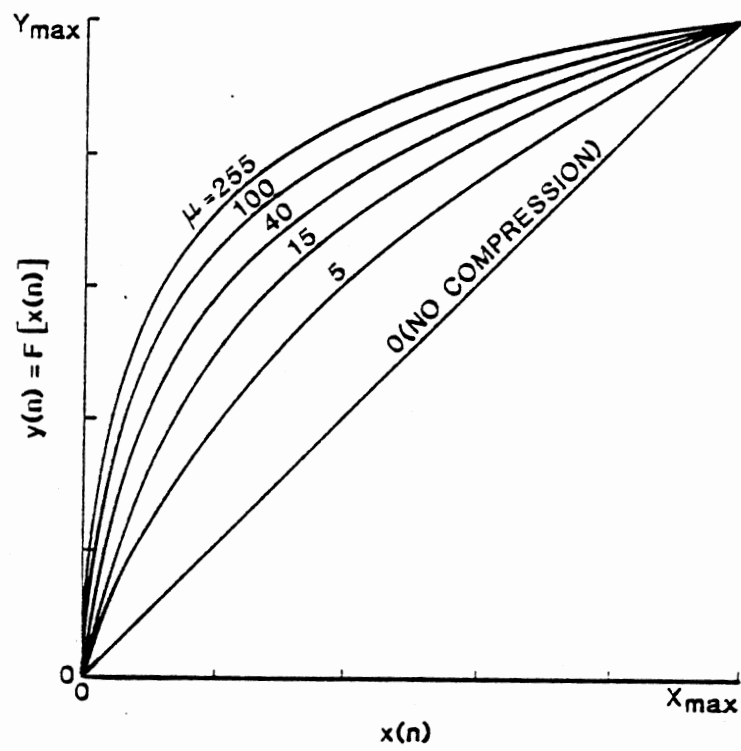
From this equation, it can be seen that the dependence of SNR upon the signal variance can be reduced by controlling the μ -value. This will be discussed in more detail later.

Both uniform and non-uniform step-size quantization has limitations when a signal has time-varying properties. For this reason, an adaptive step-size quantization technique is investigated in the next section.

Adaptive Step-Size Quantization

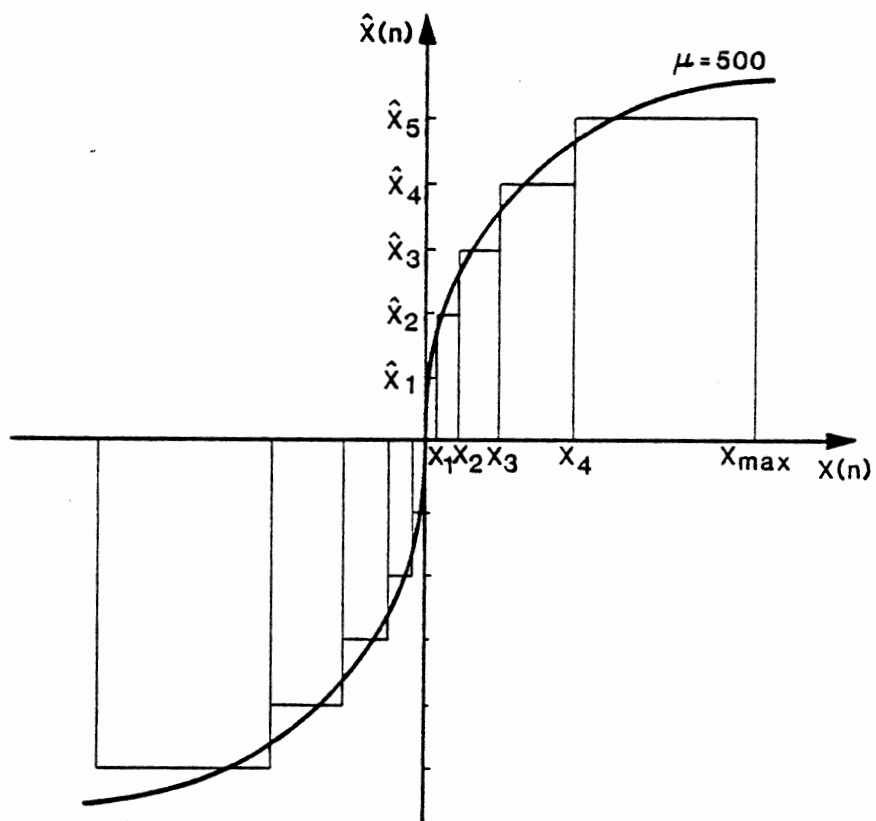
The basic idea of adaptive quantization is to let step-size vary so as to match the variance of the input signal. This implies that it is necessary to obtain an estimate of the time varying amplitude properties of the input signal. Sample-to-sample changes (or rapid changes within a few samples) and syllabic changes (or slowly varying) [6] need to be considered. For simplicity, the amplitude changes of sample-to-sample are used as the basis of the step-size adaptation for seismic data.

In general, there are two schemes in adaptive quantization, feed forward and feed backward quantization. When the step-size is adjusted according to the input itself, it is referred to as a feed-forward adaptive quantizer. When the step-size is adapted on the basis of the previous output of the quantizer, it is referred to as feed-backward quantizer. The feed-backward adaptation is based on the assumption that adjacent



(a) μ -curves (After Smith [18])

Figure 19. μ -curves and Distribution Levels



(b) Distribution Levels ($\mu = 500$)

Figure 19. (Continued)

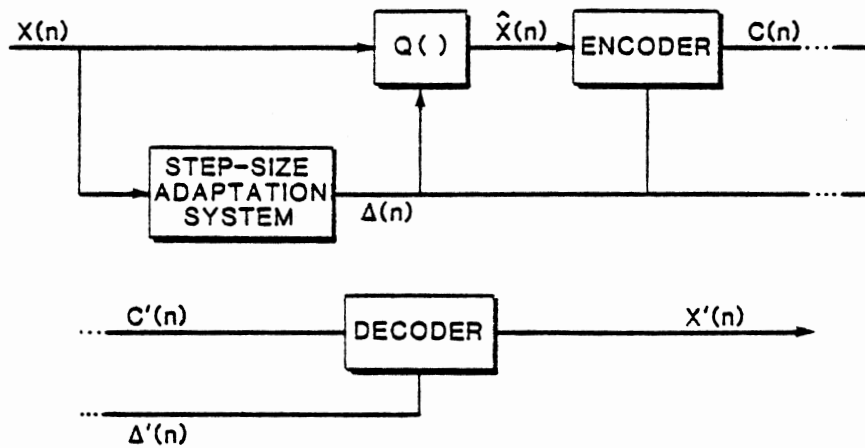
samples do not vary much, that is, sample-to-sample correlation is high. The feed forward and feed backward quantization schemes are shown in Figures 20a and 20b, respectively.

The feed-forward adaptation needs to transmit the step-size information for decoding, as the decoder cannot generate the step-size without current input data, while the feed-backward adaptation allows for the computation of step-size at the decoder in the absence of channel errors. This is a distinct advantage of feed-backward adaptation, especially when data compression is a critical issue. For this reason, feed-backward adaptation is used in this research and is discussed in detail in a later part of this section. It should be pointed out that the feed-backward adaptation has increased sensitivity to errors in the codewords, since such errors imply not only an error in the quantization level but also in the step-size [19].

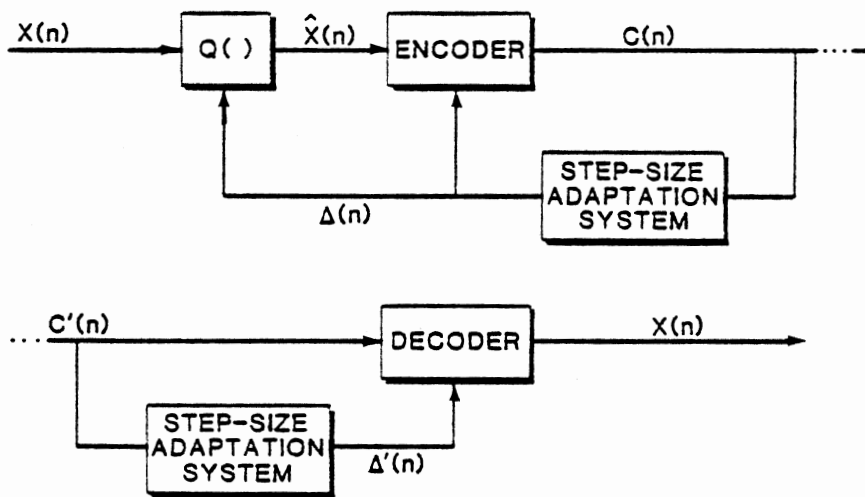
The step-size computation associated with the feed-backward quantizer is given by

$$\Delta(n) = M(|C(n-1)|) \cdot \Delta(n-1) \quad (3.17)$$

where $C(n-1)$ is the previous output code, $\Delta(n-1)$ is the previous step-size, and M is the multiplier array. It can be seen from Equation (3.17) that the current step-size is obtained by multiplying a selected multiplier and the previous step-size. The multiplier is selected from the multiplier array indexed by the absolute value of the previous output code. The multiplier is designed so that the entries in the first half of the array are less than one and the entries in the second half of the array are greater than one. Thus, the step-size will be reduced (increased) if the value of the previous code is less than (greater than) half of the code range. It



(a) FEED-FORWARD ADAPTATION



(b) FEED-BACKWARD ADAPTATION

Figure 20. Adaptive Quantization Procedures
(After Rabiner and Schafer [6])

is easy to see that the initial step-size is very critical as the succeeding step-sizes are proportional to the initial step-size. The optimum initial step-size is obtained by the following formula [20].

$$\Delta_{\text{opt}} = E[(X(n) - X(n-1))^2]^{\frac{1}{2}} \ln(2F_o) \quad (3.18)$$

$$F_o = F_s / 2 F_n$$

where F_s is the sampling frequency and F_n is the Nyquist sampling rate.

The SNR of this technique is derived in a similar manner as the uniform step-size quantization technique, as the step-size is uniform for each quantization instance. Basically, this technique pursues to meet the third assumption of the uniform quantization so that every bit is used efficiently.

Differential Input Quantization

In the last few sections, various quantization techniques have been examined with respect to their step-size decision process. It has been found that the signal variance influences SNR for uniform quantizers. Non-uniform quantizers were considered to reduce the dependency in signal variance to improve SNR. Another method of improving SNR is by coding the difference signal, which is the difference between the input and the predicted value. This method is called the differential input quantization technique.

When a predictor is designed based on a system model or based on mathematical derivation of an optimal filter, it is referred to as a predictive coding technique. It is differentiated from a simple differential input quantization technique, delta modulation, where the difference between adjacent samples is quantized and coded. Delta modulation is a

suboptimal technique; it has a significant advantage over a predictive coding technique in computational simplicity; it uses one delay instead of an optimal filter [21].

This section focuses on general considerations in measuring the SNR and the details of predictive coding techniques will be discussed in Chapter IV. The SNR of the differential input quantization is [22]

$$\text{SNR} = \frac{\sigma_x^2}{\sigma_e^2} = \frac{\sigma_x^2}{\sigma_d^2} \cdot \frac{\sigma_d^2}{\sigma_e^2} = G_p \cdot \text{SNR}_q \quad (3.19)$$

$$G_p = \frac{\sigma_x^2}{\sigma_d^2}, \quad \text{SNR}_q = \frac{\sigma_d^2}{\sigma_e^2}$$

where σ_d^2 is the variance of the difference signal. Equation (3.19) shows that signal-to-noise ratio consists of prediction gain and the signal-to-noise ratio due to quantization. Assuming the prediction gain is greater than one, it can be expected to improve the overall signal-to-noise ratio. In cases of predictive coding, higher order of filter contributes to better prediction. The prediction gain is dependent on the performance of the predictor and it can be maximized by minimizing the prediction error. For applying delta modulation, high correlation of adjacent samples is necessary in order to achieve a good prediction gain. When delta modulation employs more than two quantization levels, it is called a differential pulse code modulation (DPCM). When adaptation quantization is used, it is referred to as an adaptive differential pulse code modulation (ADPCM). Usually, nonuniform quantization techniques are not applied since the variance of difference signal is assumed to be small.

The performance of various digital coding techniques--ADPCM, DPCM,

APCM, and LPCM--are illustrated in Table II for 16 traces of vibroseis data. The ADPCM shows the best signal-to-noise ratio.

Optimum Encoding Techniques

In general, there are two types of encoding, source encoding and channel encoding [23]. The main subject of this section is source encoding, defined as the process of converting an information source signal into a binary sequence. An optimum encoder for N symbols corresponds to an average bit rate (\hat{H}_N), which approaches the source entropy (H) as N approaches infinity [23]. This can be expressed as

$$\hat{H}_N = 1/N \sum_{i=1}^q n_i p_i \rightarrow 1/N \sum_{i=1}^q p_i \log_2 (1/p_i), \quad (3.20)$$

$$H = \lim_{N \rightarrow \infty} \hat{H}_N \text{ (bits/symbol)}, \quad (3.21)$$

where q is the number of messages encoded into the sequence of N symbols, p_i is the probability of the i th message m_i , and n_i is the optimum code word length for the i th code word, c_i .

Any solution of Equation (3.20) is an optimum encoding technique. One example given by Shannon and Fano [14] is discussed below. This algorithm has the property of assigning short (long) code word lengths for high (low) probability messages. Specifically, if q messages $m_1, m_2, m_3, \dots, m_q$ are ordered in decreasing probability, the code length of the i th message is computed from

$$\log_2(1/p_i) < n_i < 1 + \log_2(1/p_i) \quad (3.22)$$

where $p_1 \geq p_2 \geq \dots \geq p_q$, and n_i is an integer.

After the optimum code word length is computed for the i th message,

TABLE II
SNR(dB) OF 6-BIT DIGITAL CODING METHODS
ON VIBROSEIS DATA

TEST	ADPCM	DPCM	APCM	LPCM
1	37.1	32.5	26.7	31.8
2	38.3	34.1	27.6	32.2
3	36.6	33.6	28.6	32.0
4	37.2	31.5	26.1	31.6
5	36.6	30.7	25.0	31.8
6	36.7	30.9	24.9	30.8
7	35.9	31.2	26.2	31.5
8	36.0	30.6	27.5	31.7
9	33.7	31.2	27.8	31.2
10	35.5	33.1	29.6	30.2
11	39.0	28.7	20.5	31.7
12	34.4	30.4	23.2	31.3
13	39.9	38.1	29.1	32.3
14	39.3	36.0	29.9	32.2
15	37.5	35.4	30.6	31.8
16	44.7	43.7	31.0	32.6
AVG	36.8	31.0	25.0	30.8

a unique code word c_i is obtained by truncating the binary expansion of the probability function F_i after a maximum of n_i bits. Let

$$F_i = \sum_{k_i=1}^{i-1} p_k$$

then

$$c_i = (F_i)_{\text{binary } n_i \text{ bits.}}$$

Another example of an optimum encoding technique is the minimum redundancy coding method developed by Huffman [24]. The length of code words is inversely related to the probability of messages as in Shannon's algorithm.

These algorithms require code word tables, and therefore may also require large amounts of memory storage. Second, the message probability computation may not be feasible for some cases. However, these techniques still provide the ideal encoding, and thus can be used as a measure of the performance of other encoding methods. For example, the efficiency of an encoding technique (e) can be obtained by the ratio between the average bit ratio of the given encoding technique, H_N^I , and \hat{H}_N . That is,

$$e = \hat{H}_N / H_N^I. \quad (3.23)$$

Summary

This chapter has evaluated general data compression techniques with respect to their utilities for seismic data compression. The results showed that entropy reducing transforms and interpolation techniques are not adequate for seismic data compression due to the reversibility considerations and the low sampling rate. The prediction techniques which allow

for transmitting the prediction error signal instead of the original input signal were considered. They were found to be impractical for large size seismic data since computational complexities involved in obtaining a predictor increase as the square of the number of data. These techniques were suggested for impulsive seismic data for a restricted region. The restricted region is obtained by a data slicing technique where all impulsive waveforms are separated from random waveforms as will be discussed later in the next chapter.

Alternative approaches for seismic data compression were discussed with orthogonal transforms and digital coding methods. For both approaches various techniques were examined and their performances on seismic data compression were evaluated.

The performances including compression ratio and signal-to-noise ratio of the WHT, the DFT, and the DCT were illustrated in Table I. Also, the performances of various digital coding techniques--DPCM, ADPCM, LPCM, and APCM--were shown in Table II.

From these evaluations, it was concluded that a technique which combines an orthogonal transform and a digital coding method may be a proper approach for vibroseis data compression. In particular, the discrete cosine transform was shown to be the best choice for the orthogonal transformation with respect to computational complexity, compression ratio, and signal-to-noise ratio.

In the next chapter, data compression techniques for vibroseis data and impulsive seismic data will be investigated based on these observations of orthogonal transforms, digital coding methods, and prediction techniques.

CHAPTER IV

SEISMIC DATA COMPRESSION

Introduction

In the previous chapter, prediction, orthogonal transforms, and digital coding methods were suggested for seismic data compression. This chapter discusses combinations of these techniques for vibroseis data and impulsive seismic data.

As pointed out earlier, vibroseis data differ from impulsive seismic data in their statistical characteristics. For this reason, the data compression technique is developed separately for each type of seismic data. For vibroseis data compression, a "hybrid technique," which combines an orthogonal transform and a digital coding method is considered. For impulsive seismic data compression, the trace is divided into two sections, an impulsive section and a nonimpulsive section. For the first section, a "predictive coding" technique is investigated; and for the second section, the hybrid technique is considered.

The hybrid technique for seismic data compression can be evaluated in the following manner. First, the signal-to-noise ratio and the compression ratio of the hybrid technique are derived. These derivations are discussed with two aspects, orthogonal transforms and digital coding methods. Second, the selection method for the significant transformed vector coefficients is examined in relation to a threshold value decision scheme. For the retained coefficients, the choice of a digital coding

method is discussed based on empirical results of various coding methods. Third, the μ -law quantization is considered as a digital coding method for the retained coefficients with respect to the relation between the μ -value and the compression results. Finally, the implementation of the hybrid technique for a given set of vibroseismic data is illustrated, and the compression results are evaluated.

The predictive coding method for the first segment of the impulsive seismic trace can be analyzed as follows. First, the data segmentation consideration is studied with respect to the waveform characteristics of each segment. Also, the signal-to-noise ratio and the compression ratio of the trace are discussed in relation to the SNR and CPR of each segment. Second, the "linear prediction" algorithm is studied with respect to optimum predictor parameters based on the autocorrelation method and the Lattice method. Finally, the implementation of the predictive coding technique for a given set of impulsive seismic data is illustrated, and the compression results are evaluated.

Vibroseis Data Compression

The signal-to-noise ratio of the hybrid technique can be derived in terms of SNR for the orthogonal transform and the digital coding method. Explicitly,

$$\text{SNR} = \frac{\sigma_x^2}{\sigma_e^2} = \frac{\sigma_x^2}{\sigma_y^2} \frac{\sigma_y^2}{\sigma_e^2} \quad (4.1)$$

where σ_x^2 is the variance of the original signal, σ_y^2 is the variance of the orthogonally transformed vector coefficients, and σ_e^2 is the variance of noise. Using Parseval's theorem, it follows that $\sigma_x^2 = \sigma_y^2$ [25].

However, in the hybrid technique the insignificant coefficients are suppressed, and the error due to this suppression needs to be included in Equation (4.1). For simplicity, this error is included with the digital coding error, and expressed in the approximate form.

$$\begin{aligned}
 \text{SNR}^V &\approx \frac{\sigma_x^2}{\sigma_e^2} = \frac{\sigma_y^2}{\sigma_{y'}^2 + \sigma_q^2} \\
 &\approx \frac{1}{\frac{\sigma_{y'}^2}{\sigma_y^2} + \frac{\sigma_q^2}{\sigma_y^2}} \\
 &\approx \frac{\text{SNR}_1^V \cdot \text{SNR}_2^V}{\text{SNR}_1^V + \text{SNR}_2^V} \tag{4.2}
 \end{aligned}$$

where $\sigma_{y'}^2$ is the variance of the suppressed coefficients, and σ_q^2 is the variance of the quantization noise. The term SNR_1^V refers to the signal-to-noise ratio obtained from the orthogonal transform technique, and the terms SNR_2^V refers to the signal-to-noise ratio obtained from a digital coding method.

In order to maintain reversibility, most of the energy of the transformed vector should be preserved, and thus SNR_1^V should be far greater than SNR_2^V . From this relation, the SNR^V can be simply expressed as

$$\text{SNR}^V \approx \frac{\text{SNR}_1^V \cdot \text{SNR}_2^V}{\text{SNR}_1^V + \text{SNR}_2^V} \approx \text{SNR}_2^V \tag{4.3}$$

where it is assumed that

$$\text{SNR}_1^V \gg \text{SNR}_2^V .$$

The compression ratio can be expressed in general as

$$\text{CPR} = \frac{\text{Bit rate of the original signal}}{\text{Bit rate of the compressed signal}} \quad (4.4)$$

Using Equation (4.4) the compression ratio of the hybrid technique is given by

$$\text{CPR}^V = \frac{N \cdot b_o}{M \cdot b_1 + N} \quad (4.5)$$

where N is the number of data points in the original signal, b_o is the number of bits used per sample, M is the number of retained coefficients, and b_1 is the number of bits used per coefficient. The last term in the denominator of CPR, N , corresponds to the number of bits necessary for the bookkeeping array. The bookkeeping array preserves the information for selection and suppression of the coefficients.

In general, $M \cdot b_1 \gg N$, and Equation (4.5) can be approximated by

$$\text{CPR}^V \approx \frac{N}{M} \cdot \frac{b_o}{b_1} = \text{CPR}_1^V \cdot \text{CPR}_2^V \quad (4.6)$$

where CPR_1^V is the compression achieved by the orthogonal transform and CPR_2^V is the compression achieved by a digital coding method.

In the next section, CPR_1^V and SNR_1^V are examined in relation to the threshold value for the insignificant coefficient decision; also, the characteristics of the retained coefficients are discussed with respect to the digital coding techniques.

Selection of the Transformed Vector Coefficients

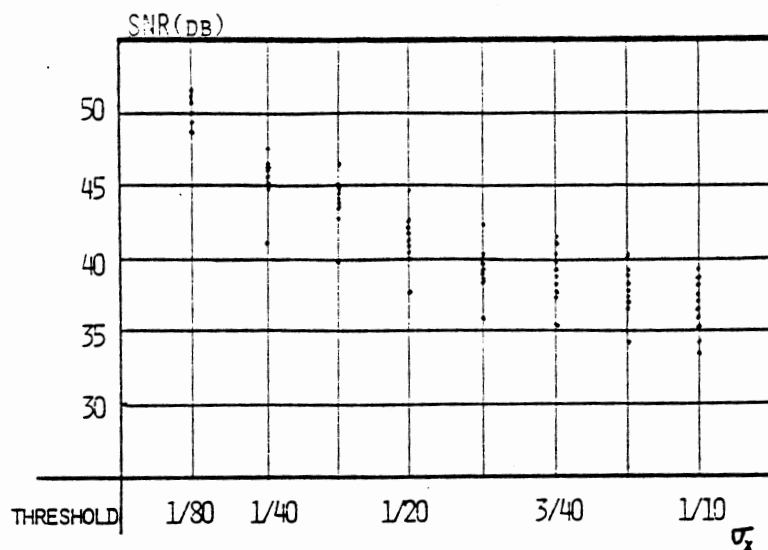
In Chapter II, it was shown that the discrete cosine transform (DCT) is the appropriate transform for the hybrid technique. For the selection of the significant coefficients, a threshold value computed from either the absolute maximum or the standard deviation of the coefficients needs

to be determined. From an empirical observation (see Figure 21), it was determined that the threshold value from the standard deviation gives more uniform results for compression and signal-to-noise ratio. Figure 21a and 21b show the distribution characteristics of the signal-to-noise ratio and the compression ratio as functions of the threshold value for 16 traces of a given set of vibroseis data. From Figure 21, it can be seen that the threshold value varies from 1/80 to 1.10 of the standard deviation, and varies from 1/800 to 1/100 of the absolute peak value. Also, the compression ratio is more sensitive to the threshold value than the signal-to-noise ratio.

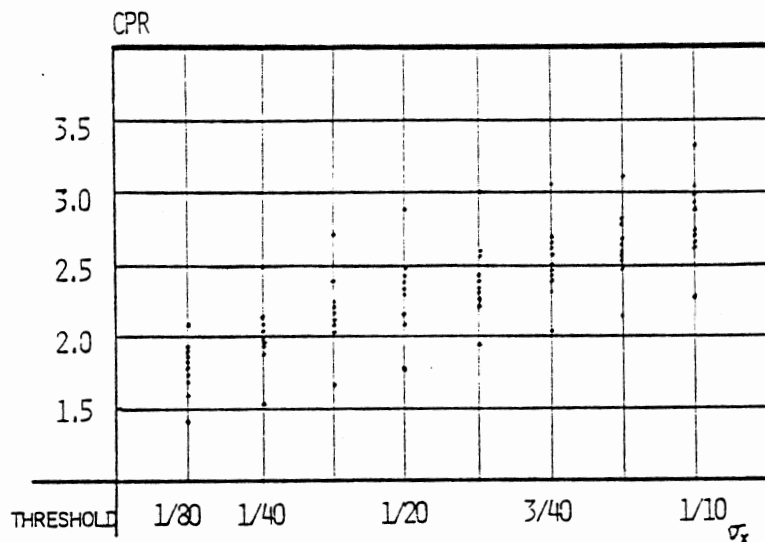
For selection of the significant coefficients, $(1/20) \sigma_x$ is used as the threshold value for the insignificant coefficients. With this threshold, the CPR ranges from 2.0 to 2.5, and the SNR is approximately 42 dB. These values indicate that by discarding $1/10^4$ of the total energy, at least two to one compression can be achieved from the DCT method. These can be denoted as $\text{SNR}_1^V = 10^4$ and $\text{CPR}_1^V > 2$.

The selection method has been applied to various lengths of partial vibroseis trace and their compression ratios are illustrated in Figure 22. It can be observed that whole trace shows the best compression result and less than 512 points of partial vibroseis trace shows almost no compression.

In order to obtain SNR_2^V and CPR_2^V , a proper digital coding method needs to be chosen to represent the retained coefficients. The compression results of various digital coding techniques, ADPCM, LPCM, APCM, and μ -law, are compared in Table III. This table shows that the μ -law quantization gives a better compression result than others. This result is due

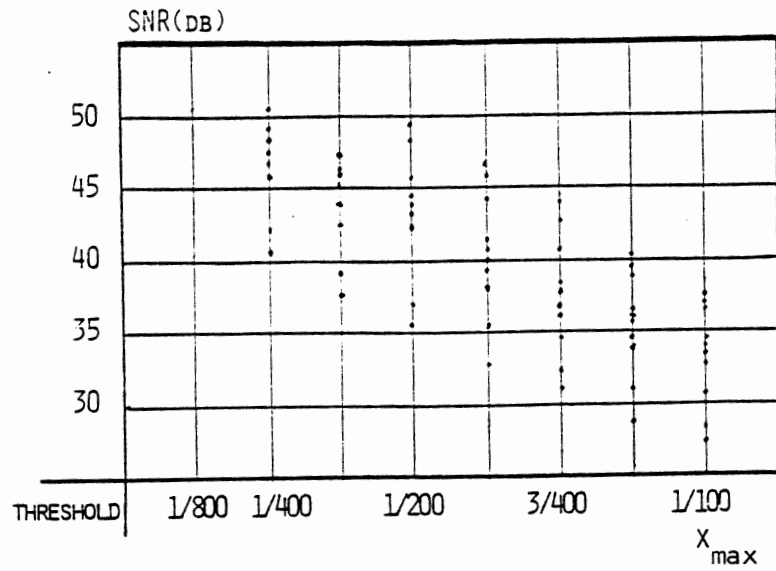


(a) SNR Distribution of Suppressed DCT Coefficients

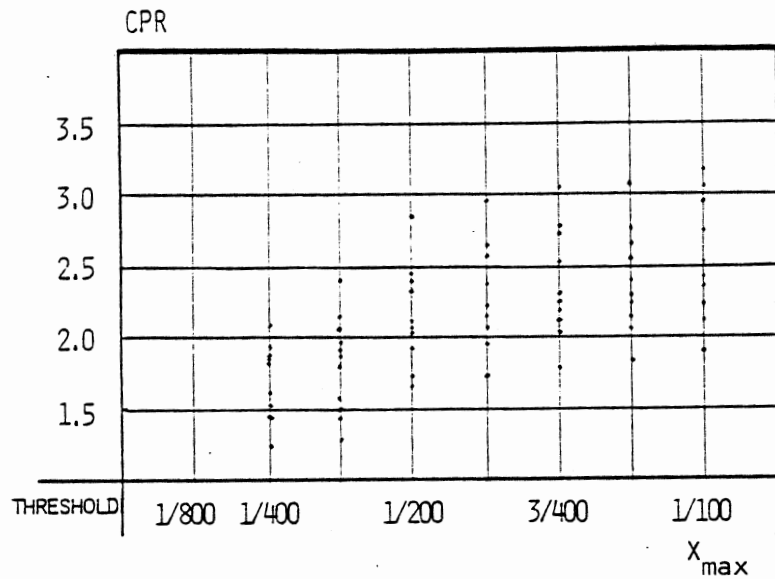


(b) CPR Distribution of Selected DCT Coefficients

Figure 21. Distributions of SNR and CPR as a Function of Threshold



(c) SNR Distribution of Suppressed DCT Coefficients



(d) CPR Distribution of Selected DCT Coefficients

Figure 21. (Continued)

TABLE III
 SNR(dB) OF 6-BIT DIGITAL CODING METHODS
 ON SELECTED DCT COEFFICIENTS

TEST	APCM	ADPCM	LPCM	μ -LAW
1	25.92	18.82	24.26	27.38
2	23.96	13.36	23.96	26.30
3	14.69	13.08	24.24	26.53
4	16.27	18.94	24.83	30.46
5	19.50	12.56	25.65	30.78
6	13.68	11.20	24.07	26.56
7	18.20	22.06	25.66	28.34
8	23.84	19.80	24.92	26.87
9	21.43	20.35	26.11	27.86
10	17.27	14.60	26.54	27.72
11	18.94	16.83	26.93	27.88
12	19.51	19.09	26.38	26.67
13	26.38	28.19	28.80	29.41
14	20.12	23.62	28.27	30.11
15	19.70	18.71	27.97	30.22
16	14.86	10.45	23.86	30.43
AVG	19.64	17.60	25.84	28.34

to the distributional characteristics of the retained coefficients and low correlation between adjacent coefficients.

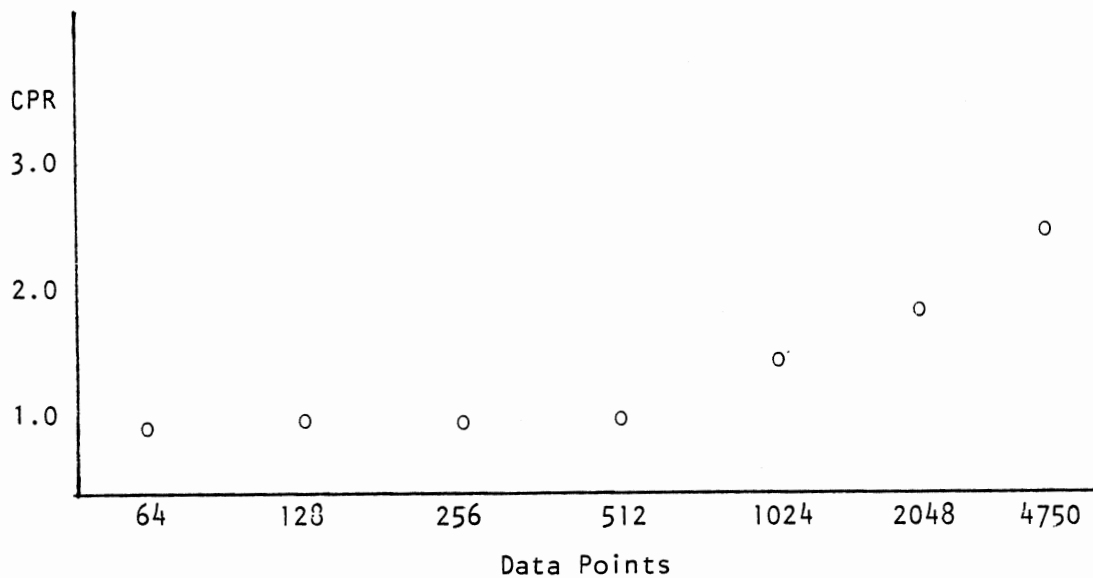


Figure 22. Compression Ratio of Partial Vibroseis Trace

μ -Law Quantization

The μ -law quantization method is applied for quantizing the retained discrete cosine transform coefficients. In Chapter III, it was pointed out that the μ -value controls the dependence of the signal-to-noise ratio on the signal variance. Figure 23 illustrates the signal-to-noise ratio on the signal variance. Figure 23 illustrates the signal-to-noise ratio of μ -law quantization as a function of the signal variance and μ -value [18]. It can be seen that the dependence of the signal-to-noise ratio decreases as the μ -value increases. For example, with $\mu = 500$, the signal-

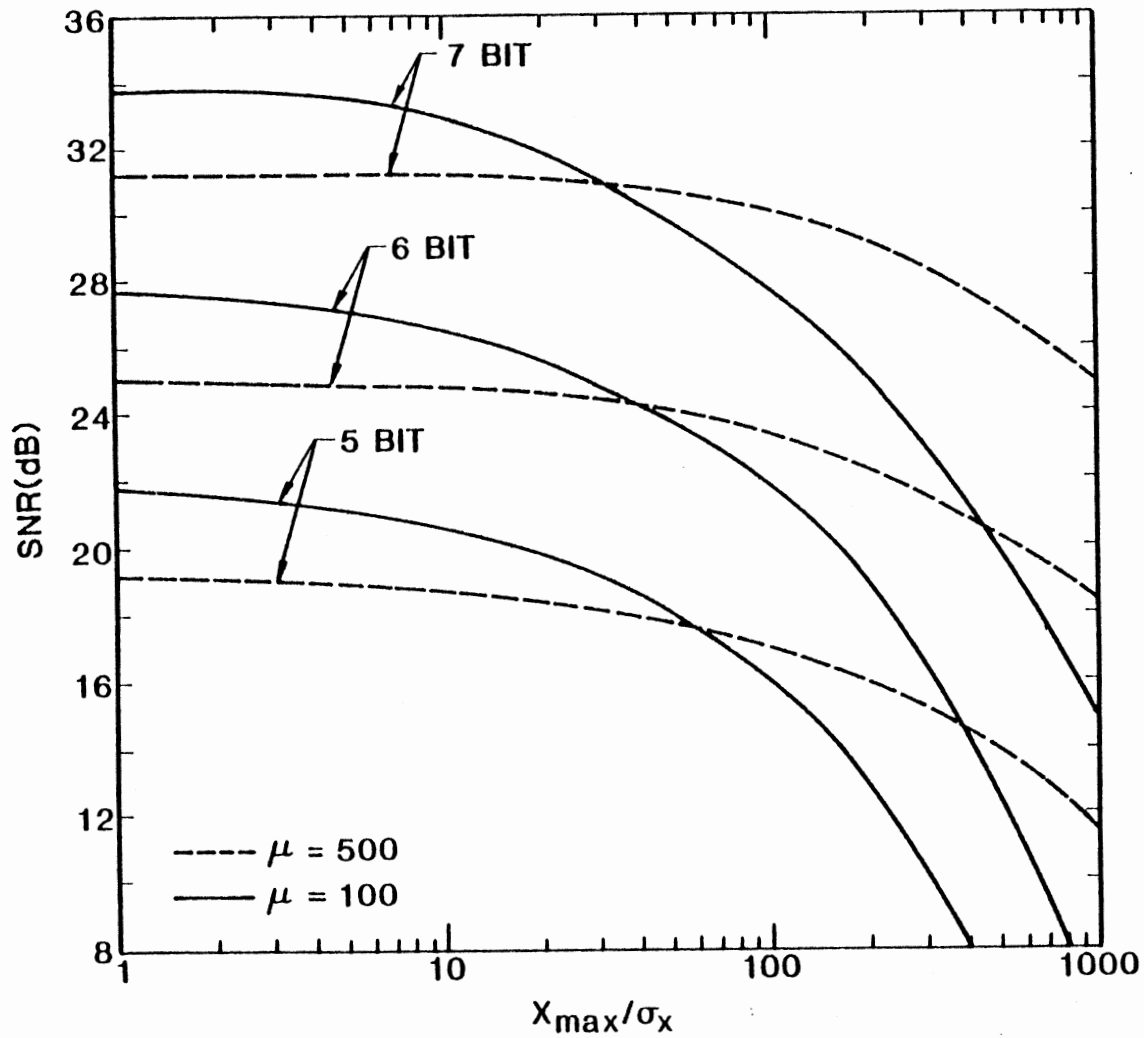


Figure 23. SNR(dB) of μ -law Quantizer as a Function of $|X_{\max}|/\sigma_x$
 (After Rabiner and Schafer [6])

TABLE IV
 μ -VALUES AND $|X_{\max}|/\sigma_x$

$ X_{\max} /\sigma_x$	μ -value
0 - 4	0.0
4 - 15	10.0
15 - 20	30.0
20 - 25	50.0
25 - 35	100.0
35 -	

to-noise ratio remains stable over $\sigma_x < X_{\max} < 100 \sigma_x$. For higher values of μ , the expected signal-to-noise ratio is lower. For this reason, a value for μ is adapted for each trace of the seismic data based on the value $(|X_{\max}|/\sigma_x)$. It is shown in Table IV and the table is obtained empirically.

For encoding the quantized levels, PCM is used. The number of bits per code word is determined based on the required signal-to-noise ratio referred to as SNR_2^V in the previous section. The results corresponding to 6- and 7-bit encoding with PCM are illustrated in Table V. It was pointed out earlier that SNR^V approaches SNR_2^V when SNR_1^V is much higher than SNR_2^V (see Equation (4.3)). Noting that SNR^V is required to be greater than 30 dB in order to maintain reversibility, SNR_2^V is desired to be at least 30 dB. Table V indicates that a 7-bit PCM is required for encoding the μ -law quantized levels when the μ -value is computed to be greater 10. It can also be seen that a 6-bit PCM can be applied for the traces whose μ -value is ≤ 10 .

From Figure 23 and Table V, it can be observed that adaptation of the μ -value prior to applying the μ -law quantization influences the compression significantly. The adaptation of the μ -value requires overhead computation of the value $(|X_{\max}|/\sigma_x)$.

Implementation of the "Hybrid Technique"

In the last two sections, the hybrid technique was investigated using the discrete cosine transform followed by a μ -law quantizer with a 7-bit PCM encoder. The encoding and decoding procedures of the hybrid technique are illustrated in Figure 24a, b, c, and d.

In the transmitter, the seismic data are transformed via the DCT, and

TABLE V
RESULTS OF μ -LAW QUANTIZATION

TEST	μ -VALUE	6-BIT		7-BIT	
		SNR(dB)	CPR	SNR(dB)	CPR
1	30.0	28.02	8.50	33.68	6.38
2	100.0	25.85	7.62	31.63	5.71
3	100.0	26.43	8.80	31.55	6.60
4	10.0	30.45	9.13	36.48	6.84
5	10.0	30.78	10.55	36.70	7.91
6	10.0	25.69	8.09	31.60	6.07
7	30.0	27.93	10.02	33.78	7.51
8	30.0	27.61	11.71	34.32	8.78
9	50.0	27.60	9.03	34.14	6.78
10	30.0	27.16	9.73	33.59	7.30
11	100.0	27.47	9.48	33.41	7.11
12	30.0	28.42	8.87	34.64	6.65
13	50.0	27.81	7.96	34.00	5.97
14	30.0	26.69	8.69	33.17	6.52
15	30.0	28.37	8.96	34.52	6.72
16	50.0	27.76	9.18	33.74	6.89
AVG		27.74	9.15	33.81	6.86

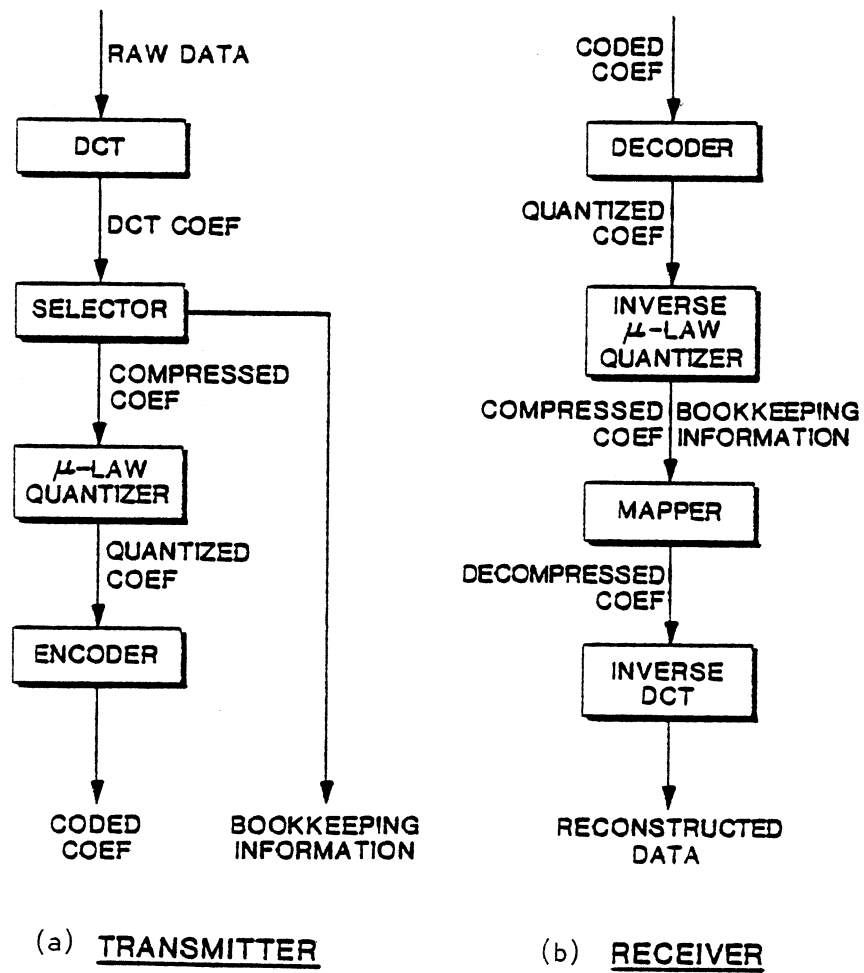
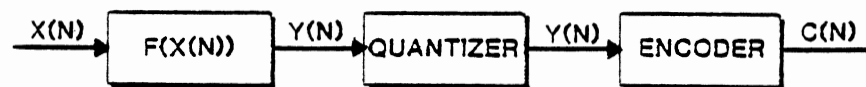


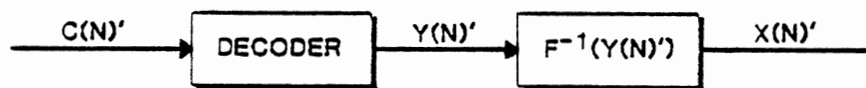
Figure 24. "Hybrid Technique" Procedures



$$Y(N) = F(X(N))$$

$$= \frac{X_{\text{MAX}} \cdot \text{LOG} \left(1 + \mu \frac{X(N)}{X_{\text{MAX}}} \right) \cdot \text{SIGN}(X(N))}{\text{LOG}(1 + \mu)}$$

(c) μ -LAW QUANTIZER



$$X(N)' = F^{-1}(Y(N)')$$

$$= \frac{Y'_{\text{MAX}}}{\mu} \left\{ \text{LOG}^{-1} \left[\frac{Y(N)' \cdot \text{LOG}(1 + \mu)}{Y'_{\text{MAX}}} \right] - 1 \right\} \cdot \text{SIGN}(Y(N)')$$

(d) INVERSE μ -LAW QUANTIZER

Figure 24. (Continued)

the transformed vector coefficients above the threshold value are selected by the selector shown in Figure 24a. The information of selection is transmitted to the receiver through the bookkeeping array. The retained coefficients are quantized with a μ -law quantizer and encoded for transmission. Details are shown in Figure 24c and d.

In the receiver, the coefficients are decoded and then dequantized with an inverse μ -law quantizer. Then, using the bookkeeping information, the mapper reconstructs the transformed coefficients by inserting zeros corresponding to the insignificant coefficients. These coefficients are inversely transformed via the IDCT. The output is the reconstructed seismic data; though it includes a slight distortion, it maintains the reversibility property.

Compression results are evaluated with respect to the signal-to-noise ratio and the compression ratio. The signal-to-noise ratio is computed from

$$\text{SNR}^V (\text{dB}) = 10 \log \frac{E[(x - x')^2]}{E[x^2]} \quad (4.7)$$

where x is the raw data and x' is the reconstructed data. The compression ratio is computed from Equation (4.5), simply replacing N with 4750,

$$\text{CPR}^V = \frac{4750 \times 20}{7M + 4750} \quad (4.8)$$

where M is the average number of retained coefficients and 4750 in the denominator indicates the size of the bookkeeping array.

Table VI presents the results of 16 traces of the vibroseis data. It should be pointed out that the SNR is approximately 2 dB lower than SNR_2^V due to the noise introduced by suppressing the insignificant DCT

TABLE VI
RESULTS OF VIBROSEIS DATA COMPRESSION

TEST	SELECT (DCT)	μ -LAW (7-BIT)	OVER-ALL SNR	OVER-ALL CPR
1	43.8	33.7	32.6	4.91
2	42.9	31.6	30.9	4.53
3	43.7	31.7	29.3	5.03
4	43.6	36.5	34.3	5.15
5	44.2	36.7	34.4	5.72
6	43.9	31.6	30.7	4.73
7	44.5	33.8	32.3	5.52
8	43.5	34.3	31.4	6.15
9	45.2	34.1	31.7	5.03
10	43.4	33.6	31.7	5.31
11	42.0	33.4	31.5	5.27
12	44.1	34.6	33.1	5.04
13	43.2	34.0	32.9	4.68
14	42.7	33.2	32.3	4.99
15	43.2	34.5	33.4	5.00
16	41.5	33.7	32.8	5.19
AVG	43.5	33.8	32.2	5.16

coefficients. The bookkeeping array may not be necessary since the transformed vector of each trace shows that most energy is concentrated in two major lobes, and the location of the lobes is common for all traces. These lobes are located at the beginning and ending of the trace. Also, for some traces, 6-bit μ -law quantization can provide over 30 dB of SNR corresponding to a small μ -value. For these cases, i.e., 6 b/sample and no bookkeeping array, the CPR^V can be computed as

$$CPR^V = \frac{4750 \times 20}{6M} \quad (4.9)$$

which gives a significant enhancement in compression ratio. From the statistical observation of the CPR, it can be concluded that compression ratios can be achieved in the range of five-to-one to eight-to-one via the hybrid technique.

Impulsive Seismic Data Compression

The impulsive seismic data have a good deal of similarity from trace to trace, which is different when compared to vibroseis data. It was pointed out in Chapter II that an impulsive trace has two distinct parts. Most of the energy is contained in the beginning of a trace and the rest of the trace is relatively insignificant. Figure 25 illustrates this aspect of an impulsive trace.

In order to examine the energy distribution characteristics, 25 traces of 1000 data points each are divided into five parts with equal duration. These 25 corresponding parts are concatenated to form a segment. For example, the first segment contains 25 first parts from the respective traces. Figure 26a through e illustrates the five segments taken from a 25-trace impulsive data. Considering the similarities, the five

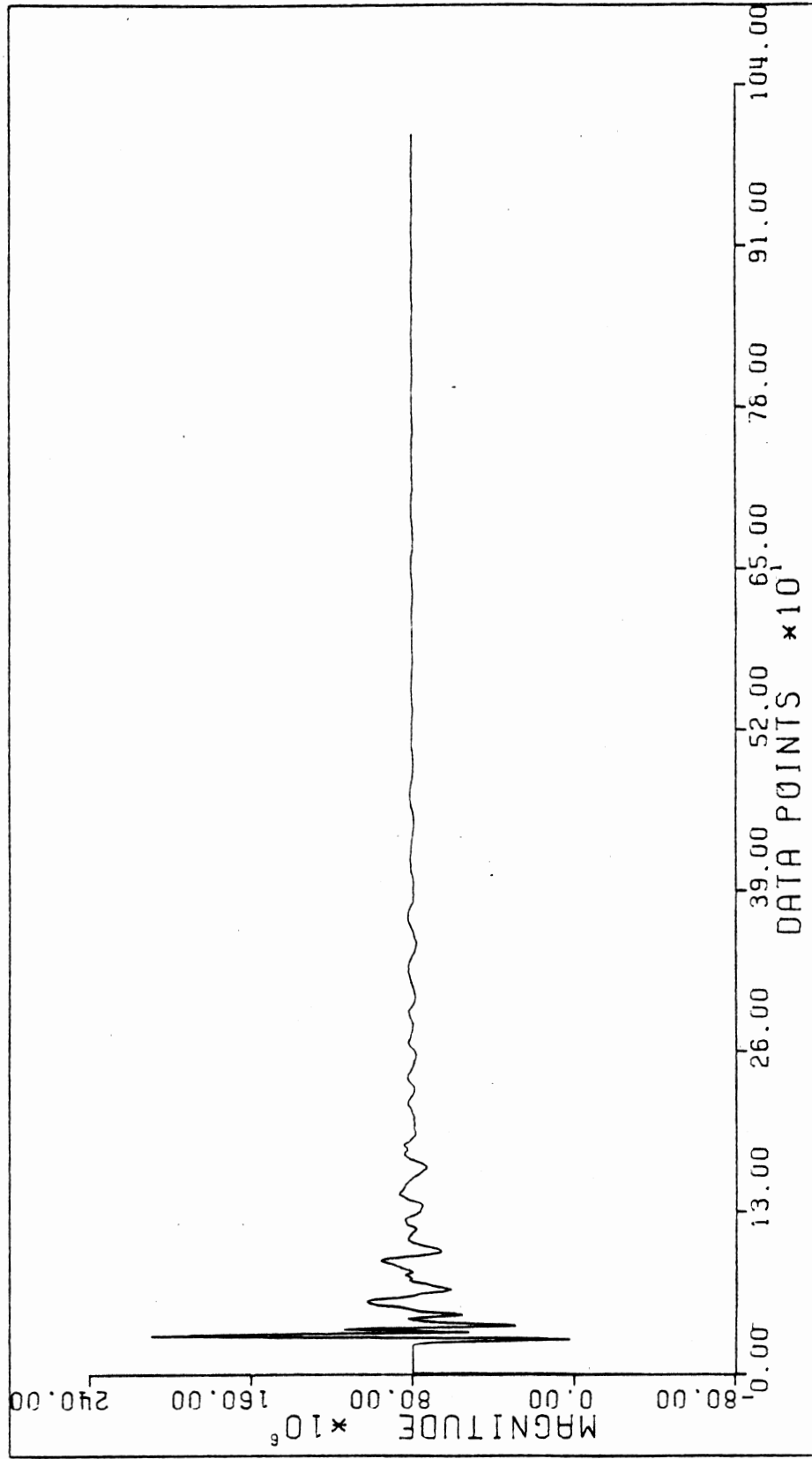
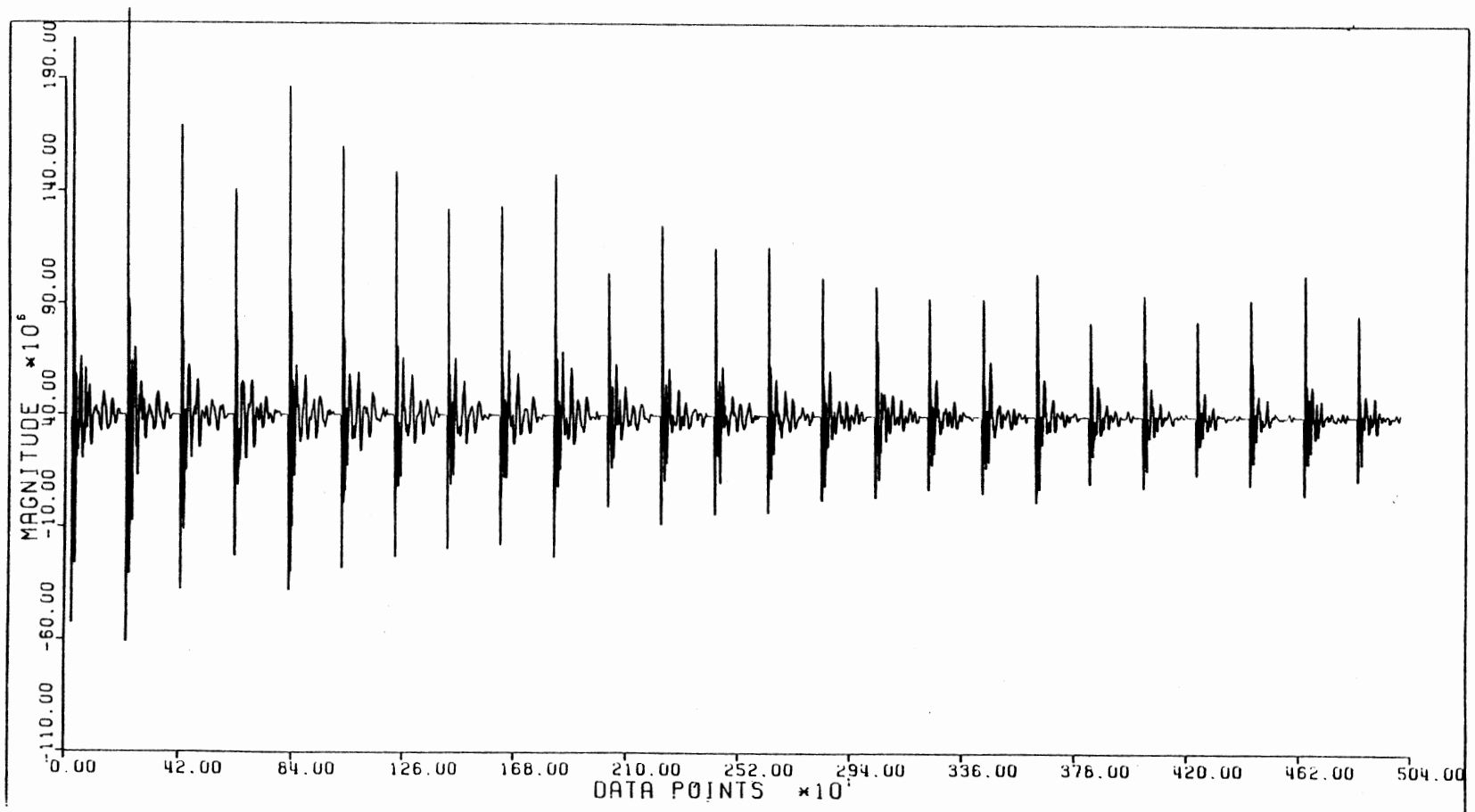
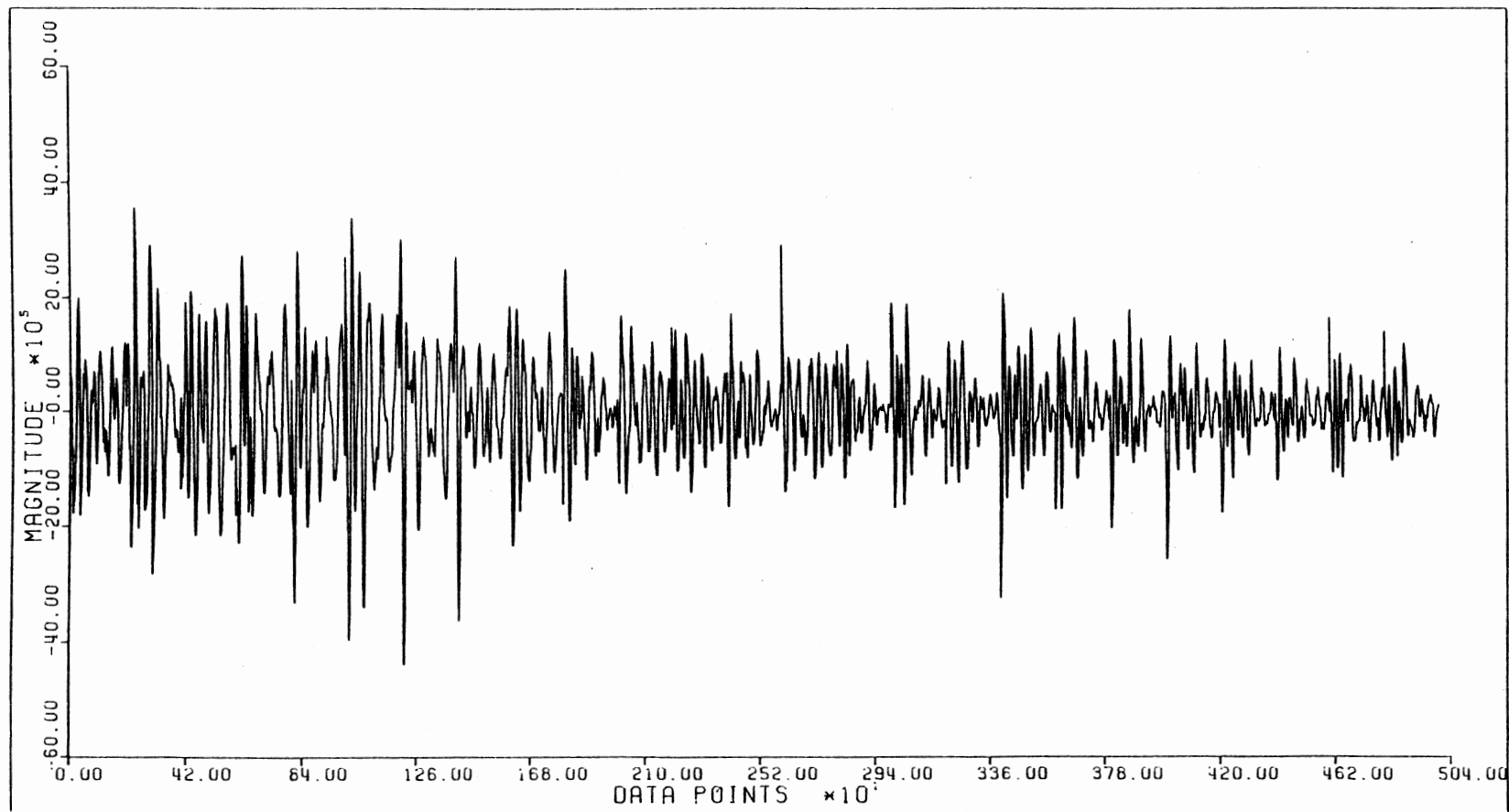


Figure 25. One Trace of Impulsive Seismic Data



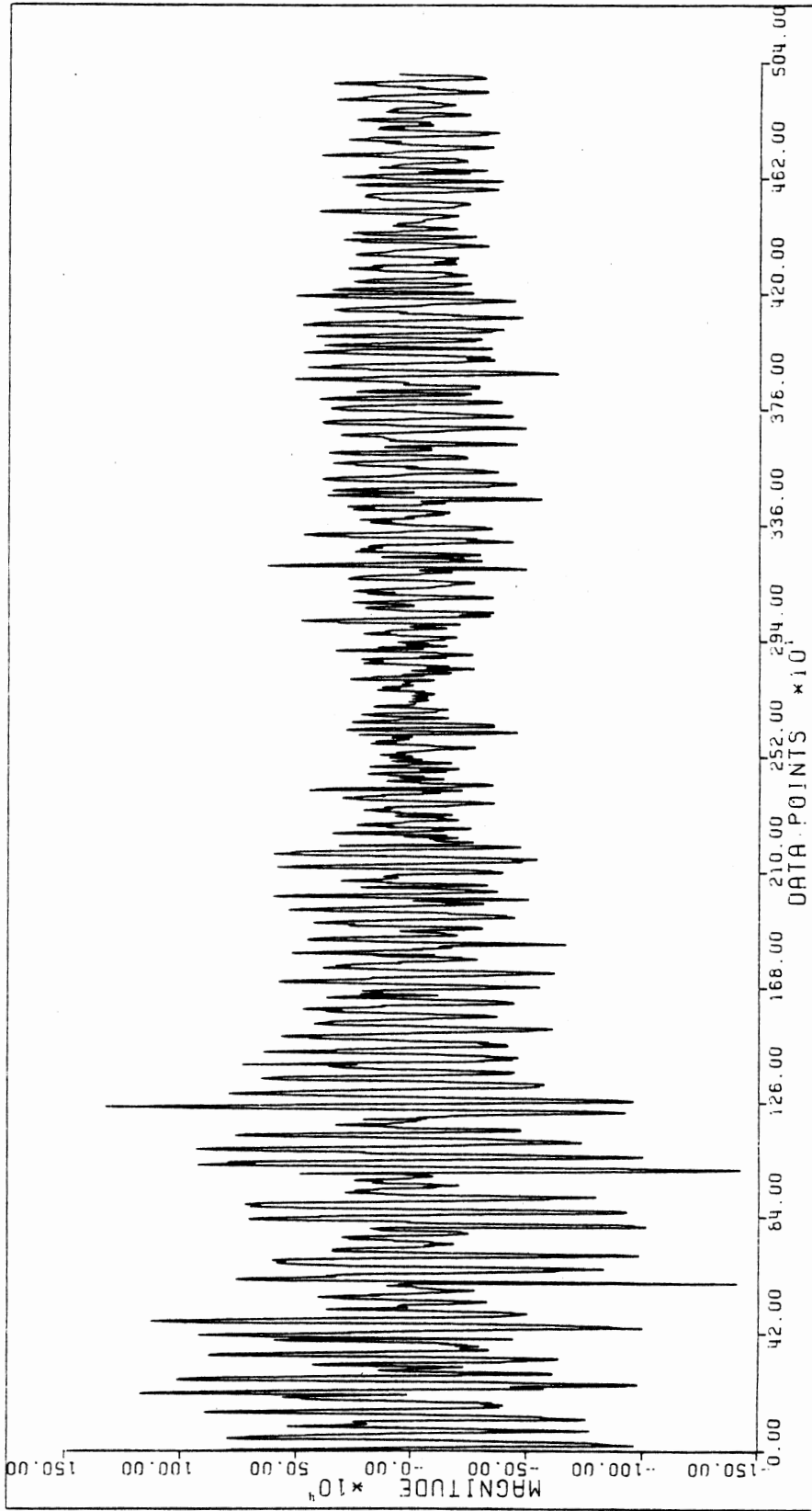
(a) First 200 Points of 25 Impulsive Seismic Data Traces

Figure 26. Illustrations of Five Segments of Impulsive Seismic Data



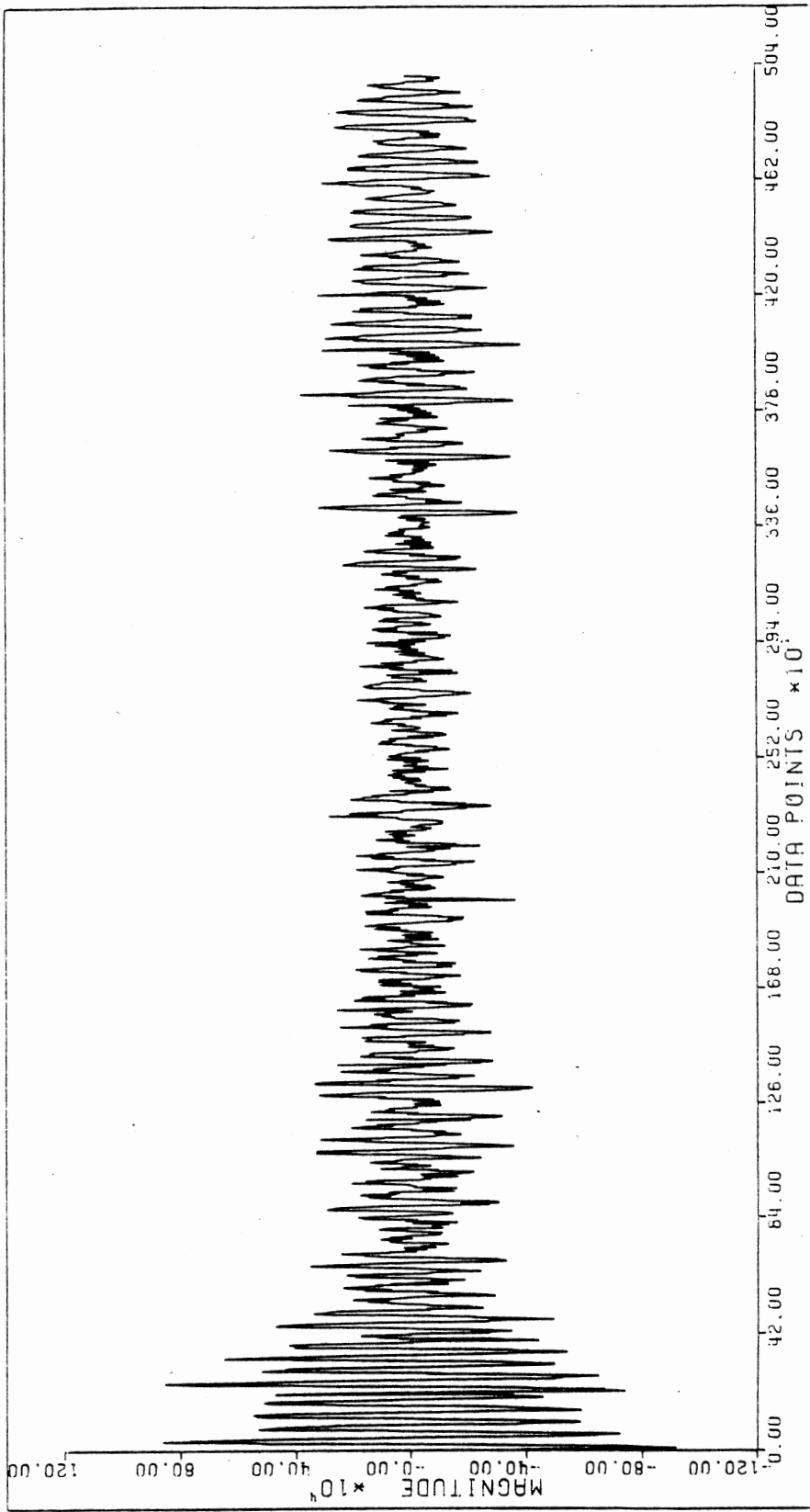
(b) Second 200 points of 25 Impulsive Seismic Data Traces

Figure 26. (Continued)



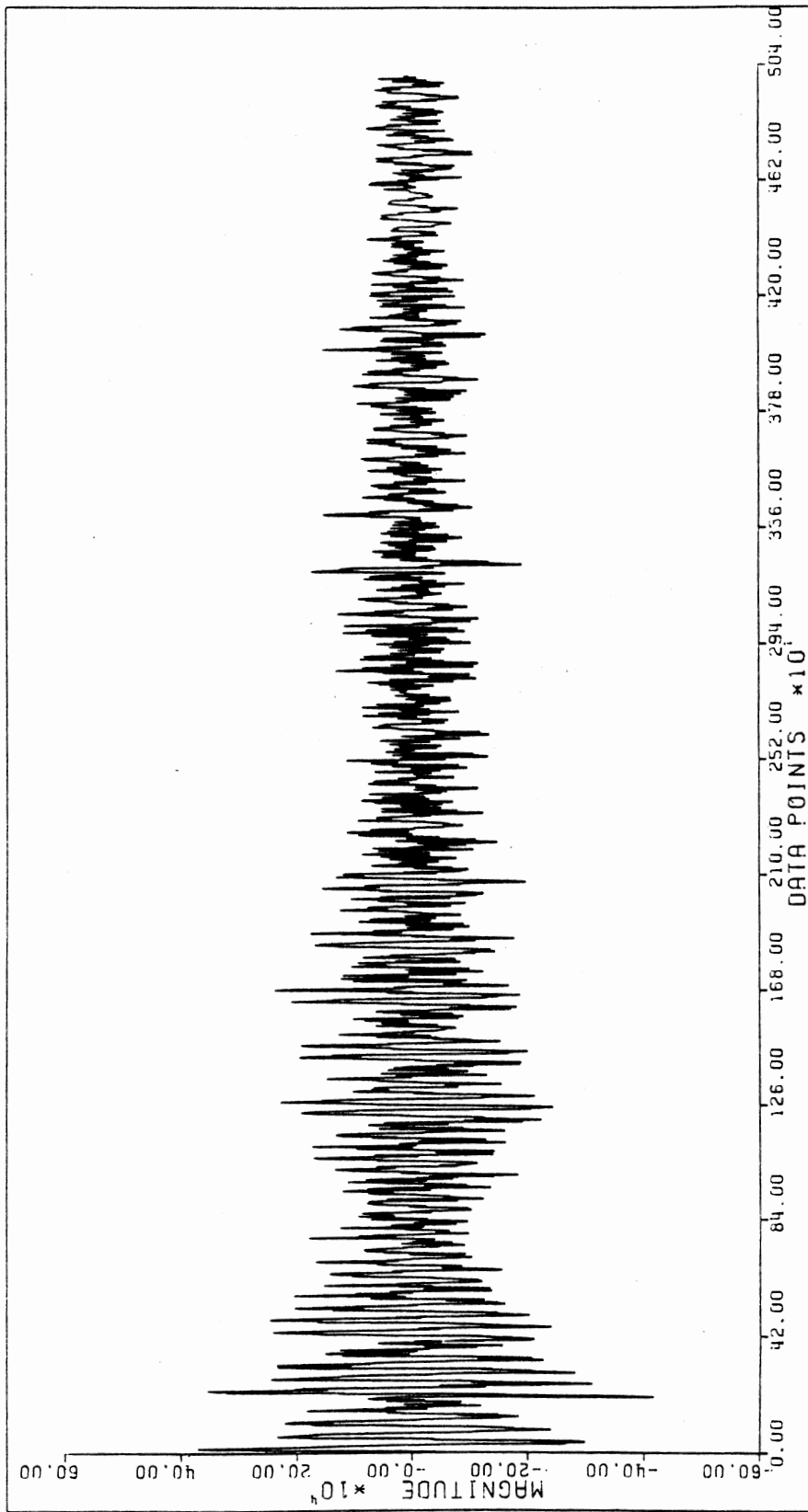
(c) Third 200 points of 25 Impulsive Seismic Data Traces

Figure 26. (Continued)



(d) Fourth 200 Points of 25 Impulsive Seismic Data Traces

Figure 26. (Continued)



(e) Fifth 200 Points of 25 Impulsive Seismic Data Traces

Figure 26. (Continued)

segments are grouped into two sections; since the first part in each trace contains most of the energy, the first segment is identified by impulsive section. The remaining four segments contain insignificant amounts of energy and are referred by nonimpulsive section.

It was found that the impulsive section contains more than 99 percent of the total energy. However, the amount of energy does not necessarily correspond to the amount of information in seismic data analysis. For this reason, signal-to-noise ratio of the both sections, SNR_1^I and SNR_2^I , need to be maintained uniformly for preserving information contained in each section. That is, $SNR_1^I \cong SNR_2^I$.

For simplicity, it will be assumed that the two sections are independent of each other. That is, it will be assumed that $\sigma_x^2 = \sigma_{x_1}^2 + \sigma_{x_2}^2$, where x_1 and x_2 correspond to the two sections. Noting $\sigma_{x_1}^2 \gg \sigma_{x_2}^2$, $\sigma_{x_1}^2 \gg \sigma_{e_2}^2$, it follows that overall SNR^I can be expressed as

$$SNR^I = \frac{\sigma_x^2}{\sigma_e^2} = \frac{\sigma_{x_1}^2 + \sigma_{x_2}^2}{\sigma_{e_1}^2 + \sigma_{e_2}^2} = \frac{1 + \frac{\sigma_{x_2}^2}{\sigma_{x_1}^2}}{\frac{\sigma_{e_1}^2}{\sigma_{x_1}^2} + \frac{\sigma_{e_2}^2}{\sigma_{x_1}^2}} \cong SNR_1^I \quad (4.10)$$

where $\sigma_{e_1}^2$ is the noise variance due to the first section compression technique and $\sigma_{e_2}^2$ is the noise variance due to the second section compression technique.

Next, the compression ratio can be expressed as using Equation (4.4),

$$CPR^I = \frac{B}{\frac{B_1}{CPR_1^I} + \frac{B_2}{CPR_2^I}}$$

$$= \frac{B \text{ CPR}_1 \cdot \text{CPR}_2}{B_1 \cdot \text{CPR}_2 + B_2 \cdot \text{CPR}_1}$$

where $B = B_1 + B_2$. Variables B_1 and B_2 are the bit rates of the first and second sections, and CPR_1 and CPR_2 are expected compression ratios for the corresponding two sections.

These two sections will be compressed separately. The first section will be compressed using the predictive coding method based on linear prediction analysis, and the second section will be compressed using the hybrid technique discussed earlier. In the next section, predictive coding is discussed with respect to the algorithms for an optimum linear predictor.

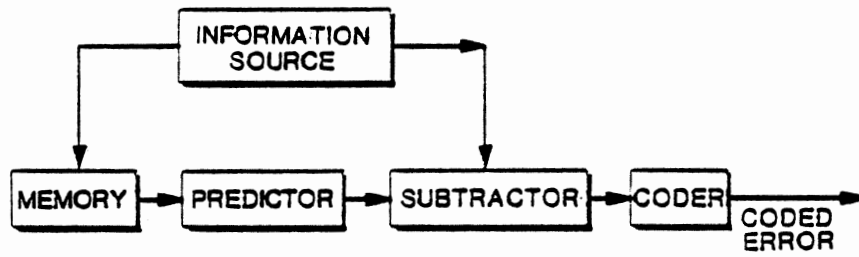
Linear Prediction Analysis

Predictive coding simply corresponds to the differential coding with an optimum predictor, which is based upon Wiener's work [26]. Let P_i be the predicted value of the i th message m_i . Then, the prediction error is given by

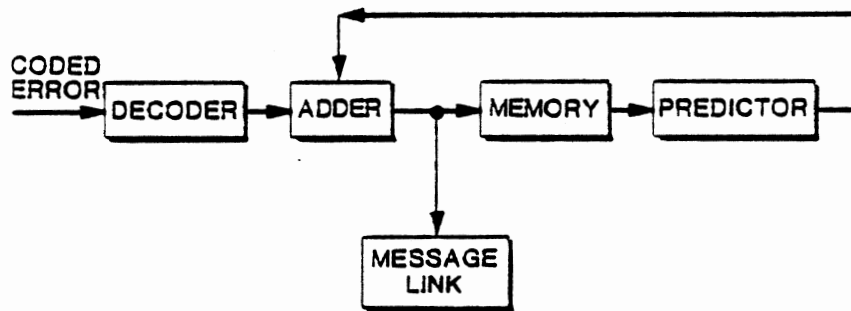
$$e_i = m_i - P_i. \quad (4.12)$$

Since P_i is a deterministic value, it has no information according to Shannon's definition. Thus by transmitting the error term, e_i , no information is lost.

The predictive coding and decoding procedures are illustrated in Figure 27, where an identical predictor appears in the transmitter and the receiver. Since the predictor operates on the past values of the message, storage of the past values is necessary. The size of this storage, or simply the order of the prediction, is an important parameter and is discussed later.



(a) TRANSMITTER



(b) RECEIVER

Figure 27. "Predictive Coding" Procedures
(After Alias [26])

The predicted value P_i is a linear combination of the previous message values. That is,

$$P_i = \sum_{j=1}^{\infty} a_j m_{i-j} \quad (4.13)$$

where the a_j 's are some constants to be determined by minimizing the root mean square value of e_i . Wiener has shown that this predictor is determined not by the message ensemble, but by the autocorrelation function of the ensemble [26]. In general, there will be many ensembles with the same autocorrelation function, and all of these will have the same linear predictor.

For obvious reasons, Equation (4.13) cannot be implemented with infinite sum, and P_i is given below with finite sum

$$P_i = \sum_{j=1}^P a_j m_{i-j} \quad (4.14)$$

where P corresponds to the order of the filter.

For this case, there are various formulations, such as the covariance method [27], the autocorrelation method [28], the lattice method [29], the inverse filter formulation [20], the spectral estimation formulation [30], the maximum-likelihood formulation [31], and the inner product formulation [23]. The most simple method is the autocorrelation method, which uses Durbin's recursive solution [28].

Consider the autocorrelation equations [23]

$$\sum_{k=1}^P \alpha_k R(i-k) = R(i) \quad 1 \leq i \leq p \quad (4.15)$$

$$R(i) = \sum_{m=0}^{N-i-1} X(m) X(m+i) \quad (4.16)$$

where $R(i)$ is the i th autocorrelation lag, and N is the number of data points. The solution of Equation (4.15), by Durbin's method, is given by

$$E^{(0)} = R(0) \quad (4.17)$$

$$k_i = (R(i) - \sum_{j=1}^{i-1} \alpha_j^{(i-1)} R(i-j)) / E^{(i-1)} \quad (4.18)$$

$$\alpha_j^{(i)} = k_i \quad (4.19)$$

$$\alpha_j^{(i)} = \alpha_j^{(i-1)} - k_i \alpha_{i-j}^{(i-1)} \quad (4.20)$$

$$E^{(i)} = (1 - k_i^2) E^{(i-1)} \quad (4.21)$$

where E is the variance of the prediction error, $\alpha_j^{(i)}$ is the j th parameter of the i th order predictor, and i is contained in the closed interval $[1, P]$ while j is contained in the closed interval $[1, i-1]$. It should be noted that

$$\alpha_j = \alpha_j^{(p)}$$

which implies that α_j of the p th order predictor is equal to α_j of an i th order predictor where $i \leq p$.

The most popular method of implementation of a linear predictor is by the Lattice method, which uses k_i in Equation (4.19). The variable k_i is often referred to as partial correlation coefficients or PARCOR coefficients [29], and k_i 's can be computed recursively. It is guaranteed to yield a stable filter without requiring the use of a window [6].

The Lattice method computes the prediction error as follows [32].

The prediction error sequence, $e^{(i)}(n)$, can be expressed as

$$e^{(i)}(m) = e^{(i-1)}(m) - k_i b^{(i-1)}(m-1). \quad (4.22)$$

Recursively, the i th stage backward prediction error is computed from

$$b^{(i)}(m) = b^{(i-1)}(m-1) - k_i e^{(i-1)}(m) \quad (4.23)$$

where

$$e^{(0)}(m) = b^{(0)}(n) = s(m).$$

This procedure is illustrated in Figure 28.

Implementation of "Predictive Coding"

It was discussed earlier that impulsive seismic data are divided into two sections due to the energy distribution characteristics. The predictive coding technique is used for the impulsive section in order to achieve better signal-to-noise ratio by taking advantage of the prediction gain discussed in Chapter III (see Equation (3.8)). For the second section, referred to as non-impulsive section, the hybrid technique with a 6-bit μ -law quantizer is used for the DCT coefficients. The details of the hybrid technique were discussed earlier.

In the following implementation, considerations of predictive coding are discussed. The predictor coefficients are derived from the linear prediction analysis using Durbin's autocorrelation method. The residual signal is computed from the Lattice method [29]. It should be noted that the prediction parameters and the PARCOR coefficients are computed only

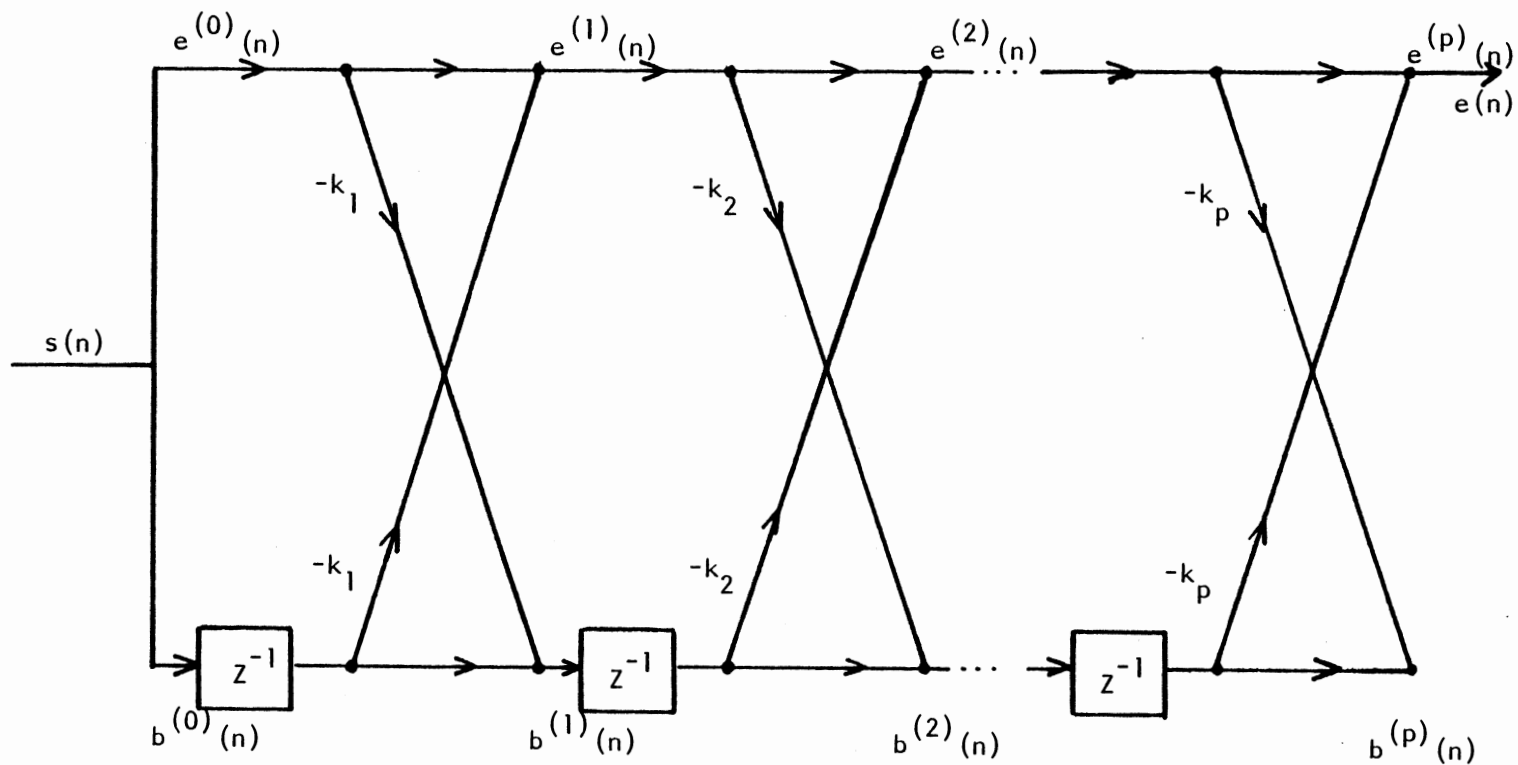


Figure 28. Block Diagram of Implementation of the lattice Method
 (After Rabiner and Schafer [6])

once and shared by all traces. Also, the PARCOR coefficients are transmitted to the receiver for synthesizing the original seismic data.

Table VII gives ten parameters corresponding to a tenth-order linear predictor for eight impulsive seismic traces. It can be seen that the first five parameters for each trace are approximately same. It was found from empirical result that a fifth-order predictor is sufficient. Using the Lattice method residual signal, or simply prediction error, is computed where the PARCOR coefficients are obtained from the fifth-order linear predictor. For the residual signal coding, a 6-bit μ -law quantizer is used. These two coded sections are edited into one array for transmission. Table VIII illustrates the quantization signal-to-noise ratios of prediction method comparing with LPCM and 6-bit μ -law coding. It is shown that μ -law coding method approaches the result of prediction method by 1 dB.

At the receiver, the coded residual signal and the coded transformed coefficients are separately decoded. The decoded residual signal is synthesized and the impulsive section is reconstructed. For reconstructing the nonimpulsive section, the hybrid technique is used. Then, both reconstructed sections are concatenated for obtaining the full trace of the impulsive seismic data. These procedures are illustrated in Figure 29. The impulsive sections and their residual signals are shown in Figure 30, where the prediction gain can be easily noted.

The compression results are evaluated with respect to the signal-to-noise ratio and the compression ratio. The SNR of two sections are computed from

$$\text{SNR}_1^1 = E[(X_1 - X_1')^2] / E[(X_1^2)]$$

TABLE VII
PREDICTION PARAMETERS

trace	1	2	3	4	5	6	7	8	9	10
1	-0.594	0.781	-0.697	0.764	-0.583	0.629	-0.267	0.061	0.199	-0.050
2	-0.572	0.768	-0.714	0.731	-0.551	0.540	-0.193	0.166	0.165	-0.077
3	-0.680	0.770	-0.728	0.738	-0.555	0.619	-0.268	0.210	0.043	-0.127
4	-0.719	0.726	-0.741	0.730	-0.560	0.698	-0.202	0.182	0.116	-0.029
5	-0.627	0.836	-0.736	0.752	-0.637	0.653	-0.342	0.179	0.028	-0.090
6	-0.656	0.769	-0.755	0.686	-0.600	0.624	-0.344	0.453	-0.216	-0.191
7	-0.673	0.773	-0.742	0.674	-0.493	0.367	-0.101	-0.101	0.024	0.058
8	-0.755	0.751	-0.712	0.763	-0.567	0.574	-0.329	0.291	-0.146	0.118

TABLE VIII
RESULTS OF IMPULSIVE SEISMIC DATA
COMPRESSION/DECOMPRESSION

TRACE	IMPULSIVE SECTION			NON-IMPULSIVE SECTION	OVER-ALL CPR
	LPCM	μ -LAW	PREDICTION	HYBRID	
1	27.16	28.56	30.09	34.62	6.48
2	27.40	28.62	28.57	30.76	5.47
3	27.17	29.43	30.44	30.17	5.99
4	27.08	29.02	30.30	31.04	6.13
5	27.43	29.58	30.44	32.53	6.83
6	25.49	28.32	29.92	32.21	6.82
7	27.95	30.01	30.37	32.37	6.63
8	28.11	29.83	30.40	31.83	5.61
AVG	27.22	29.18	30.06	31.94	6.25

IMPULSIVE SEISMIC DATA COMPRESSION

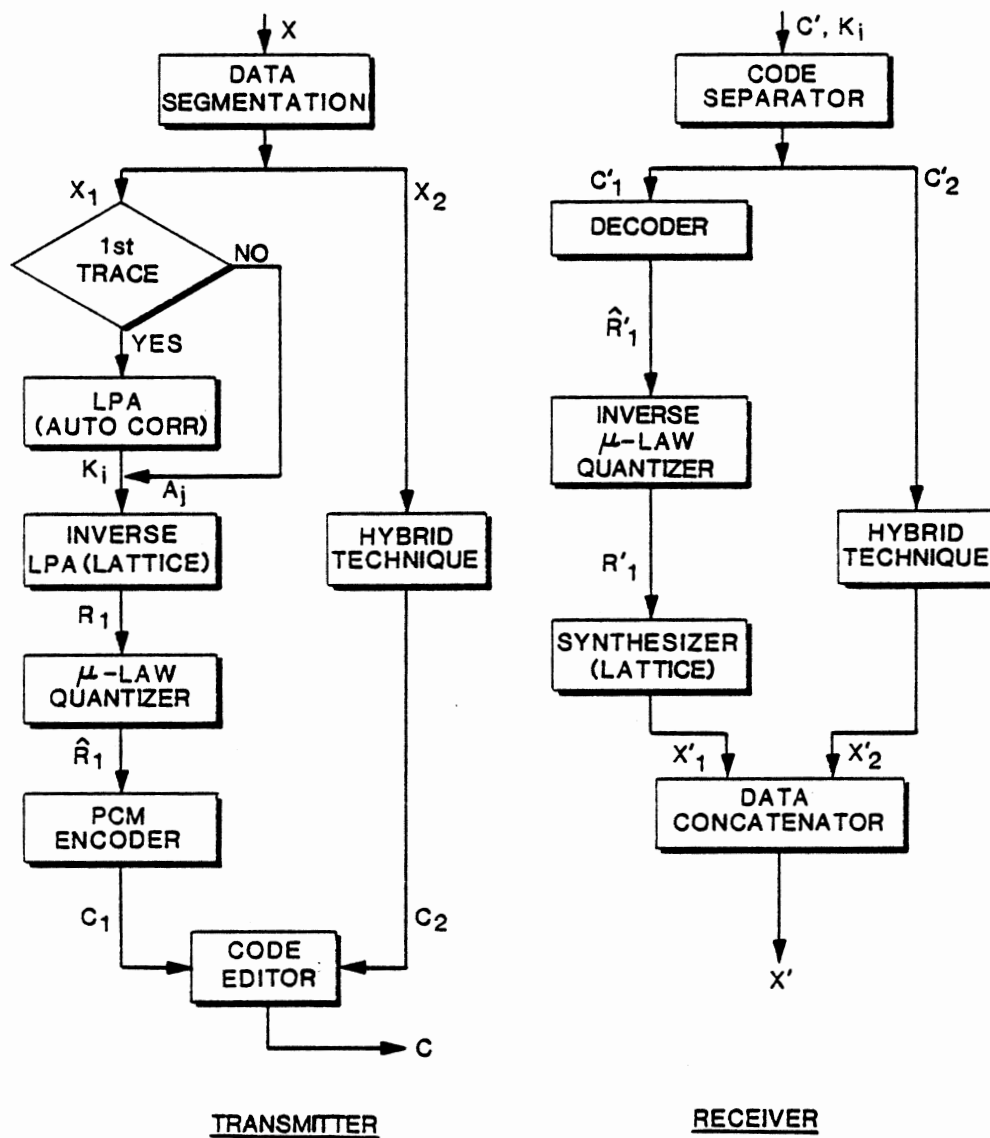
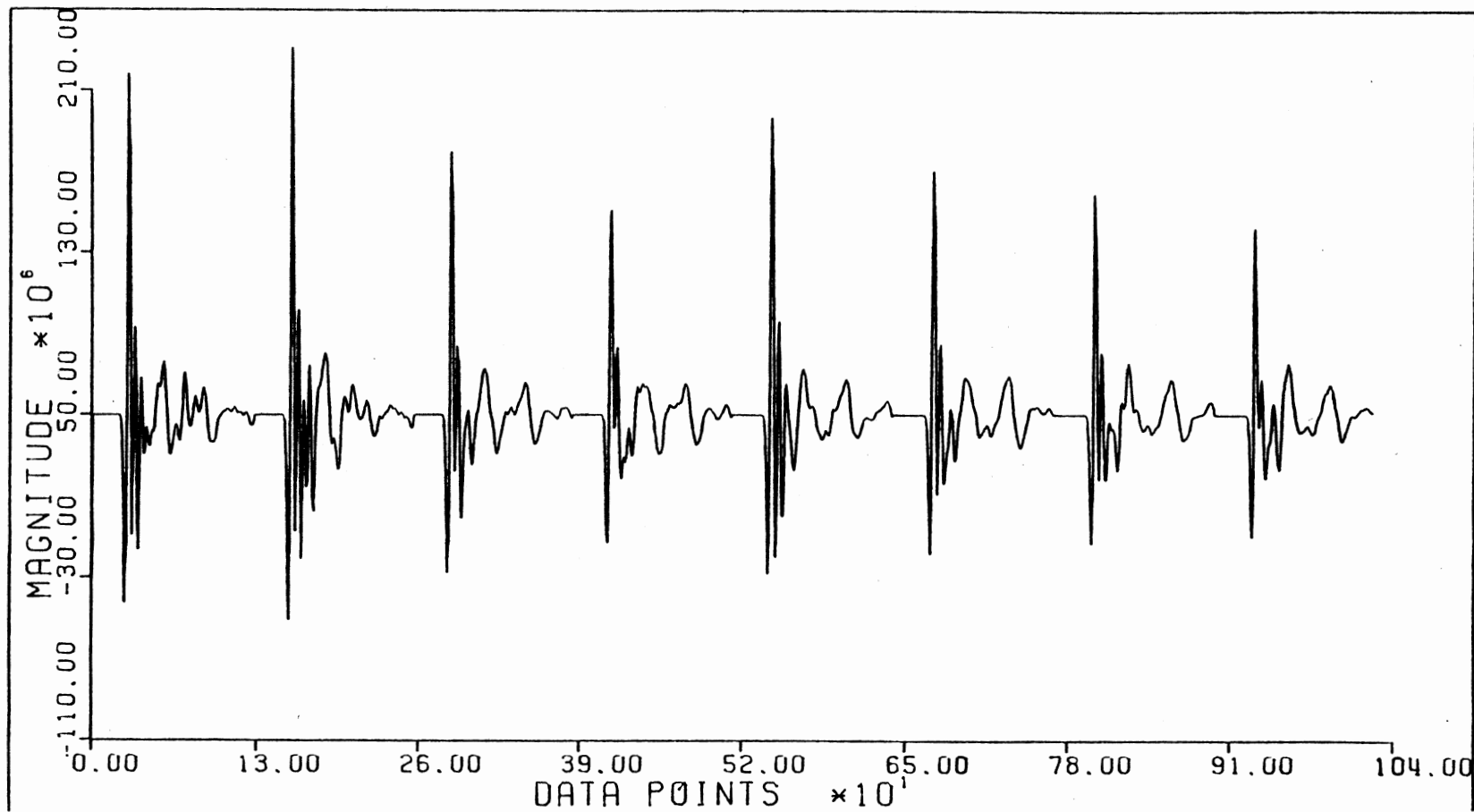
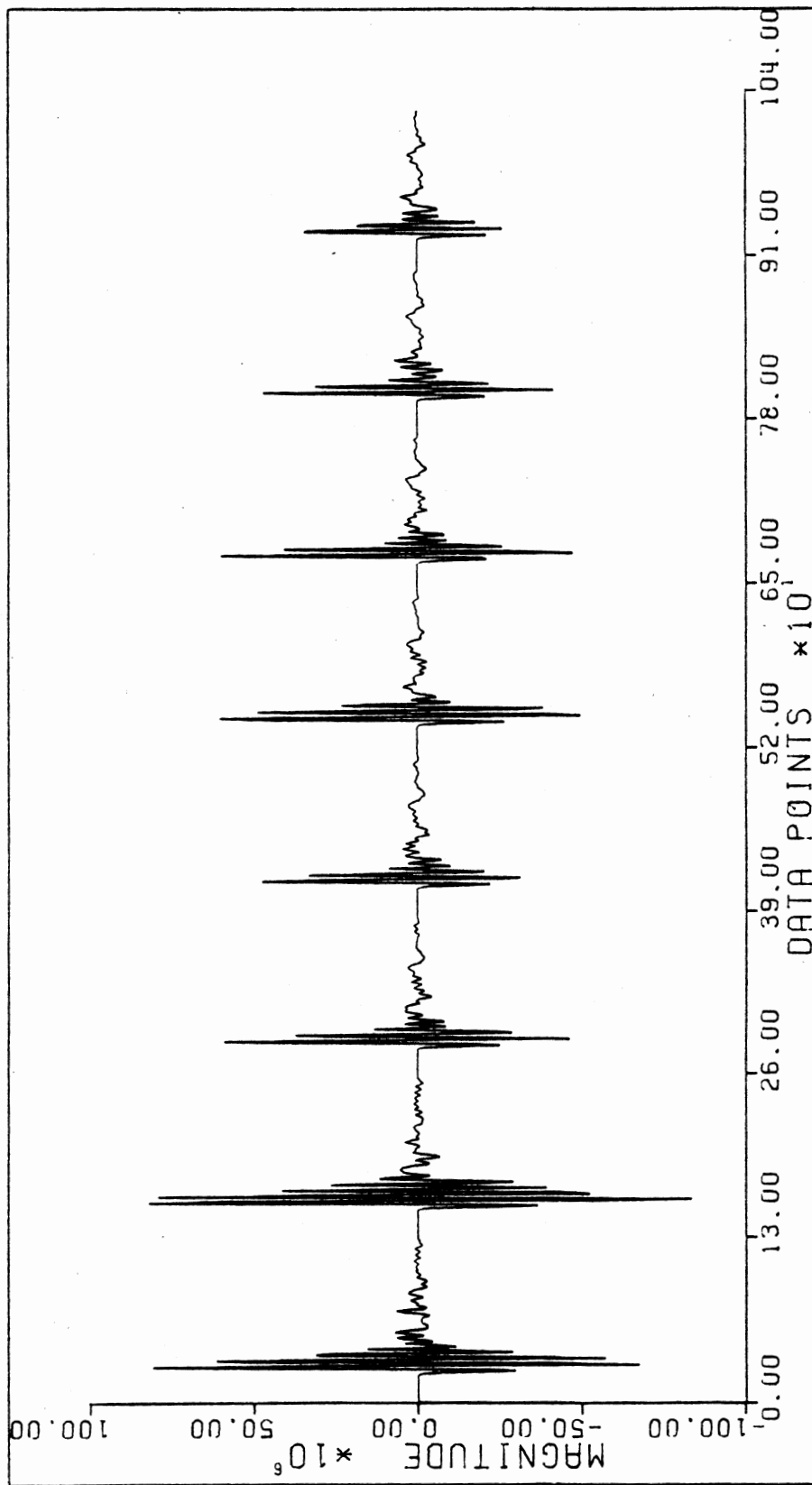


Figure 29. Impulsive Seismic Data Compressive/Decompressive Procedures Using "Predictive Coding" and "Hybrid Technique"



(a) Impulsive Section of eight impulsive seismic data traces

Figure 30. Illustration of Prediction residual signal of Impulsive Section



(b) Prediction Residual Signal of Eight Impulsive Sections

Figure 30, (Continued)

$$\text{SNR}_2^I = E[(X_2 - X_2')^2]/E[(X_2^2)]$$

where X_1 is the original impulsive section, X_2 is the original non-impulsive section, and X_1' and X_2' are their reconstructed data.

The compression ratio is computed from

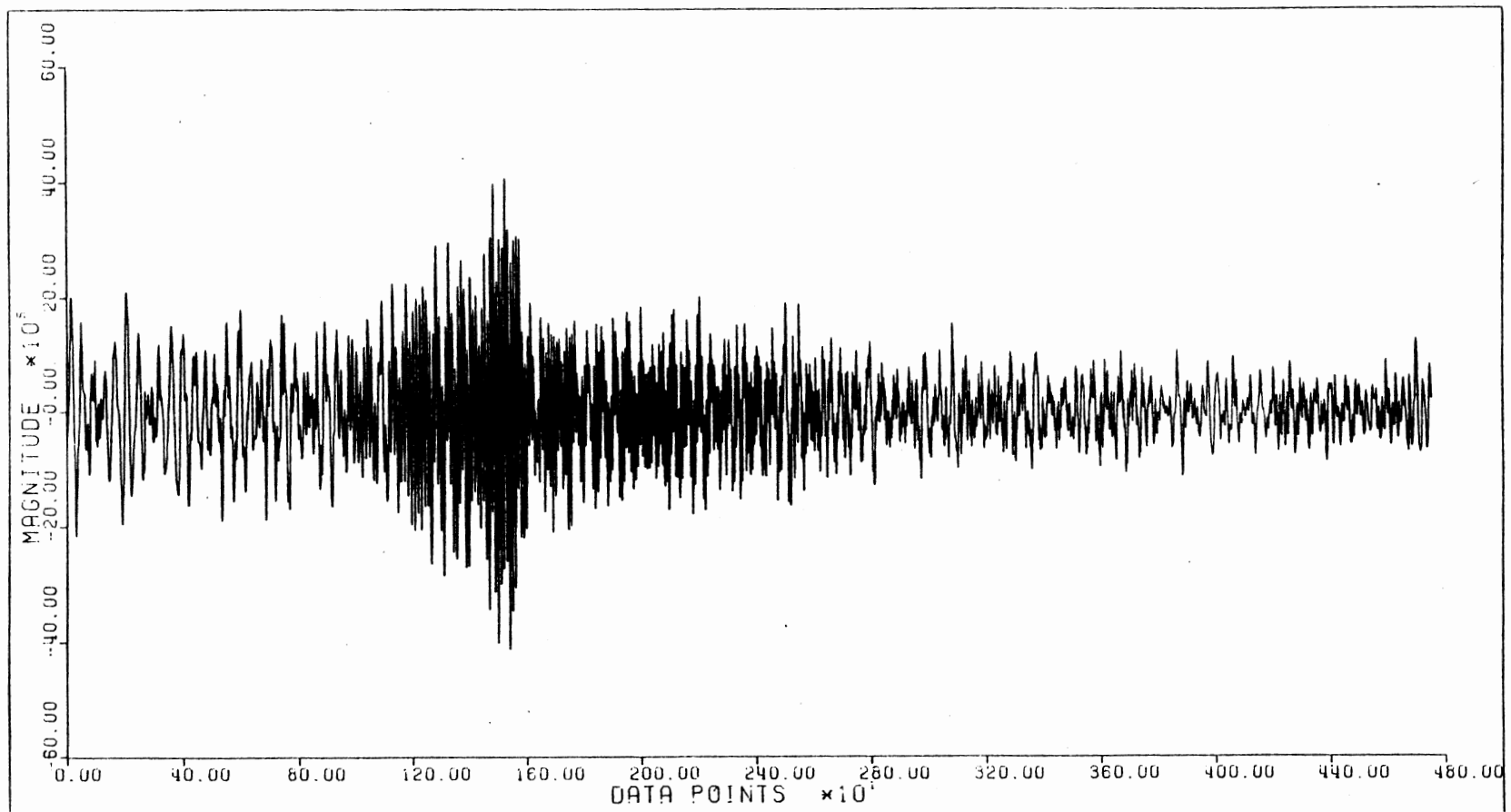
$$\text{CPR} = (N_1 \times 6 + M \times 6 + N_2)/(1000 \times 20)$$

where N_1 is the number of data in the impulsive section, N_2 is the number of data in the non-impulsive section, and M is the number of the retained coefficients of the non-impulsive section. N_2 is included as a bookkeeping information array. SNR_1^I , SNR_2^I , SNR^I , and CPR are illustrated in Table VIII, and the average compression ratio is approximately 6.25 to 1.00 for the impulsive seismic data at hand.

Examples

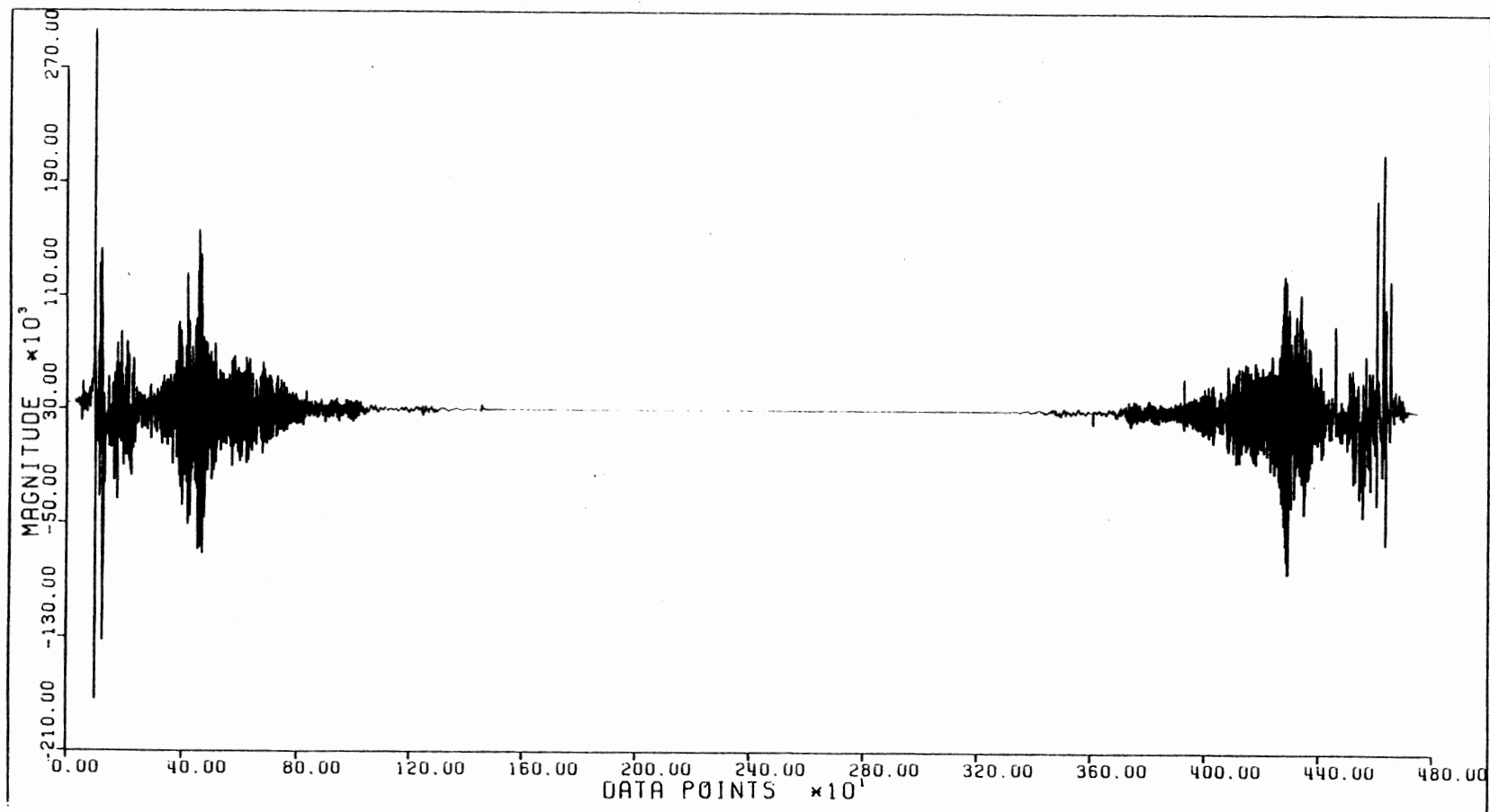
Compression techniques for vibroseis data and impulsive seismic data have been developed and their performances were evaluated earlier. In order to visualize these compression procedures on seismic data, Figures 31 and 32 illustrate outputs of each functional block of the hybrid technique and predictive coding shown in Figures 24a, 24b and 27. Figure 31a is the plot of one vibroseis trace with 4750 data points, Figure 31b is the DCT coefficients of the vibroseis trace, Figure 31c is the selected DCT coefficients reduced to 2216 points, and Figure 31d is the μ -law quantized and a 7-bit PCM coded result of the selected DCT coefficients. Exact inverse procedures are performed (see Figure 24b) at the receiver and Figure 31e illustrates the reconstructed vibroseis trace.

The impulsive seismic data case is illustrated in Figure 32. Figure 32a is the plot of one impulsive seismic trace with 1000 data points.



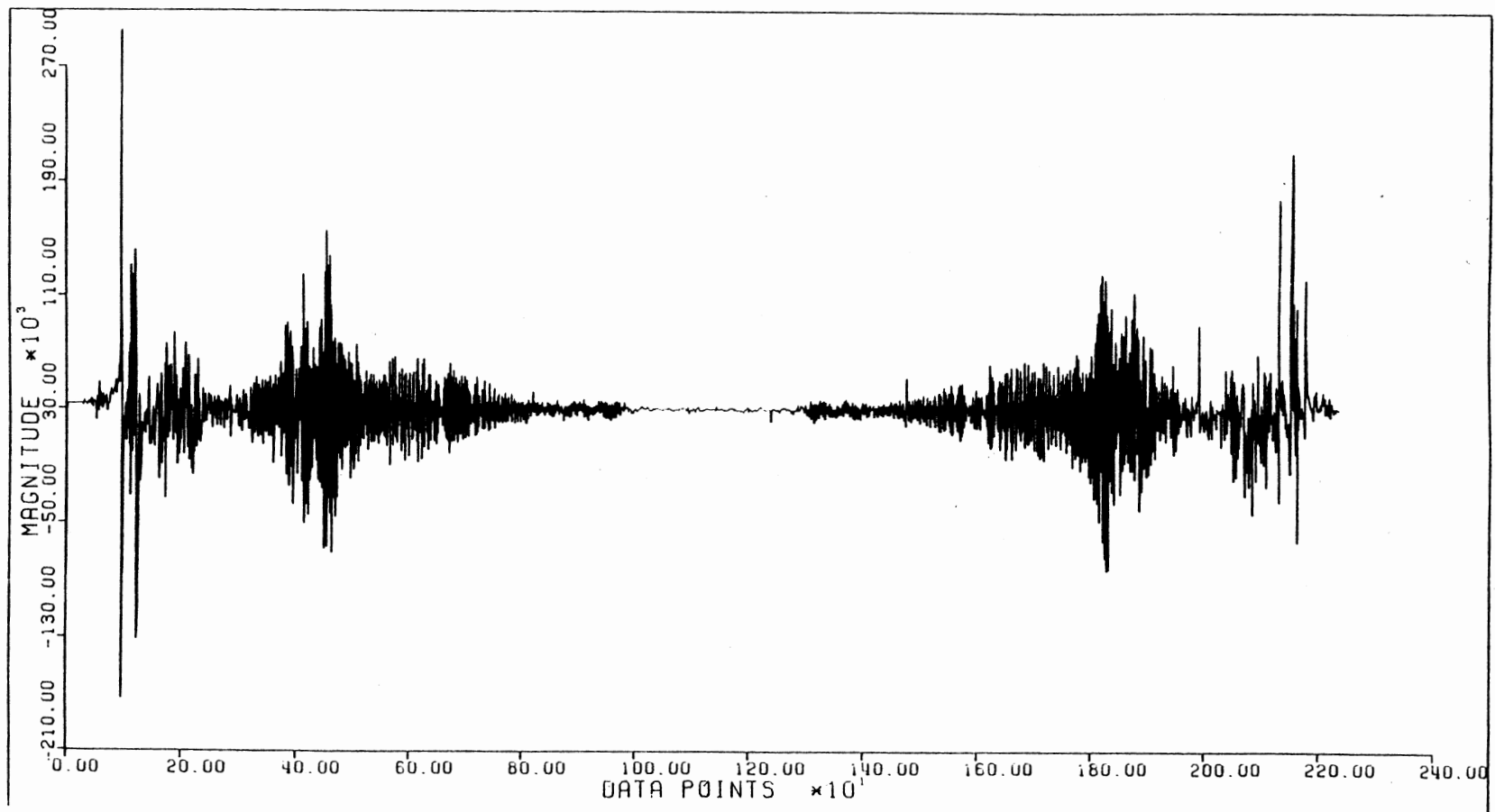
(a) One Trace of Vibroseis Data

Figure 31. Illustration of Vibroseis Data Compression/Decompression



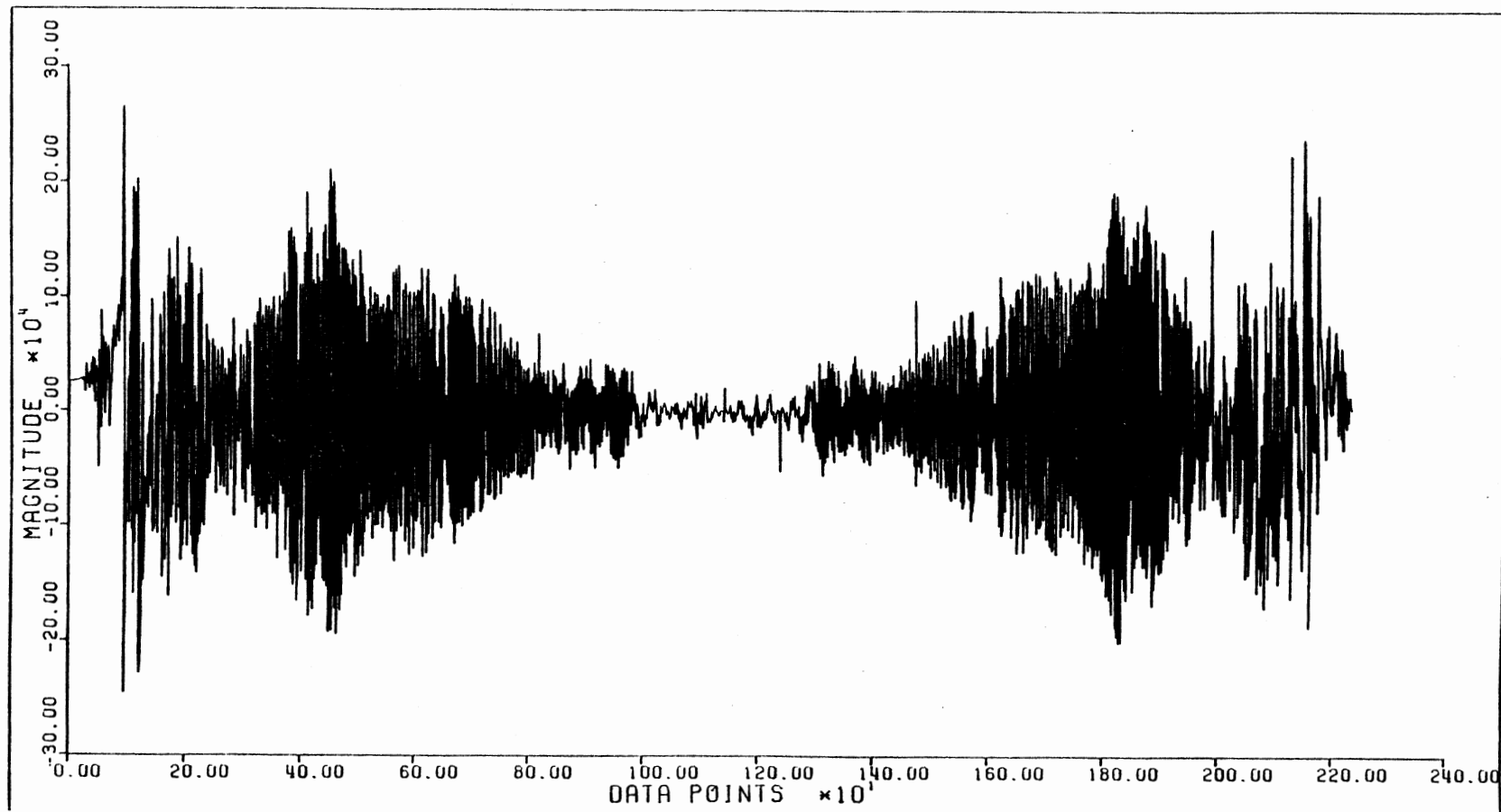
(b) Discrete Cosine Transformed Vector Coefficients

Figure 31. (Continued)



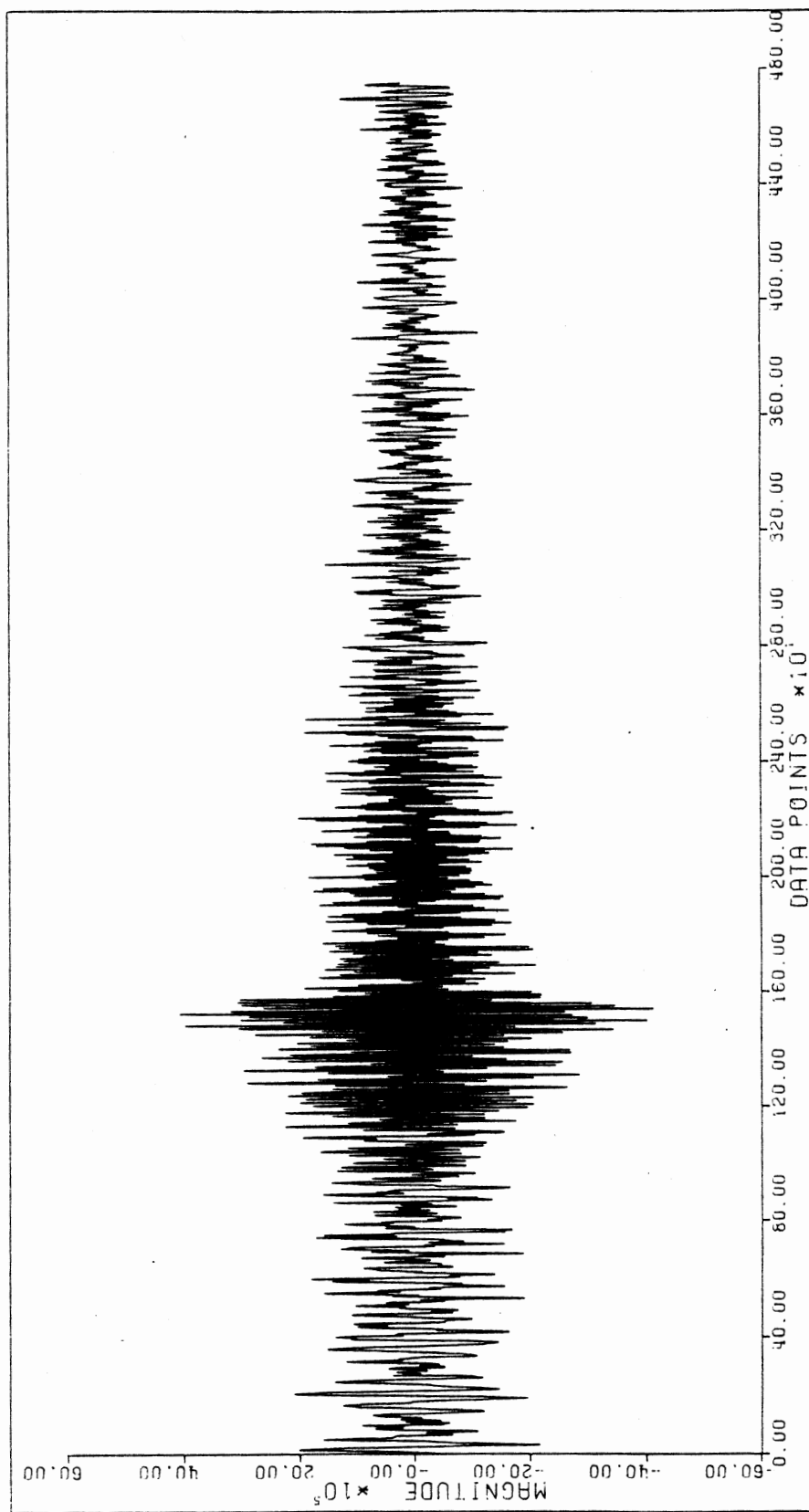
(c) Selected Discrete Cosine Vector Transformed Coefficients

Figure 31. (Continued)



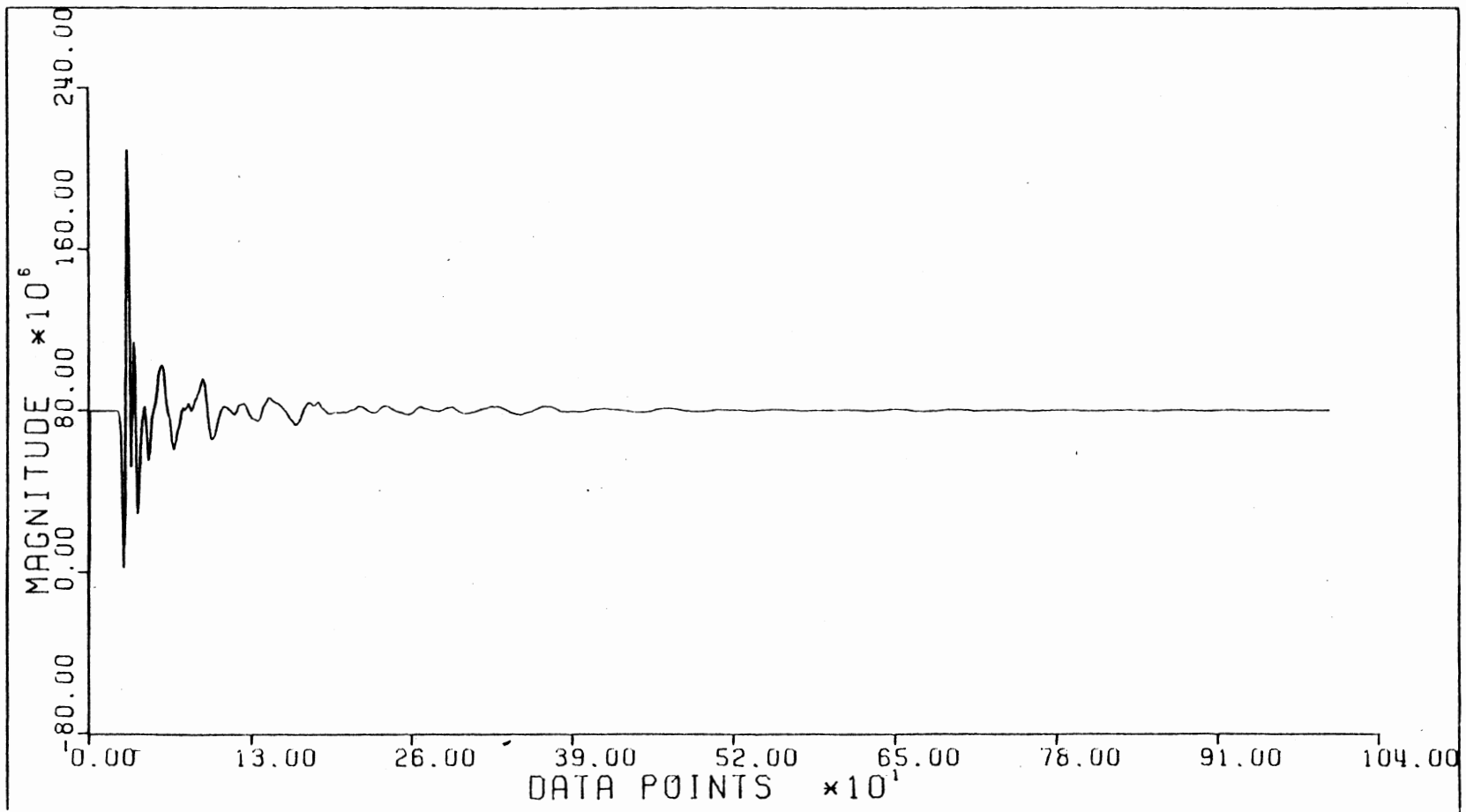
(d) μ -law Quantized Output of the Selected DCT Coefficients

Figure 31. (Continued)



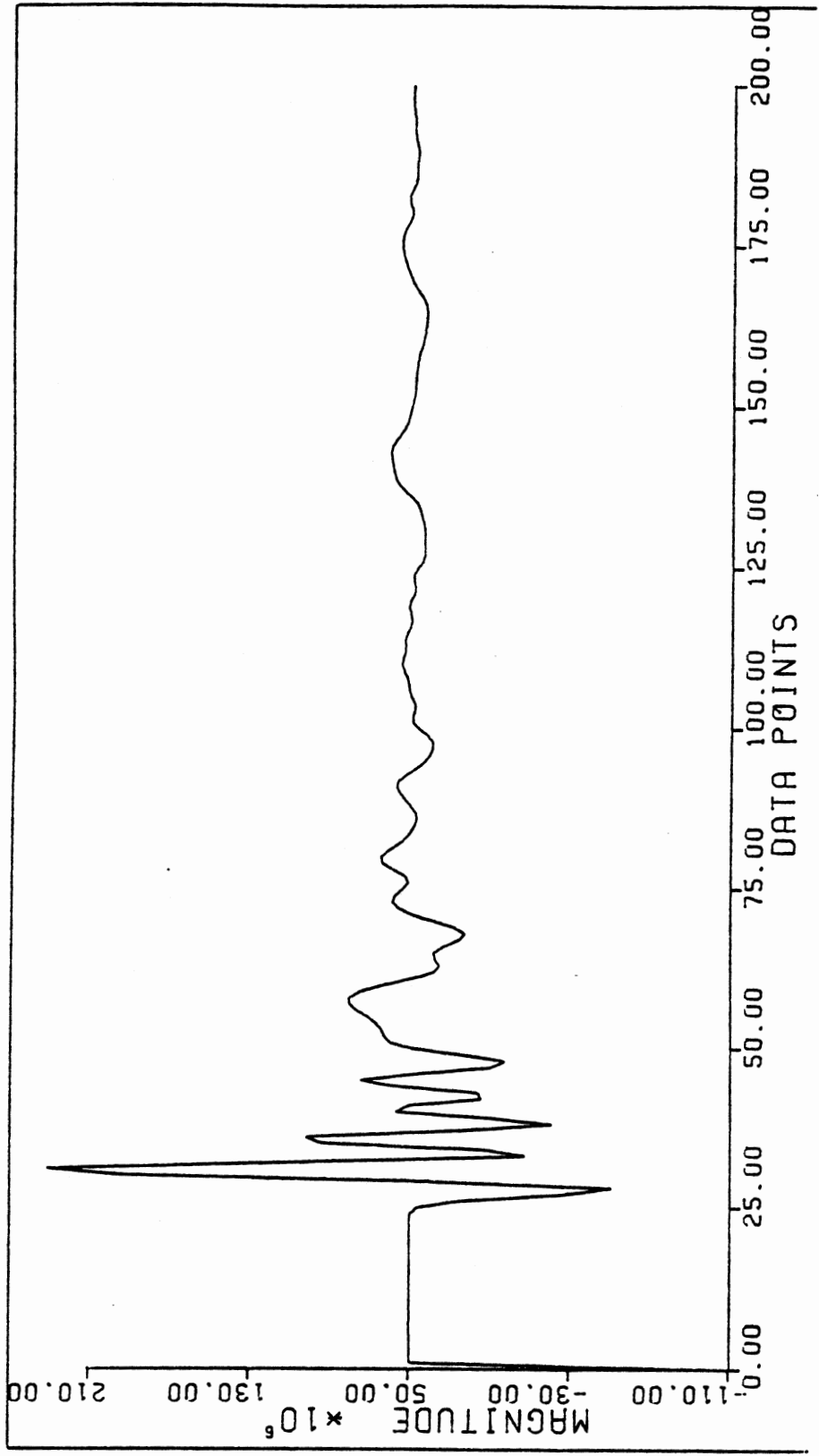
(e) Reconstructed Vibroseis Data Trace

Figure 31. (Continued)



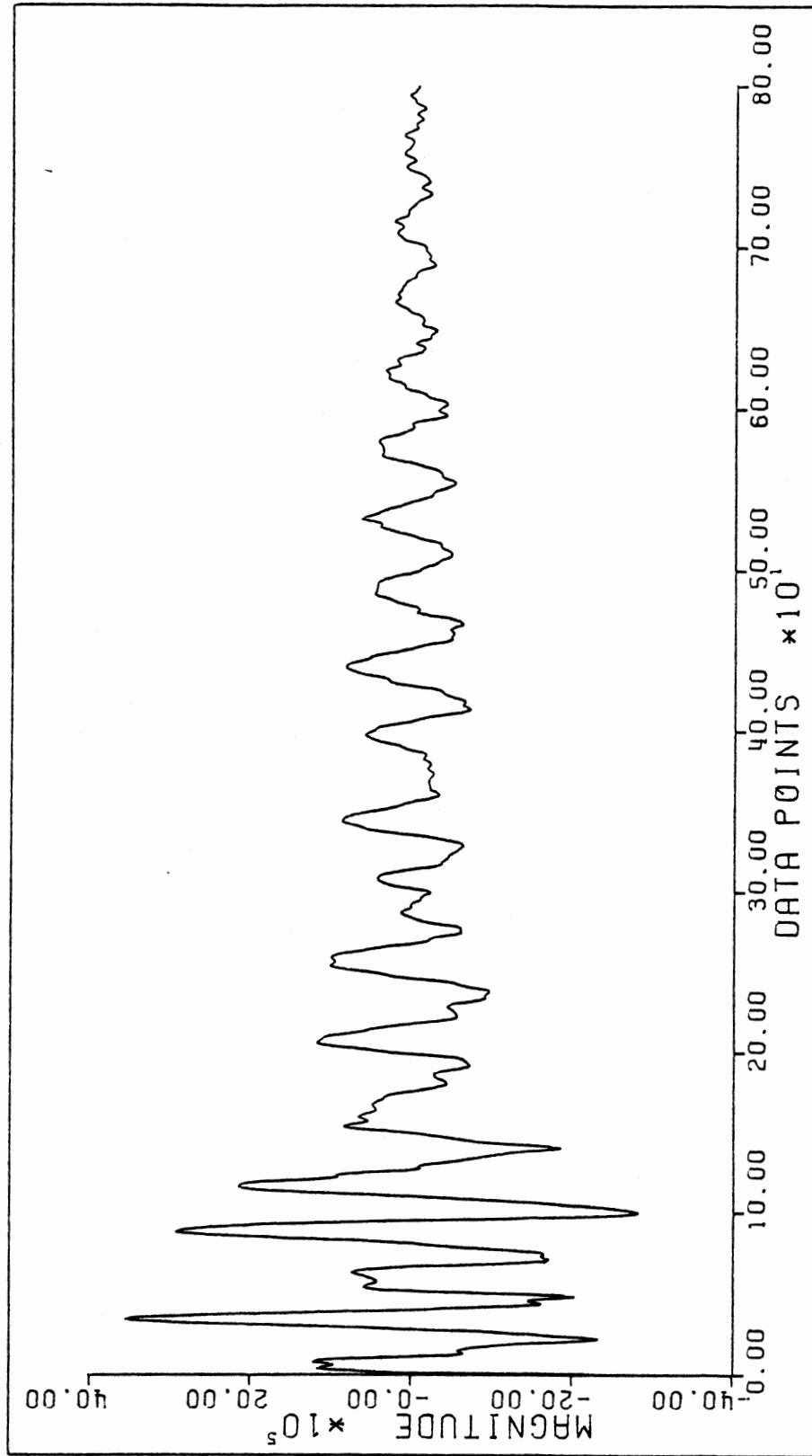
(a) One Trace of Impulsive Seismic Data

Figure 32. Illustrations of Impulsive Seismic Data Compression/Decompression



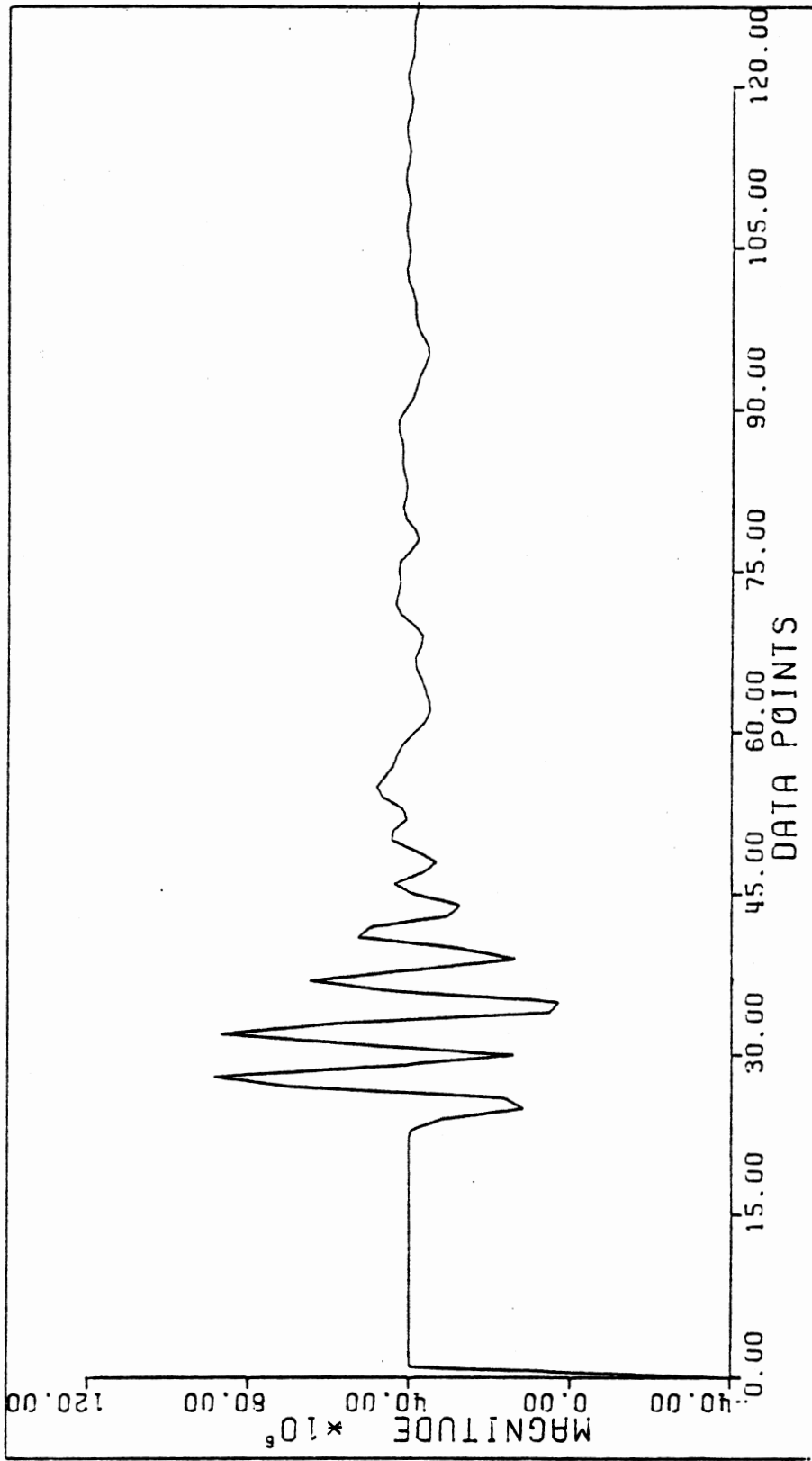
(b) Impulsive Section of Impulsive Seismic Data

Figure 32. (Continued)



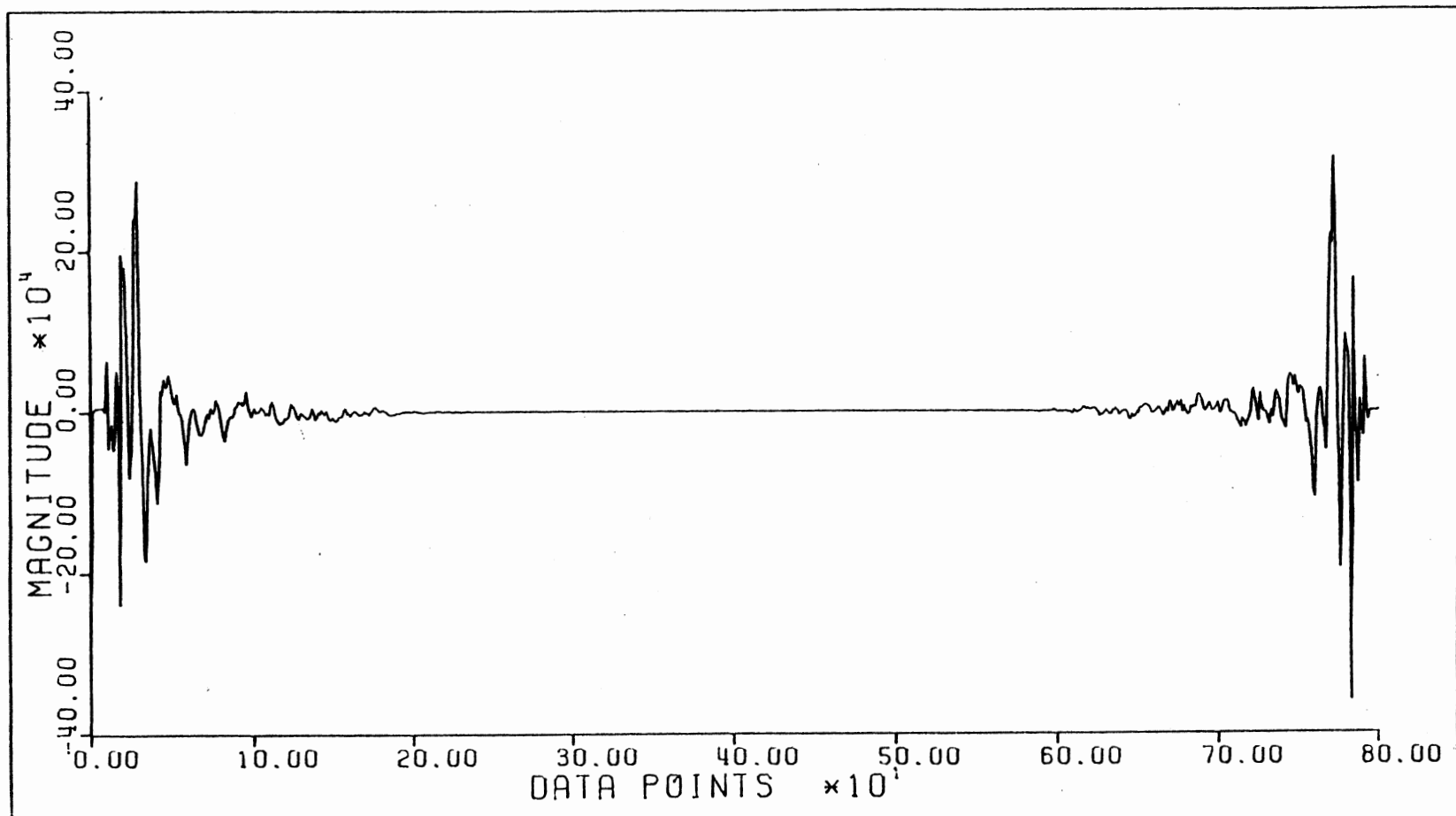
(c) Nonimpulsive Section of Impulsive Seismic Data

Figure 32. (Continued)



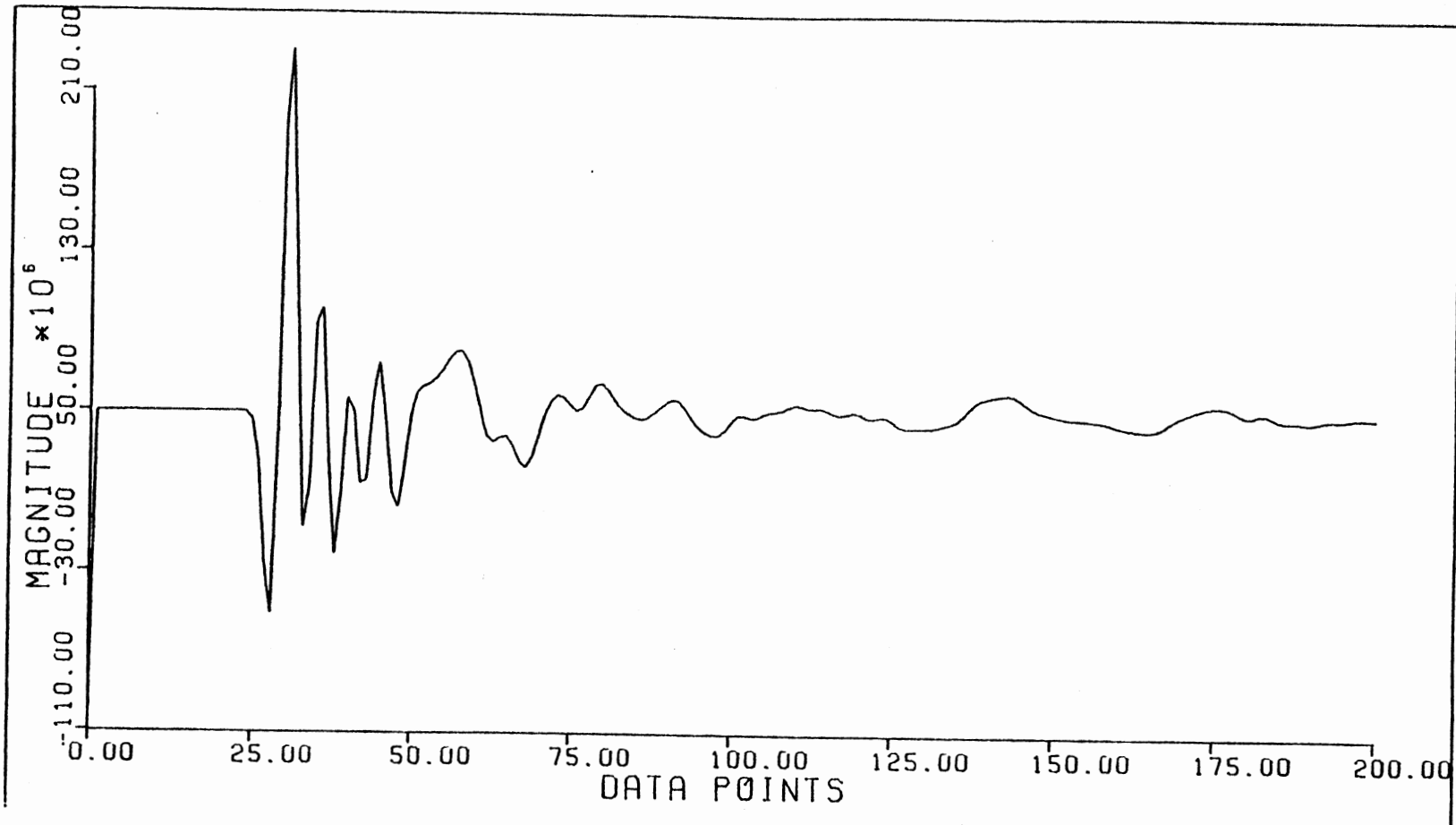
(d) Prediction Residual Signal of Impulsive Section

Figure 32. (Continued)



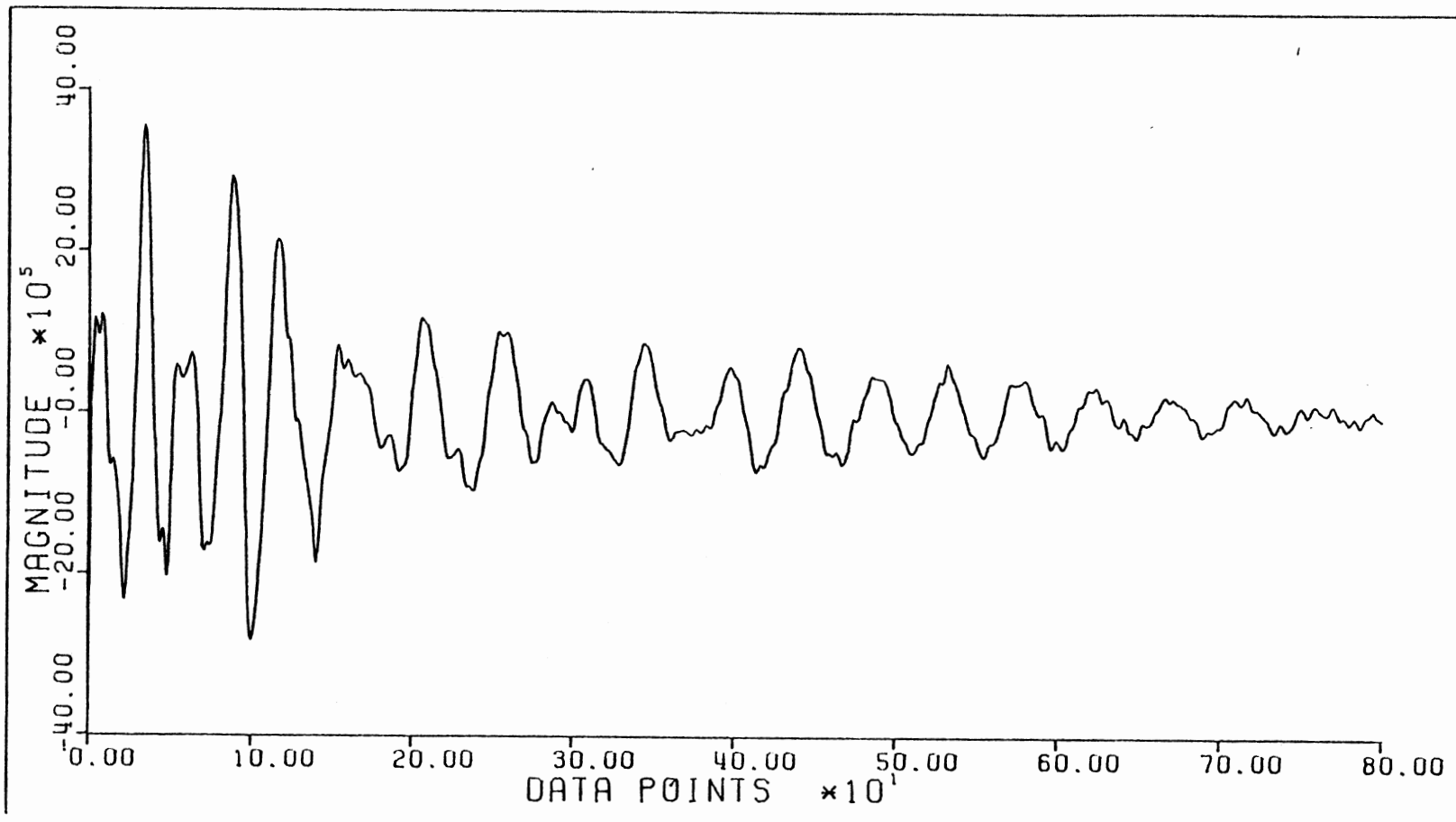
(e) Discrete Cosine Transformed Vector Coefficients of Non-impulsive Section

Figure 32. (Continued)



(f) Reconstructed Impulsive Section of Impulsive Seismic Data

Figure 32. (Continued)



(g) Reconstructed Nonimpulsive Section of Impulsive Seismic Data

Figure 32. (Continued)

This trace is segmented into two sections, impulsive and nonimpulsive, as illustrated in Figures 32b and 32c. Figure 32d shows the residual signal of the impulsive section, and 32e is the DCT coefficients of nonimpulsive section. Selection of the significant DCT coefficients and coding of the selected coefficients are performed in a similar manner as the vibroseis data case. The reconstructed data of the both sections at the receiver are shown in Figures 32f and 32g, respectively.

From Figures 31 and 32, it can be seen that the reconstructed waveforms have no noticeable distortion. Furthermore, the objective data given in terms of SNR's in Table VI and VIII indicate that 30 dB SNR requirement is, in general, satisfied for the data in hand.

In the next chapter, simulation program steps will be evaluated with respect to their execution time and core size requirement for real-time implementation considerations.

CHAPTER V

SEISMIC DATA TRANSMISSION VIA SATELLITE

Introduction

Most satellite systems provide a variety of bandwidths so that the interface equipment for the terrestrial links can also be used for satellite circuits. In Chapter II, the basic idea of seismic data transmission via satellite was proposed due to the bandwidth constraints of channels, storage limits, and transmission costs, data compression has been considered. Specific compression methods have been developed for vibroseis data and impulsive seismic data in Chapter IV.

This chapter considers some important aspects of the compressed seismic data transmission over a satellite channel. Figure 33 illustrates a digital communication system [23]. This includes several functional blocks, such as source encoder/decoder, channel encoder/decoder, modulator/demodulator, and the communication channel. Satellite systems provide a variety of services so that the user can interface with this system properly. This chapter discusses only the portion of the overall system that needs to be interfaced.

First, the source encoder and decoder, the seismic data compressor and decompressor in this thesis, are examined with respect to their execution time and core size requirement. These are evaluated from the simulation program of the hybrid technique and predicting coding discussed in

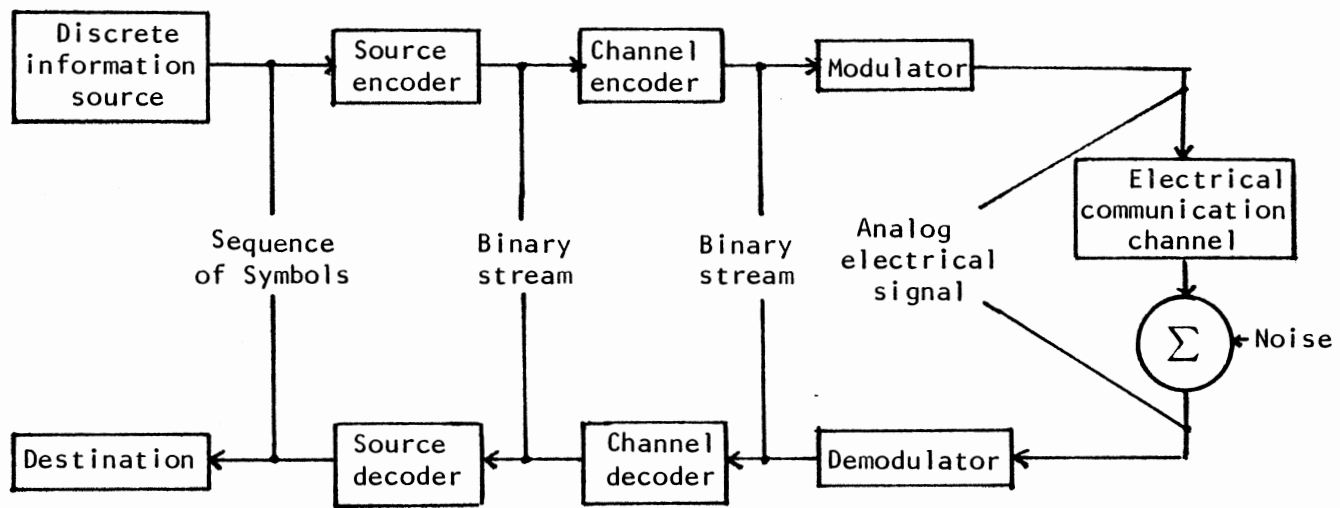


Figure 33. Digital Communication System
(After Shanmougam [23])

Chapter IV. Also, some important factors of real-time system design considerations are discussed.

Second, access methods such as frequency division multiple access and time division multiple access of the satellite channel are explored. Finally, network protocols for satellite channels are examined with two types of link control procedures, half-duplex and full-duplex, and their transmission efficiencies are evaluated.

Time and Space Configuration

In Chapter IV, the hybrid technique and predictive coding have been developed for compressing vibroseis data and impulsive seismic data. These techniques have been simulated in FORTRAN language with the IBM 370. All FORTRAN subroutines that are used in this thesis are included in Appendix C. The hybrid technique involves seven subroutines: FDCT (fast discrete cosine transform), SELECT (significant DCT coefficients selector), CONMEU (μ -law quantization), ENCODE (PCM coder), DECODE (mapping and decoding), INVMEU (inverse μ -law quantization), and FIDCT (fast inverse DCT). The first four subroutines are used for compression at the transmitter, and the remaining three subroutines are used for compression at the receiver for vibroseis data and/or nonimpulsive section of impulsive seismic data. The predictive coding involves three subroutines: AUTO (autocorrelation method for LPA), INVERS (lattice method for inverse LPA), and SYNTHZ (synthesizer), in addition to four coding and decoding subroutines: CONMEU, ENCODE, DECODE, and INVMEU. AUTO is used only for the first trace in order to obtain prediction parameters. INVERS, CONMEU, and ENCODE are used for compression at the transmitter, and DECODE, INVMEU, and SYNTHZ are used for decompression of the impulsive section of the

impulsive seismic data. In the following, these subroutines are evaluated with their execution time and core size requirements.

First, all the routines mentioned above are compiled with the FORTRAN H compiler and cross-referenced with IBM 370 assembler language (see Appendix C). Since execution time of each instruction may vary among computer systems, the execution time is expressed as a combination of instruction types rather than absolute numbers for the IBM 370 case.

Second, the instructions are divided into six different types, such as load/store, multiply/divide, add/subtract, branch, compare, and external subroutine calls. These divisions are based on characteristics of instructions, and it should be pointed out that the differences of execution time within one type of instruction, such as load, load address, and load resistor, are ignored for simplicity. The execution time of the above six types are denoted as t_l , t_m , t_a , t_b , t_c , and t_x , respectively.

Third, instructions inside "DO" loops are evaluated rather than entire program steps, since the execution time is mostly occupied by "DO" loops. This assumption will not introduce a great deal of error in time estimation if the number of iterations of a "DO" loop is high.

For illustrative purposes, these procedures are implemented for a subroutine CONMEU. The source listing and cross-referenced assembler listing of CONMEU are shown in Table IX. It can be seen that there are five-load, three-store, three-multiply, two-divide, one-add, and three-branch instructions inside "DO" loop 10. The "DO" loop is iterated as many times as the value of the variable NUM. The variable NUM indicates the number of selected DCT coefficients in this case. Thus, the execution time can be estimated as

TABLE IX
COMPILED LISTING OF SUBROUTINE "CONMEU"

```

/ STRUCTURED SOURCE LISTING /
(002) ISN 0002      SUBROUTINE CONMEU(RDATA,NLN,EPSLCN,XMAX,RMEU)      00000910
C                  MEU-LAW CONVERTER                          00000920
C                                                            00000930
C      RDATA      :   INPUT ARRAY DCT COEFFICIENTS OF SEISMIC DATA TRACE  00000940
C      RMEU       :   MEU-VALUE                                           00000950
C      XMAX       :   ABSOLUTE MAXIMUM VALUE OF RDATA ARRAY              00000960
C      NLN        :   NO. OF DATA POINTS PER SWEEPS                     00000970
C      EPSLCN     :   THRESHOLD VALUE OF INSIGNIFICANT DCT COEFFICIENTS  00000980
C                                                            00000990
      ISN 0003      DIMENSION RDATA(1)                                  00001000
C                  APPLY MEU-LAW                                         00001010
      ISN 0004      RMEU=RMEU+1.0                                         00001020
      ISN 0005      DEN=ALG(RMEU)                                          00001030
      ISN 0006      EPSLCN=EPSLCN/XMAX                                     00001040
      ISN 0007      EPSLCN=1.0+RMEU*EPSLCN                                 00001050
      ISN 0008      EPSLCN=XMAX*ALG(EPSLCN)                               00001060
      ISN 0009      EPSLCN=(EPSLCN/DEN)                                    00001070
      ISN 0010      DO 10 I=1,NLN                                         00001080
(001) ISN 0011      RDATA(I)=0.0                                         00001090
      ISN 0012      SIGN=1.0                                              00001100
      ISN 0013      ABSR=ABS(RDATA(I))                                    00001110
      ISN 0014      IF (RDATA(I).LT.0.0) SIGN=-1.0                       00001120
      ISN 0016      ABSR=ABSR/XMAX                                         00001130
      ISN 0017      AESF=1.0+RMEU*AESR                                     00001140
      ISN 0018      ABSR=XMAX*ALG(ABSR)                                    00001150
      ISN 0019      ABSR=(ABSR/DEN)*SIGN                                   00001160
      ISN 0020      RDATA(I)=AESF                                         00001170
      ISN 0021      10 CONTINUE                                          00001180
(001) C                                                                C
(002) ISN 0022      20 RETURN                                           00001190
      C                                                                C
      ISN 0023      END                                                  00001200

```

TABLE IX (Continued)

*LEVEL 2.3.0 (JUNE 78)		CS/360 FCRTAN F EXTENDED	DATE 81.153/21.05.18	
000000	47 F0 F 00C	CCNMEU EC	15,12(0,15)	
000004	07	CC	XL1'07'	
000005	C3CED5DACEE440	DC	CL7'CCNMEU'	
00000C	50 EC D 00C	STM	14,12,12(13)	
000010	1E 4D	LR	4,13	
000012	58 CD F 020	LM	12,12,32(15)	
000016	E0 40 D C04	ST	4,4(0,13)	
00001A	50 D0 4 00B	ST	13,8(0,4)	
00001E	07 FC	EXR	15,12	
CONSTANTS				
000080	00C00000	DC	XL4'CCCC00C0'	
000084	00C00001	DC	XL4'00000001'	
000088	41100000	DC	XL4'4110C000'	
00008C	00000000	DC	XL4'00000000'	
000090	00C00000	DC	XL4'00000000'	
ADCONS FOR VARIABLES AND CONSTANTS				
ADCONS FOR EXTERNAL REFERENCES				
0000C8	00C00000	DC	XL4'CCCC0000'	ALCG
0000CC	00000000	DC	XL4'00000000'	RDATA
0000EC	5E 80 D 0A4	100001 L	E, 164(0,13)	
0000F0	58 A0 D 0B4	L	10, 180(0,13)	4
0000F4	5E 70 D 074	L	7, 116(0,13)	NUM
0000F8	78 00 D C60	LE	0, 96(0,13)	41100000
0000FC	7A 00 D 07C	AE	C, 124(0,13)	FMEL
000100	70 00 D 088	STE	0, 136(0,13)	RMEU1
000104	58 F0 D CA0	L	15, 160(0,13)	ALCG
000108	41 10 D 04C	LA	1, 76(0,13)	
00010C	05 EF	EALR	14,15	
00010E	47 00 0 005	BC	C, 5(0, 0)	
000112	70 00 D 0BC	STE	C, 186(0,13)	0100
000116	70 00 D C70	STE	C, 112(0,13)	CEM
00011A	78 20 D 090	LE	2, 144(0,13)	EPSLON
00011E	70 20 D CE4	DE	2, 132(0,13)	XMAX
000122	70 20 D C50	STE	2, 144(0,13)	EPSLON
000126	7C 20 D 07C	ME	2, 124(0,13)	RMEU
00012A	7A 20 D 060	AE	2, 96(0,13)	41100000
00012E	70 20 D 090	STE	2, 144(0,13)	EPSLON
000132	58 F0 D 0A0	L	15, 160(0,13)	ALCG
000136	41 10 D C50	LA	1, 80(0,13)	

TABLE IX (Continued)

00013A	05 EF		EALR	14,15	
00013C	47 00 0 C08		EC	0, 8(0, 0)	
000140	78 20 D C84		LE	2, 132(0,13)	XMAX
000144	3C 20		MER	2, 0	
000146	70 20 D 090		STE	2, 144(0,13)	EPSLEN
00014A	7D 20 D 08C		DE	2, 188(0,13)	0T00
00014E	70 20 D 090		STE	2, 144(0,13)	EPSLEN
000152	78 00 D C80		IE	0, 96(0,13)	41100000
000156	33 00		LCER	(, 0)	
000158	70 00 D 080		STE	0, 176(0,13)	0C02
00015C	18 B7		LR	11, 7	
00015E	89 E0 0 002		SLL	11, 2	
000162	18 5A		LF	5,10	
000164	78 00 D C58	100002	LE	0, 88(0,13)	0
000168	70 09 8 000		STE	0, 0(9, 8)	FCATA
00016C	78 00 D C80		LE	0, 96(0,13)	41100000
000170	70 00 D 080		STE	0, 128(0,13)	SIGN
000174	78 29 8 000		LE	2, 0(5, 8)	FCATA
000178	30 02		LPER	0, 2	
00017A	70 00 D C7E		STE	0, 120(0,13)	ABSF
00017E	32 22		LIER	2, 2	
000180	47 A0 D 164		EC	10, 256(0,13)	100004
000184	78 00 D 080	100003	LE	0, 176(0,13)	0C02
000188	70 00 D 080		STE	0, 128(0,13)	SIGN
00018C	78 20 D C78	100004	LE	2, 120(0,13)	ABSF
000190	70 20 D 084		DE	2, 132(0,13)	XMAX
000194	7C 20 D C7C		ME	2, 124(0,13)	RNEL
000198	7A 20 D C60		AE	2, 96(0,13)	41100000
00019C	70 20 D 078		STE	2, 120(0,13)	AESF
0001A0	58 F0 D 0A0		L	15, 160(0,13)	ALOG
0001A4	41 10 D 054		LA	1, 84(0,13)	
0001A8	05 EF		EALR	14,15	
0001AA	47 00 0 012		EC	0, 18(0, 0)	
0001AE	78 20 D C84		LE	2, 132(0,13)	XMAX
0001B2	3C 20		MER	2, 0	
0001B4	7D 20 D C70		DE	2, 112(0,13)	DEA
0001B8	7C 20 D 080		ME	2, 128(0,13)	SIGN
0001BC	70 29 8 000		STE	2, 0(5, 8)	FCATA
0001C0	87 9A D 13C	10	EXLE	9,10, 316(13)	100002
0001C4	18 FF	20	SR	15,15	
0001C6	58 E0 D 000		L	14, 0(0,13)	
0001CA	07 FE		ECR	15,14	

TABLE IX (Continued)

ADDRESS OF EPILOGUE					
0001CC	58 A0 D 004	L	10,	4(0,13)	
0001D0	58 E0 A 00C	L	14,	12(0,10)	
0001D4	58 B0 A 018	L	11,	24(0,10)	
0001D8	5E 10 B C08	L	1,	8(0,11)	
0001DC	78 20 D 090	LE	2,	144(0,13)	EPFLCN
0001E0	70 20 1 000	STE	2,	0(0, 1)	
0001E4	18 DA	LR	13,10		
0001E6	92 FF A 00C	MVI	12(10),255		
0001EA	5E 2C A 01C	LM	2,12,	28(10)	
0001EE	07 FE	BCR	15,14		
ADDRESS OF PROLOGUE					
0001F0	98 7A 1 C04	LM	7,10,	4(1)	
0001F4	58 20 7 000	L	2,	0(0, 7)	
0001F8	50 20 D C74	ST	2,	116(0,13)	NUM
0001FC	78 2C 8 000	LE	2,	0(0, 8)	
000200	70 20 D 090	STE	2,	144(0,13)	EPFLCN
000204	7E 20 9 000	LE	2,	0(0, 9)	
000208	70 20 D 084	STE	2,	132(0,13)	KMAX
00020C	78 20 A 000	LE	2,	0(0,10)	
000210	7C 20 D C7C	STE	2,	124(0,13)	PMEL
000214	58 20 1 000	L	2,	0(0, 1)	
000218	41 30 2 000	LA	3,	0(0, 2)	
00021C	41 50 0 C04	LA	5,	4	
000220	1E 25	SR	2,	5	
000222	50 20 D CA4	ST	2,	164(0,13)	
000226	50 30 D 0A8	ST	3,	168(0,13)	BLATA
00022A	47 F0 D 0C4	BC	15,	196(0,13)	

$$E_t(\text{CONMEU}) = \text{NUM} \cdot (8t_l + 5t_m + t_a + 3t_b). \quad (5.1)$$

In a similar manner, execution time for the remaining routines can be evaluated. The number of occurrences of each type of instruction set inside the "DO" loops for the remaining subroutines are given in Table X. The subroutines FDCT and FIDCT call FFT subroutines of IMSL (International Mathematical Subroutine Library), FFTRC, and FFTCC. Since the FFT function can be performed by a dedicated processor and there are numerous software packages of FFT, the execution time of these particular FFT routines are not evaluated here. For FDCT, the execution time can be expressed as

$$E_t(\text{FDCT}) = \text{NUM} \cdot (37t_l + 2t_m + 5t_a + 7t_b) + t_x, \quad (5.2)$$

where the time t_x is solely dependent on a chosen FFT algorithm processor, and NUM is the number of data points to be transformed.

From these analyses, total execution time of the hybrid technique at the transmitter, that is, a compression only, can be computed from

$$\begin{aligned} TE_t^H(\text{Tr}) &= E_t(\text{FDCT}) + E_t(\text{SELECT}) + E_t(\text{CONMEU}) \\ &\quad + E_t(\text{ENCODE}). \end{aligned} \quad (5.3)$$

Similarly, at the receiver, corresponding to the decompression, the execution time is given by

$$TE_t^H(\text{Re}) = E_t(\text{DECODE}) + E_t(\text{INVMEU}) + E_t(\text{FIDCT}) \quad (5.4)$$

Also, the execution time for the predictive coding at the transmitter is given by

$$\begin{aligned} TE_t^P(\text{Tr}) &= E_t(\text{AUTO}) + E_t(\text{INVERS}) + E_t(\text{CONMEU}) \\ &\quad + E_t(\text{ENCODE}). \end{aligned} \quad (5.5)$$

TABLE X
 FREQUENCY LIST OF 6-TYPE INSTRUCTIONS
 INSIDE "DO" LOOP

SUBROUTINE	LOAD/ STORE	MULT/ DIV.	ADD/ SUB.	BRANCH	COMP.	EXT. CALL
FDCT	37	2	5	7		FFTRC
SELECT	18	1	8	6	2	
CONMEU	15	5	1	3		
ENCODE	20	2	7	5	1	
DECODE	8	1	1	2	1	
INVMEU	22	3	2	6		
FIDCT	26	1	3	6		FFTCC
AUTO	114	6	36	12	6	
INVERS	89	10	10	11		
SYNTHZ	94	10	20	11		

At the receiver, the execution time for reconstructing the impulsive section from the coded residual signal is given by

$$TE_t^P(\text{Re}) = E_t(\text{DECODE}) + E_t(\text{INVMEU}) + E_t(\text{SYNTHZ}). \quad (5.6)$$

It should be pointed out that CONMEU, ENCODE, DECODE, and INVMEU, which are used in both the predictive coding and the hybrid technique, are functions of the iteration value. Also, note that predictive coding is applied only for the impulsive section of an impulsive seismic trace. The subroutine AUTO will be used only for the first trace of impulsive seismic data, and therefore Equation (5.5) can be rewritten as

$$TE_t^P(\text{Tr}) = E_t(\text{INVERS}) + E_t(\text{CONMEU}) + E_t(\text{ENCODE}). \quad (5.7)$$

Finally, the core size requirement can be obtained by examining the address range of the assembler listing of each routine. For example, the address range of CONMEU, in terms of base 16, is from 0_{16} to $22E_{16}$, or, in terms of base 10, from 0_{10} to 558_{10} . The unit of memory is, in general, a byte (8-bit) and the address range shows that the subroutine CONMEU needs 558 bytes of memory size. Table XI illustrates the core size requirement and the number of program statements for all routines involved in the hybrid technique and predictive coding.

For estimating the required memory size of the computer system for implementing these techniques, the following need to be considered. If all routines are loaded simultaneously in the memory, the memory size needs to be a sum of all required core sizes of subroutines and a main routine which handles subroutines, that is,

$$C = \sum_{i=a}^x C_i + C(\text{Main}), \quad (5.8)$$

TABLE XI
MEMORY SIZE REQUIREMENT

SUBROUTINE	CORE SIZE (BYTES)	PROGRAM STATEMENT
FDCT	59582	16
SELECT	914	56
CONMEU	558	22
ENCODE	558	22
DECODE	418	11
INVMEU	560	23
FIDCT	59460	13
AUTO	1026	44
INVERS	662	19
SYNTHZ	694	19

where C_i is the core size for the subroutine i . However, in the proposed method, one routine is used at one time and the remaining routines are stored in the secondary storage. With this assumption, the memory size can be computed from

$$C = \text{MAX} (C_a, C_b, \dots, C_x) + C(\text{MAIN}). \quad (5.9)$$

For example, the required core size of the hybrid technique for compressing the vibroseis data trace (4750 data points) is given by

$$\begin{aligned} C^H(\text{Tr}) &= \text{MAX} (914, 558, 59582, 558) + 38260 \\ &= 97842 \text{ (bytes)}. \end{aligned} \quad (5.10)$$

The large core size requirement for the FDCT and FIDCT can be explained by the working area that FFT routines use. This working area can be significantly reduced by using chirp z-transform (CZT) algorithm [9]. The size of the MAIN routine is a function of the data points and it includes all global variable storages, such as input array, code array, and book-keeping array. It can be concluded that the total core size requirement may be down to 64k-byte for implementing the compression technique for vibroseis trace if CZT is used for FDCT and FIDCT.

Real-Time Design Considerations

The definition of real time is obviously application dependent and can vary anywhere from milliseconds to hours [32]. For the proposed seismic data transmission via satellite, general relationships among processing time, transmission time, and depth of queue are investigated to some extent. The term "queue" indicates in general the place where a job waits to be serviced by the system. There may be several queues in one system, such as input queue waiting for CPU or output queue waiting for printer or

display terminal. In this section, the queue is mainly used for the data to be transmitted, as shown in Figure 34.

The processing time is a flexible factor and very much dependent on the structure of the system, such as multiprocessor structure, dedicated processor, or hardware multiplier system, and so forth. The response time, in general, dictates the type of processing. The processing time of the hybrid technique using a multiprocessor system can be obtained from Equation (5.3), and is expressed as

$$T_p = TE_t^H (Tr)/N , \quad (5.11)$$

where N is the number of processors in the system.

The depth of queue can be computed from [33]

$$Q_d = \text{Input rate} - \text{Output rate (bits/second)} \quad (5.12)$$

where the input rate is the incoming data rate (bits/second) into the queue and is given by

$$\begin{aligned} \text{Input rate} &= \text{Original seismic data rate/CPR} \cdot T_p \\ &= \text{Compressed seismic data rate}/T_p , \end{aligned} \quad (5.13)$$

and the output rate is the data rate which is taken out of the queue. It is given by

$$\text{Output rate} = \text{MIN (Input rate [Speed of circuit} \cdot e]), \quad (5.14)$$

where e is the transmission efficiency and is defined by [34]

$$e = \frac{\text{Data transmission time}}{\text{Total transmission time}} . \quad (5.15)$$

The speed of circuit indicates the bandwidth of the channel. Satellite circuits include network protocols which control satellite channels. It

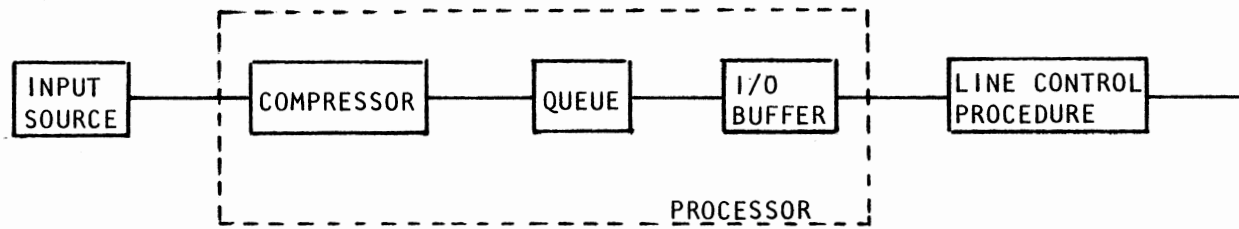


Figure 34. Data Compression and Transmission System

is clear that the speed of the circuit multiplied by the transmission efficiency should be less than the input rate in order that the channel bandwidth is used efficiently. From Equations (5.13), (5.15), and (5.15) the depth of queue in Equation (5.12) can be rewritten as

$$Q_d = \frac{\text{Compressed seismic data rate}}{T_p} - (\text{Channel bandwidth} \cdot e). \quad (5.16)$$

The transmission time (T_t) of the compressed seismic data is computed from

$$T_t = \frac{\text{Compressed seismic data}}{(\text{Channel bandwidth} \cdot e)}. \quad (5.17)$$

The response time of the system can be obtained from

$$R = \text{MAX} (T_p, T_t), \quad (5.18)$$

provided there is enough space in the queue so that processing time can be independent from the transmission time. It indicates that the processing time may be as long as the transmission time without degrading the response time.

The above considerations have not taken into account the effects of transmission delay and/or the overhead control bits to be transmitted. In the next section, detail of the satellite links and constraints of block size in relation to the bandwidth will be discussed. Also, the transmission efficiency will be examined with two types of line control procedures.

Satellite Communication

Satellite links have the following general properties [34]. There

is approximately 270 msec of propagation delay and transmission cost is independent of distance within the range of a satellite. A signal sent to a satellite is transmitted to all receivers within the range of the satellite antenna, simply referred to as broadcast property. Also, a satellite provides large bandwidths and uses digital transmission. The broadcast property of satellites may cause a serious security problem. However, it can be economic for transmitting data to geographically dispersed places.

Access methods of satellite links are different from terrestrial links due to their propagation delay and large bandwidth. Some of the basic terms associated with satellite communications are defined in Appendix D.

Network Protocols

Network protocols are sets of rules that govern the flow of data in a network. This involves automatic error detection and correction as well as recovery procedures [35]. In general, protocols are divided into three levels. The lowest-level protocol is the hardware level, such as hardware interfaces, where "handshaking" sequences can be achieved. The high-level protocol, referred to as a link control procedure, is a set of rules that ensures a block of data gets from one end of a data link to the other without errors. The highest-level protocol is another set of rules related to message flow. This level interacts with the line control procedure for complete message reception. This protocol is often referred to as a network handler.

A system which detects an error in data and has those data automatically retransmitted is called ARQ (automatic repeat request). ARQ systems

are of two types: stop and wait ARQ with half-duplex line and continuous ARQ with full-duplex line. These two types of line control procedures are discussed in this section in relation to their message exchange sequence and transmission efficiencies. Also, the optimum frame size (1/0 buffer size) for continuous ARQ is examined with respect to the speed of the circuit and error rate.

Half-Duplex Transmission. This is the most common line control procedure in use today, and stop and wait ARQ uses this transmission. In stop and wait ARQ, the source waits for an acknowledgment from the receiver before transmitting the next block of data. If the source receives ACK (no error acknowledgment), it continues transmitting the next block; but if NAK (error acknowledgment) is received, the source retransmits the last block of data. Figure 35 illustrates this procedure.

The total transmission time of this type of line control procedure can be computed, as illustrated in Table XII. In this table, 4800 bps (600 bytes/sec) channel is used to transmit 240 bytes of data and it gives a transmission efficiency of $400/934 = 43$ percent. Improvement on transmission efficiency of such a system can be achieved by increasing the size of message blocks. However, this may create other problems, as long messages are more error-prone than short messages.

Full-Duplex Transmission. Another method of improving the throughput is to use full-duplex circuit and to transmit messages without acknowledgments. This method is examined with continuous ARQ, shown in Figure 36. In continuous ARQ, while the blocks are being transmitted the stream of acknowledgments is examined by the transmitting terminal. When the transmitting terminal receives a negative acknowledgment or fails to receive a

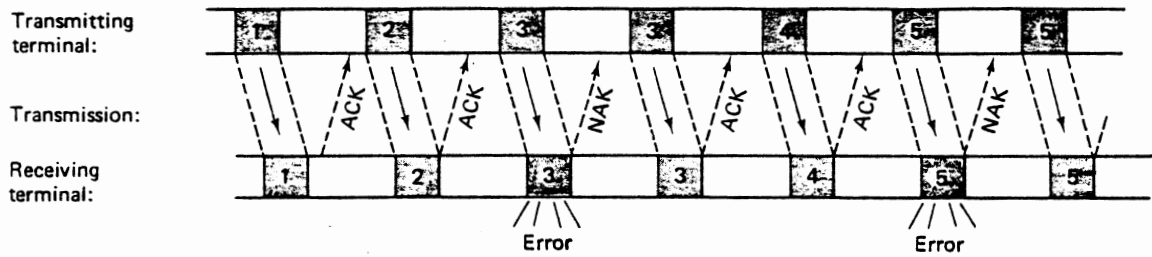
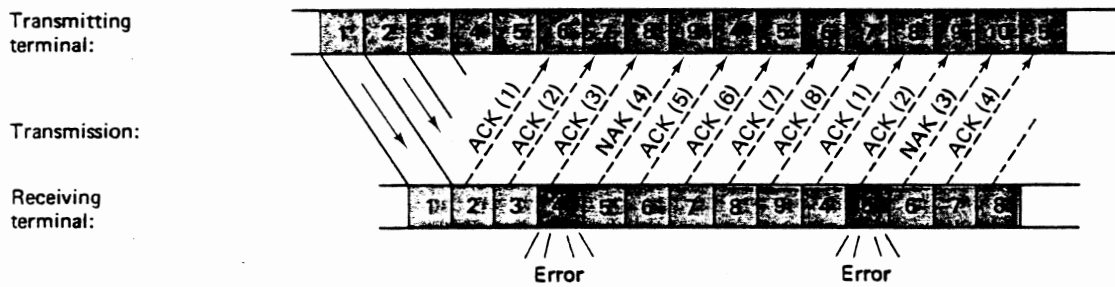
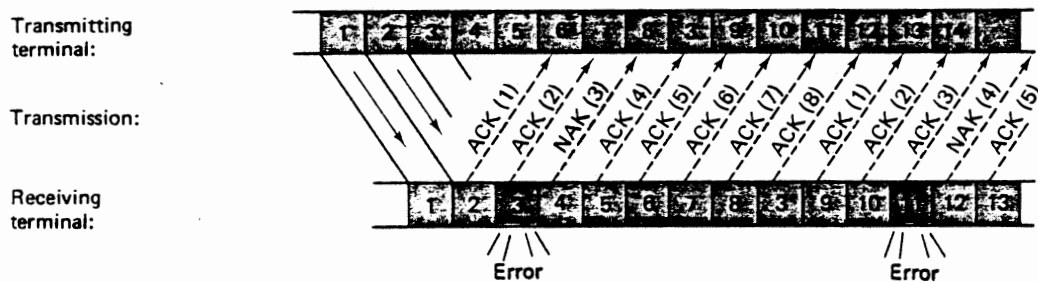


Figure 35. Stop and Wait ARQ (After Martin [35])



(a) Full Duplex with Pull-back



(b) Full Duplex with Selective Repeat

Figure 36. Continuous ARQ (After Martin [35])

TABLE XII
TRANSMISSION TIME LIST

Message transmission time	$240/600 = 400$
Propagation delay	250
Modem delay	10
Receiver reaction time	2
Transmission time for acknowledgement	10
Propagation delay	250
Modem delay	10
Computer reaction time	2
	<hr/>
	Total : 934 msec

After Housley [34].

positive acknowledgment, it must determine which block was incorrect. The blocks are therefore numbered.

Considering the transmission time delay of the satellite, the acknowledgment may be received more than several blocks after it was transmitted. For this reason, usually a 7-bit counter is used for the satellite channels while a 3-bit counter is used for the terrestrial links. With a 7-bit counter, 128 blocks can be numbered before it reinitializes the counter.

The transmission efficiency of the continuous ARQ is computed from

$$e = N_d / (N_d + N_h), \quad (5.19)$$

where N_d is the number of data bits in a frame, and N_h is the number of overhead bits. The transmission efficiencies of various speeds are shown in Figure 37 as a function of a frame size.

The minimum frame size of an efficient transmission is shown in Figure 38 as a function of the speed of the circuit and the block counter. Figure 39 illustrates the optimum frame sizes for various error rates and circuit speed. For example, compressed seismic data transmission time with a channel bandwidth of 9600 bps can be computed from Figure 37. It can be seen that frame size should be greater than 5000 bits in order to approach 100 percent efficiency. From this constraint, the compressed seismic data need to be segmented into frames with a size of 5000 bits. Then transmission time can be estimated by

$$T_t = \frac{N \cdot 5000 \text{ bits}}{9600 \text{ bps}} \quad (5.20)$$

where

$$N = \frac{\text{Compressed seismic data}}{5000} .$$

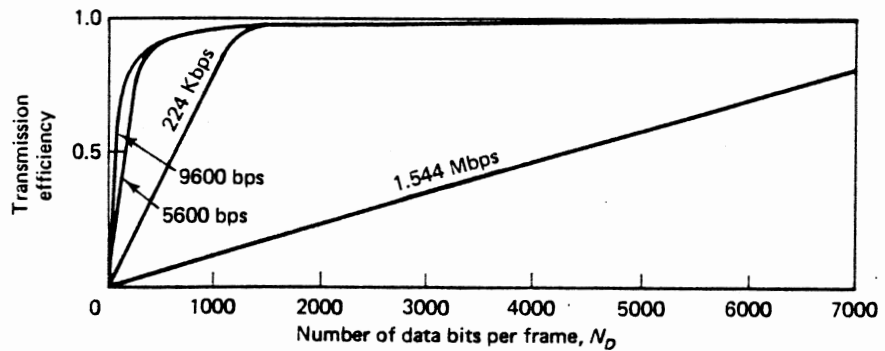


Figure 37. Transmission Efficiency of Satellite Circuits with Common Data Link Controls ($M=127$) (After Martin [35])

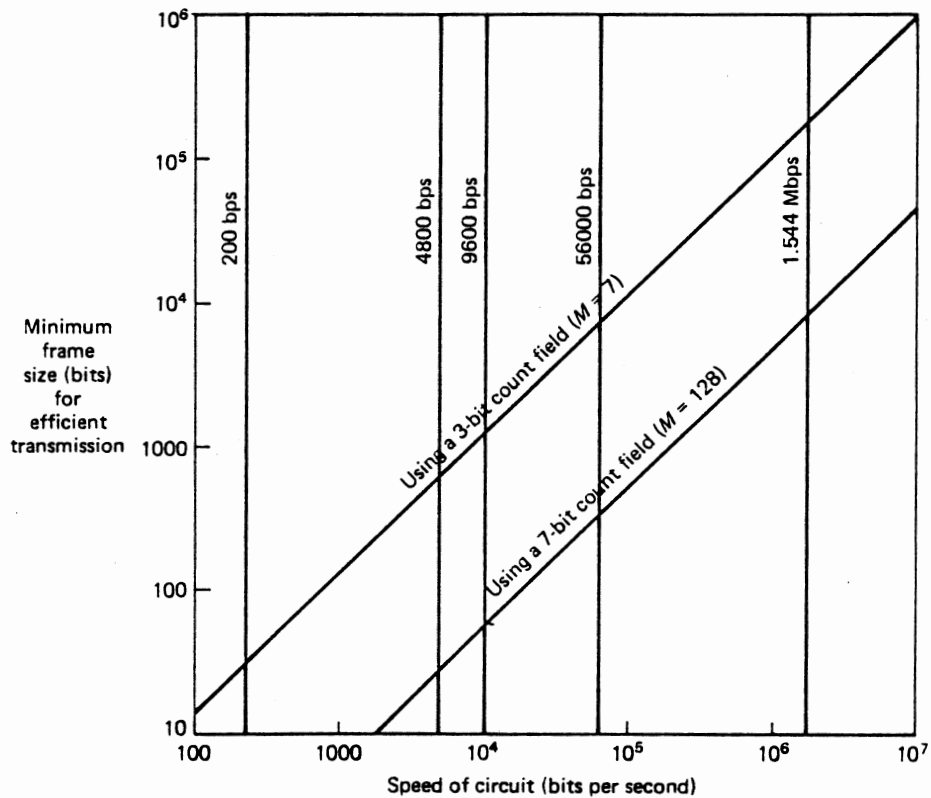


Figure 38. The Minimum Frame Size for Efficient Transmission via Satellite at Different Speeds (After Martin [35])

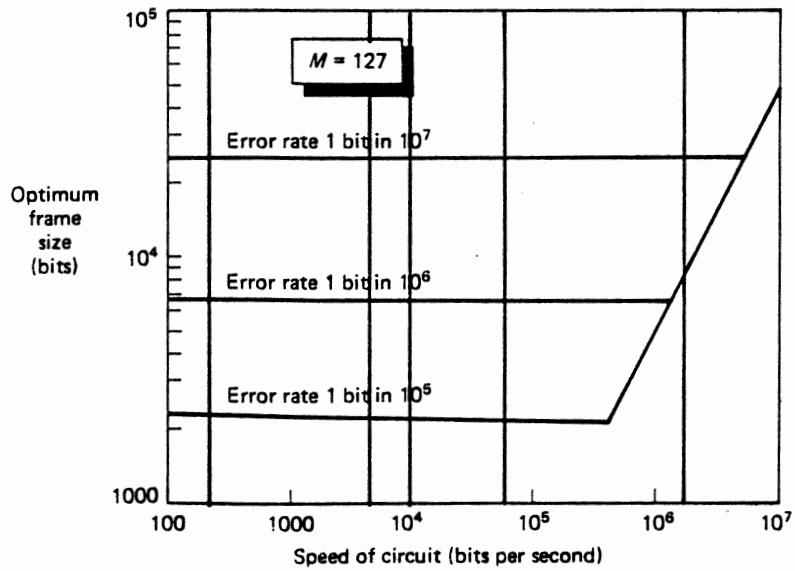


Figure 39. The Optimum Frame Size for Full Duplex Line Control (After Martin [35])

Also, the error rate can be found in Figure 39 such that the chosen frame size can maintain approximately 10^{-6} error rate.

Summary

This chapter has considered general aspects of implementing seismic data compression techniques with respect to execution time, memory size, depth of queue, and transmission time. The execution time and memory size requirements were evaluated in the assembler listing of all subroutines used in the compression techniques. It was shown that time and space requirements are functions of data points, and response time dictates the system structure depth of queue and speed of transmission circuit.

Transmission time is determined by transmission efficiency and speed of circuit when the data rate is fixed. Transmission efficiency was examined with two types of line control procedures: stop and wait ARQ and continuous ARQ. From these examinations, the following aspects were found: stop and wait ARQ cannot achieve high transmission efficiency since it has to wait for the ACK signal response and the response is delayed by the propagation delay of the satellite circuit. High transmission efficiency (above 95%) can be achieved by continuous ARQ, provided that the size of frame is chosen properly. The optimum frame size can be obtained according to the speed of circuit and a block counter. A 7-bit block counter is necessary for a satellite link due to a long propagation delay.

In general, the transmission time of continuous ARQ line control procedure can be obtained from

$$T_t = \frac{N \cdot F}{B \cdot e}, \quad (5.21)$$

where N is the number of frames, F is the frame size, B is the speed of the circuit (or channel bandwidth), and e is the transmission efficiency.

CHAPTER VI

CONCLUSIONS

New avenues into seismic data compression techniques have been developed. Compression techniques have been considered based on orthogonal transforms, digital coding methods, and prediction methods. A hybrid technique, which is a combination of DCT and a μ -law quantization has been developed for vibroseis data and nonimpulsive section of impulsive seismic data. Prediction coding has been developed for the impulsive section of impulsive seismic data, where a predictor is designed by the linear prediction analysis.

A range between five-to-one and eight-to-one compression has been achieved with at least a 30 dB signal-to-noise ratio. This result is obtained empirically, and it is dependent upon sampling rate and signal distribution.

This study differs from previous work [8] in the following aspects. First, earlier work focused on vibroseis data sampled at 500 Hz. This study investigates vibroseis data sampled at 250 Hz and impulsive seismic data sampled at 1000 Hz. Second, in the earlier work, the reconstructed signal at the receiver was measured objectively. In this study, the error in the compression-decompression process is maintained within a 30 dB SNR for every trace. This is defined as reversible in this study. Third, the seismic signal was treated as a two-dimensional image in earlier work. In this study, the seismic signal is treated as a one-dimensional vector.

Finally, earlier work was based on frequency limiting of the Walsh domain. Compression techniques of this study are based on redundancy and entropy reduction. Compression ratios of these two studies need not be compared, since compression is a function of the signal-to-noise ratio.

In this thesis, all sensors are weighted equally and a 30 dB SNR is maintained for each trace. A two-dimensional analysis was not performed due to energy differences among traces. An energy normalization method may be required to pursue two-dimensional data compression for further research.

All compression techniques are simulated in FORTRAN with the IBM 370. For implementing these techniques in mini- or microcomputers, execution time and core size requirements of all subroutines are evaluated with a cross-referenced list in assembler language. The evaluation is based on types and frequency of occurrences of instructions inside "DO" loops, where the iteration of "DO" loops is determined by the number of seismic data points. In order to minimize the core size requirement, a "load and swap" method is recommended, which allows for loading only the executing subroutine into the core memory. Also, general relationships among processing time, depth of queue, transmission time, and response time are discussed. Further research is required for completing a real-time system with respect to system structure and response time.

Data compression techniques have been investigated due to constraints in channel bandwidth, channel capacity, limited storage, and/or transmission cost. However, compression techniques do require computer systems, and thus compression cost needs to be analyzed in order to evaluate the trade-offs between compression cost and transmission cost. Further research in the areas, two-dimensional seismic data compression, real-time

system design, cost analysis, and evaluation of computational optimality of algorithm versus optimality of hardware design should be fruitful. Optimality in terms of computation along with hardware implementation should be considered in parallel.

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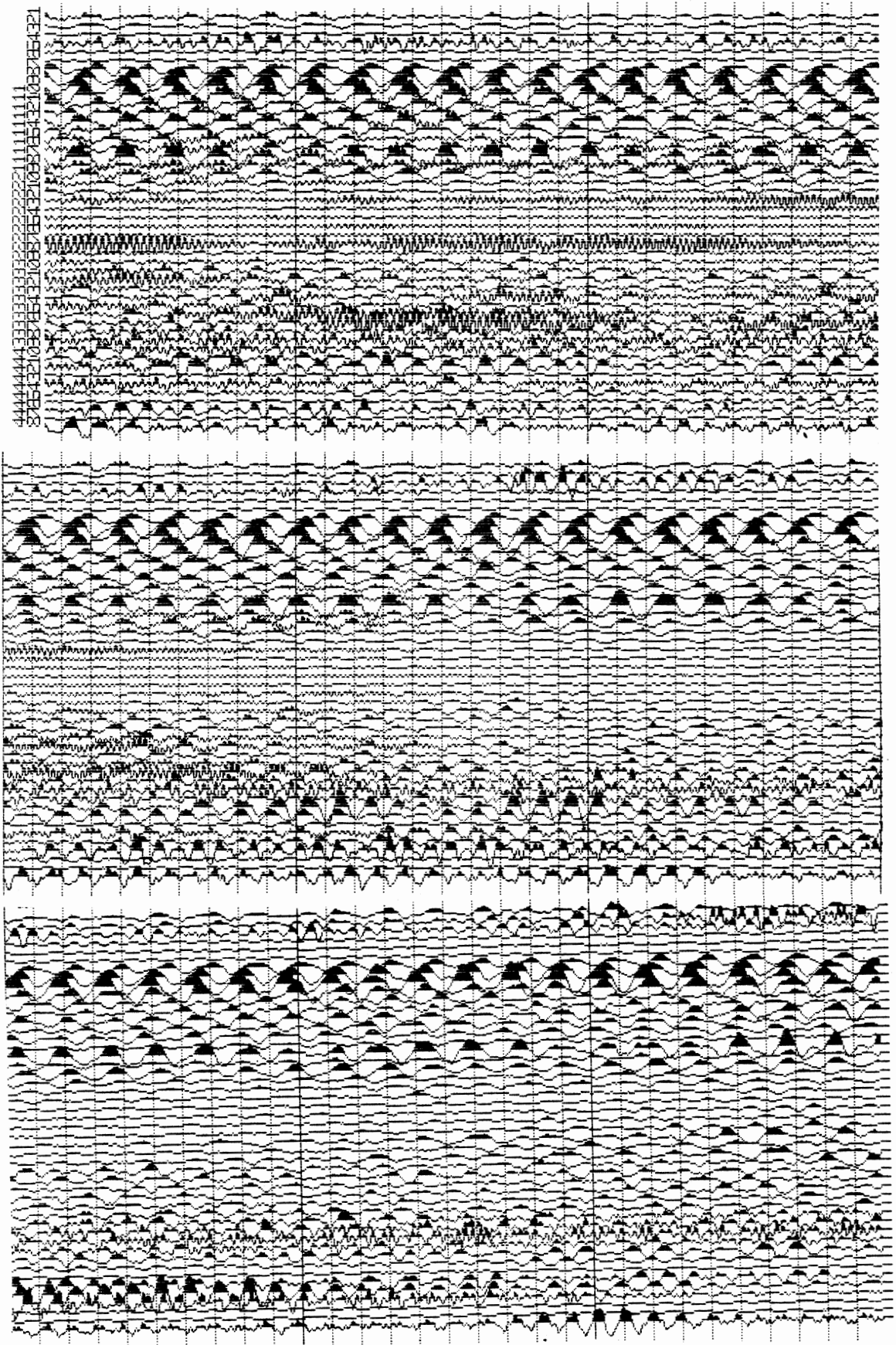
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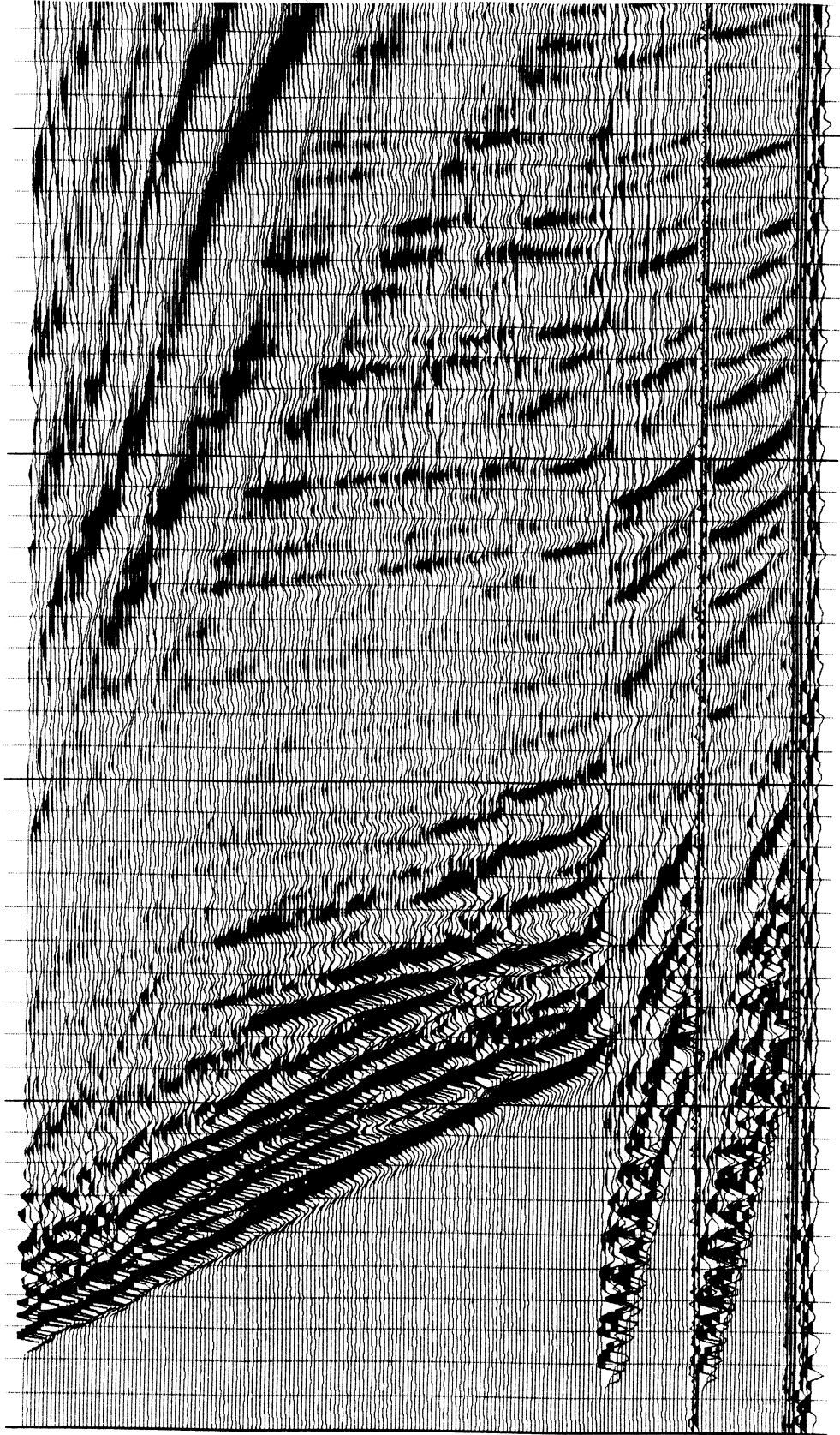
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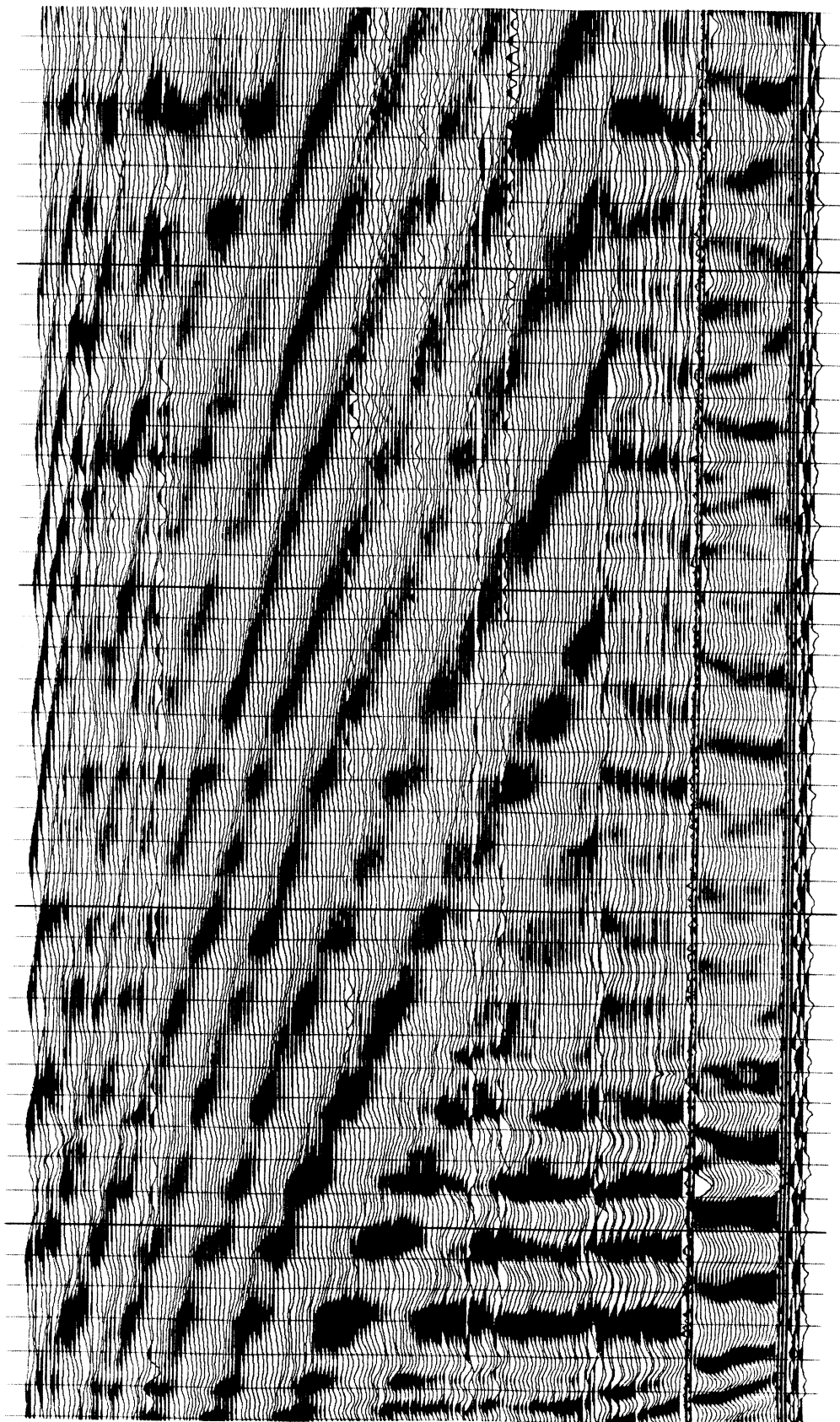
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APPENDIX A

SEISMIC DATA PLOTS







APPENDIX B

CROSS CORRELATION TABLES

CORRELATION COEFFICIENTS / PRGR > |F| UNDER HO: F

	TEST1	TEST2	TEST3	TEST4	TEST5	TEST6	TEST7	TEST8
TEST1	1.00000 0.0000	0.01290 0.3740	0.00928 0.5228	0.00915 0.5284	-0.01023 0.4810	-0.01073 0.4559	-0.00146 0.9198	-0.00865 0.5514
TEST2	0.01290 0.3740	1.00000 0.0000	0.22363 0.0001	0.13457 0.0001	-0.32770 0.0001	-0.13741 0.0001	0.07056 0.0001	-0.18383 0.0001
TEST3	0.00928 0.5228	0.22363 0.0001	1.00000 0.0000	-0.04224 0.0036	0.35505 0.0001	0.07982 0.0001	-0.24053 0.0001	-0.01549 0.2857
TEST4	0.00915 0.5284	0.13457 0.0001	-0.04224 0.0036	1.00000 0.0000	-0.09241 0.0001	0.21682 0.0001	-0.09634 0.0001	0.25413 0.0001
TEST5	-0.01023 0.4810	-0.32770 0.0001	0.35505 0.0001	-0.09241 0.0001	1.00000 0.0000	0.34325 0.0001	0.01698 0.2415	0.16591 0.0001
TEST6	-0.01073 0.4559	-0.13741 0.0001	0.07882 0.0001	0.21682 0.0001	0.34325 0.0001	1.00000 0.0000	-0.13009 0.0001	0.15432 0.0001
TEST7	-0.00146 0.9198	0.07056 0.0001	-0.24053 0.0001	-0.09634 0.0001	0.01698 0.2415	-0.13009 0.0001	1.00000 0.0000	-0.33200 0.0001
TEST8	-0.00865 0.5514	-0.18383 0.0001	-0.01549 0.2857	0.25413 0.0001	0.16591 0.0001	0.15432 0.0001	-0.33200 0.0001	1.00000 0.0000
TEST9	-0.00233 0.3722	0.13273 0.0001	-0.05342 0.0001	-0.30177 0.0001	-0.08934 0.0001	0.04416 0.0023	0.30656 0.0001	-0.54240 0.0001
TEST10	0.00752 0.6045	-0.09432 0.0001	0.01484 0.2065	0.02165 0.1357	0.19623 0.0001	-0.25016 0.0001	0.02629 0.0700	0.35329 0.0001
TEST11	-0.00130 0.9286	0.03054 0.0353	-0.04344 0.0027	0.05689 0.0001	-0.18822 0.0001	0.15242 0.0001	-0.02753 0.0578	-0.03028 0.0369
TEST12	0.01212 0.4033	-0.11076 0.0001	0.01950 0.1790	-0.21376 0.0001	0.17588 0.0001	-0.25899 0.0001	0.35239 0.0001	-0.45169 0.0001
TEST13	-0.01043 0.4703	-0.09080 0.0001	0.03564 0.0140	0.20481 0.0001	0.05570 0.0001	0.25937 0.0001	-0.29625 0.0001	0.23167 0.0001
TEST14	-0.01768 0.2232	-0.02133 0.0303	-0.04155 0.0042	-0.03029 0.0369	-0.05867 0.0001	0.00296 0.8383	0.04704 0.0012	0.23509 0.0001
TEST15	-0.00201 0.8900	-0.05645 0.0001	-0.11508 0.0001	-0.01857 0.2006	0.11404 0.0001	0.05608 0.0001	-0.04545 0.0017	0.05516 0.0001
TEST16	-0.01391 0.3377	-0.03234 0.0236	-0.03896 0.0072	-0.00450 0.7565	-0.00886 0.5417	0.01141 0.4319	-0.00460 0.7513	-0.00180 0.9013

FPO=0 / N = 4750

TEST9	TEST10	TEST11	TEST12	TEST13	TEST14	TEST15	TEST16
-0.00233	0.00752	-0.00130	0.01212	-0.01048	-0.01768	-0.00201	-0.01391
0.8722	0.6045	0.9286	0.4033	0.4703	0.2232	0.8900	0.3377
0.13273	-0.09432	0.03054	-0.11096	-0.09080	-0.03133	-0.05645	-0.03284
0.0001	0.0001	0.0001	0.0001	0.0001	0.0308	0.0001	0.0236
-0.05842	0.01484	-0.04244	0.01950	0.03564	-0.04155	-0.11508	-0.03896
0.0001	0.3065	0.0027	0.1790	0.0140	0.0042	0.0001	0.0072
-0.30177	0.02165	0.05689	-0.21376	0.20481	-0.03029	-0.01857	-0.00450
0.0001	0.1357	0.0001	0.0001	0.0001	0.0309	0.2006	0.7565
-0.08934	0.19623	-0.18622	0.17588	0.05570	-0.05867	0.11404	-0.00886
0.0001	0.0001	0.0001	0.0001	0.0001	0.0001	0.0001	0.5117
0.04416	-0.26016	0.15242	-0.25859	0.25937	0.00290	0.05608	0.01141
0.0023	0.0001	0.0001	0.0001	0.0001	0.8383	0.0001	0.4319
0.30596	0.02629	-0.02753	0.35239	-0.28625	0.04704	-0.04545	-0.00460
0.0001	0.0700	0.0578	0.0001	0.0001	0.0012	0.0017	0.7513
-0.58240	0.35325	-0.03028	-0.45109	0.23167	0.23508	0.09516	-0.00180
0.0001	0.0001	0.0365	0.0001	0.0001	0.0001	0.0001	0.9013
1.00000	-0.59711	0.35055	0.16366	-0.02231	-0.17160	-0.27690	0.00476
0.0000	0.0001	0.0001	0.0001	0.1241	0.0001	0.0001	0.7432
-0.59711	1.00000	-0.54762	0.16913	-0.12580	0.12381	0.28619	-0.01216
0.0001	0.0000	0.0001	0.0001	0.0001	0.0001	0.0001	0.4022
0.35055	-0.54762	1.00000	-0.17766	0.30768	-0.21733	-0.38187	0.01512
0.0001	0.0001	0.0000	0.0001	0.0001	0.0001	0.0001	0.2974
0.16366	0.16513	-0.17766	1.00000	-0.46665	-0.19623	0.29209	0.00695
0.0001	0.0001	0.0001	0.0000	0.0001	0.0001	0.0001	0.6322
-0.02231	-0.12580	0.30768	-0.46665	1.00000	-0.17891	-0.19543	0.00926
0.1241	0.0001	0.0001	0.0001	0.0000	0.0001	0.0001	0.5236
-0.17160	0.12351	-0.21733	-0.15628	-0.17891	1.00000	-0.08172	0.00998
0.0001	0.0001	0.0001	0.0001	0.0001	0.0000	0.0001	0.4918
-0.27690	0.28619	-0.38187	0.25209	-0.19543	-0.08172	1.00000	-0.01446
0.0001	0.0001	0.0001	0.0001	0.0001	0.0001	0.0000	0.3192
0.00476	-0.01216	0.01512	0.00655	0.00926	0.00998	-0.01446	1.00000
0.7432	0.4022	0.2574	0.6322	0.5236	0.4918	0.3192	0.0000

CORRELATION COEFFICIENTS / PROB > |R| UNDER H0:F

	TEST1	TEST2	TEST3	TEST4	TEST5	TEST6	TEST7	TEST8
TEST1	1.00000 0.0000	-0.58097 0.0001	-0.58202 0.0001	-0.58114 0.0001	0.57968 0.0001	0.58829 0.0001	-0.57720 0.0001	-0.58227 0.0001
TEST2	-0.58097 0.0001	1.00000 0.0000	0.99998 0.0001	1.00000 0.0001	-0.55555 0.0001	-0.55871 0.0001	0.55582 0.0001	0.55556 0.0001
TEST3	-0.58202 0.0001	0.99998 0.0001	1.00000 0.0000	0.55959 0.0001	-0.55551 0.0001	-0.55892 0.0001	0.55570 0.0001	1.00000 0.0001
TEST4	-0.58114 0.0001	1.00000 0.0001	0.55995 0.0001	1.00000 0.0000	-0.55554 0.0001	-0.55874 0.0001	0.55981 0.0001	0.55557 0.0001
TEST5	0.57968 0.0001	-0.55555 0.0001	-0.99991 0.0001	-0.55954 0.0001	1.00000 0.0000	0.55834 0.0001	-0.55990 0.0001	-0.55589 0.0001
TEST6	0.58829 0.0001	-0.55871 0.0001	-0.55852 0.0001	-0.55874 0.0001	0.55834 0.0001	1.00000 0.0000	-0.55975 0.0001	-0.55851 0.0001
TEST7	-0.57720 0.0001	0.55582 0.0001	0.55570 0.0001	0.55981 0.0001	-0.55990 0.0001	-0.55775 0.0001	1.00000 0.0000	0.55986 0.0001
TEST8	-0.58227 0.0001	0.55556 0.0001	1.00000 0.0001	0.99957 0.0001	-0.55585 0.0001	-0.55891 0.0001	0.55588 0.0001	1.00000 0.0000
TEST9	-0.58278 0.0001	0.55994 0.0001	0.55999 0.0001	0.55955 0.0001	-0.55584 0.0001	-0.55901 0.0001	0.55558 0.0001	1.00000 0.0001
TEST10	0.57892 0.0001	-0.55588 0.0001	-0.55578 0.0001	-0.55987 0.0001	0.55584 0.0001	0.55827 0.0001	-0.55585 0.0001	-0.99972 0.0001
TEST11	0.55988 0.0001	-0.55151 0.0001	-0.55227 0.0001	-0.55182 0.0001	0.55081 0.0001	0.55540 0.0001	-0.58899 0.0001	-0.99251 0.0001
TEST12	0.58084 0.0001	-1.00000 0.0001	-0.55997 0.0001	-1.00000 0.0001	0.55994 0.0001	0.55870 0.0001	-0.55984 0.0001	-0.55955 0.0001
TEST13	0.58230 0.0001	-0.55557 0.0001	-0.99999 0.0001	-0.99998 0.0001	0.55987 0.0001	0.55903 0.0001	-0.55587 0.0001	-0.55558 0.0001
TEST14	-0.57894 0.0001	0.55993 0.0001	0.55984 0.0001	0.55952 0.0001	-0.55551 0.0001	-0.55827 0.0001	0.55955 0.0001	0.55579 0.0001
TEST15	-0.57909 0.0001	0.55993 0.0001	0.55585 0.0001	0.55952 0.0001	-0.55552 0.0001	-0.55823 0.0001	0.55993 0.0001	0.99581 0.0001
TEST16	0.57872 0.0001	-0.99993 0.0001	-0.55586 0.0001	-0.55952 0.0001	0.55995 0.0001	0.55811 0.0001	-0.55955 0.0001	-0.55583 0.0001

:FO=0 / N = 475C

TEST9	TEST10	TEST11	TEST12	TEST13	TEST14	TEST15	TEST16
-0.58278 0.0001	0.57852 0.0001	0.55688 0.0001	0.58084 0.0001	0.98230 0.0001	-0.97894 0.0001	-0.97909 0.0001	0.57872 0.0001
0.55994 0.0001	-0.59588 0.0001	-0.99151 0.0001	-1.00000 0.0001	-0.99997 0.0001	0.99993 0.0001	0.99993 0.0001	-0.99993 0.0001
0.99999 0.0001	-0.99978 0.0001	-0.99227 0.0001	-0.99997 0.0001	-0.99999 0.0001	0.99984 0.0001	0.99985 0.0001	-0.99986 0.0001
0.99995 0.0001	-0.99987 0.0001	-0.99162 0.0001	-1.00000 0.0001	-0.99998 0.0001	0.99992 0.0001	0.99992 0.0001	-0.99992 0.0001
-0.55984 0.0001	0.55981 0.0001	0.55981 0.0001	0.55984 0.0001	0.99987 0.0001	-0.99991 0.0001	-0.99992 0.0001	0.55995 0.0001
-0.99901 0.0001	0.99927 0.0001	0.99940 0.0001	0.55987 0.0001	0.99903 0.0001	-0.99987 0.0001	-0.99923 0.0001	0.55981 0.0001
0.55958 0.0001	-0.55985 0.0001	-0.55985 0.0001	-0.55984 0.0001	-0.55967 0.0001	0.99995 0.0001	0.99993 0.0001	-0.55995 0.0001
1.00000 0.0001	-0.55972 0.0001	-0.55925 0.0001	-0.55955 0.0001	-0.99998 0.0001	0.99979 0.0001	0.99981 0.0001	-0.55983 0.0001
1.00000 0.0000	-0.55988 0.0001	-0.55984 0.0001	-0.55952 0.0001	-0.99997 0.0001	0.99974 0.0001	0.99976 0.0001	-0.55977 0.0001
-0.99966 0.0001	1.00000 0.0000	0.58975 0.0001	0.55950 0.0001	0.99980 0.0001	-0.99998 0.0001	-0.55995 0.0001	0.55989 0.0001
-0.55984 0.0001	0.98975 0.0001	1.00000 0.0000	0.55138 0.0001	0.99230 0.0001	-0.98996 0.0001	-0.55907 0.0001	0.55913 0.0001
-0.55992 0.0001	0.55950 0.0001	0.55138 0.0001	1.00000 0.0000	0.99997 0.0001	-0.55994 0.0001	-0.55994 0.0001	0.55993 0.0001
-0.55997 0.0001	0.99980 0.0001	0.99230 0.0001	0.55997 0.0001	1.00000 0.0000	-0.55984 0.0001	-0.55985 0.0001	0.99982 0.0001
0.99974 0.0001	-0.99998 0.0001	-0.58996 0.0001	-0.55994 0.0001	-0.99984 0.0001	1.00000 0.0000	0.99999 0.0001	-0.99996 0.0001
0.55976 0.0001	-0.99995 0.0001	-0.55907 0.0001	-0.55994 0.0001	-0.99985 0.0001	0.99999 0.0001	1.00000 0.0000	-0.55995 0.0001
-0.55977 0.0001	0.55989 0.0001	0.55913 0.0001	0.55993 0.0001	0.99982 0.0001	-0.55996 0.0001	-0.99995 0.0001	1.00000 0.0000

APPENDIX C

COMPILED LISTING OF COMPUTER PROGRAMS

REQUESTED OPTIONS: OPT=2,FORMAT,XRFF,LIST,MAP,SIZE(75CK)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LINECOUNT(20) SIZE(0750K) AUTOCOLL(NONE)
SOURCE EBCCIC LIST NCRCK INJECT MAP PERPAT GOSTMT XRFF NCALC NOANSF TERM IBM FLAG(1)

```

15H 0002      SUBROUTINE FDC1(A,NUM)                                00000690
              DCI TRANSFORM USING FFT(IBM VERSION : FFTFC)        00000700
              C          NUM : NUMBER OF DATA POINTS TO BE TRANSFORMED 00000710
              C          A   : INPUT ARRAY OF FFT AND OUTPUT ARRAY OF FFT 00000720
              C          INK : WORK AREA OF FFT                      00000730
              C          WK  : WORK AREA OF FFT                      00000740
              C          X   : OUTPUT ARRAY OF FFT                  00000750
15H 0003      DIMENSION INK(2600),WK(2000),A(1)                    00000760
15H 0004      COMPLEX X(4750)                                       00000770
15H 0005      NUM2=NUM/2                                           00000780
15H 0006      WRWP=3.141592/FLGAT(2*NUM)                            00000790
15H 0007      J=0                                                  00000800
15H 0008      CALL FFTFC(A,NUM,X,INK,WRP)                          00000810
15H 0009      DO 10 I=2,NUM2                                       00000820
15H 0010      10 X(NUM+2-I)=CONJG(X(I))                             00000830
15H 0011      DO 20 I=1,NUM                                         00000840
15H 0012      ARG=WRWP*FLCA(I-1)                                    00000850
15H 0013      X(I)=CMPLX(COS(ARG),SIN(ARG))*X(I)                   00000860
15H 0014      20 A(I)=REAL(X(I))/NUM2                              00000870
15H 0015      A(I)=A(I)/2.0                                         00000880
15H 0016      RETURN                                               00000890
15H 0017      END                                                  00000900
    
```

*****FORTRAN CROSS REFERENCE LISTING*****

SYMBOL	INTERNAL STATEMENT NUMBERS
A	0002 0003 0008 0014 0015 0015
I	0005 0010 0010 0011 0012 0013 0013 0014 0014
J	0007
X	0004 0003 0010 0010 0013 0013 0014
WK	0003 0008
ARG	0012 0013 0013
COS	0013
INK	0003 0008
NUM	0002 0005 0006 0008 0010 0011
SIN	0013
FDC1	0002
NUM2	0005 0007 0014
REAL	0014
WRWP	0006 0012
CMPLX	0013
CONJG	0010
FFTC	0008
FLGAT	0006 0012

*****FORTRAN CROSS REFERENCE LISTING*****

LABEL	DEFINED	REFERENCES
10	0010	0005
20	0014	0011

000000 47 TO F 00C FDC1 FC 15,12(0,1')

000004 07
 000005 CCC4CJE3404C4C
 00000C 90 EC D 00C
 000010 1E AD
 000012 98 CD F 020
 000016 50 40 D 004
 00001A 5C D0 4 CCE
 00001E 07 FC

DC XL1*07*
 DC CL1*FDC1 *
 SIM 14,12,12(13)
 LR 4,13
 LM 12,13,32(1E)
 ST 4,4(0,13)
 ST 12,8(0,4)
 BCR 1E,12

TEMPORARY FOR FIX/FLOAT

00009B 4E000000
 00009C 00C00000

DC XL4*4E000000*
 DC XL4*00000000*

CONSTANTS

0000A0 4E000000
 0000A4 80000000
 0000A8 CCC00000
 0000AC 00000001
 0000B0 00C00002
 0000B4 412C0000
 0000B8 412243F6
 0000BC 00C00000
 0000C0 00000000

DC XL4*4E000000*
 DC XL4*E0000000*
 DC XL4*00000000*
 DC XL4*CCCCC001*
 DC XL4*00000002*
 DC XL4*412C0000*
 DC XL4*412243F6*
 DC XL4*00000000*
 DC XL4*CC000000*

ADDONS FOR VARIABLES AND CONSTANTS

00E6A0 00CC2580
 00E6A4 00005220

DC XL4*00002580*
 DC XL4*00005220*

ADDONS FOR EXTERNAL REFERENCE

00E6A8 00C00000
 00E6B0 00C00000
 00E6B4 00C00000
 00E6B8 00C00000
 00E6BC 00C00000
 00E6F0 5E 50 C 00B
 00E6F4 58 90 C 004
 00E6F8 58 80 C 000
 00E6FC 58 A0 D 0E4
 00E700 58 80 C 02B
 00E704 5E 70 D 0B8
 00E708 58 20 D 0A8
 00E70C 8E 20 0 020
 00E710 10 27
 00E712 50 30 D 010
 00E716 5E 40 D 0A8
 00E71A 5E 40 0 001
 00E71E 18 04
 00E720 57 00 D 07C
 00E724 50 00 D 074
 00E728 68 20 D 070
 00E72C 68 20 D 078
 00E730 78 00 D 090
 00E734 30 C2
 00E736 70 C0 D 014
 00E73A 58 F0 C 013
 00E73E 41 10 D 04C
 00E742 0E EF
 00E744 47 00 0 00B
 00E748 18 4E
 00E74A 58 60 D 0A9
 00E74E 58 F0 C 034

10CC01

DC XL4*00000000*
 DC XL4*CCCC0000*
 DC XL4*00000000*
 DC XL4*00000000*
 DC XL4*00000000*
 L 5, 8(0,12)
 L 5, 4(0,12)
 L 11, 0(0,12)
 L 10, 132(0,13)
 L 5, 40(0,12)
 L 7, 136(0,13)
 L 2, 166(0,13)
 SFDA 2, 32
 DP 2, 7
 ST 3, 176(0,13)
 L 4, 166(0,13)
 SLL 4, 1
 LR 0, 4
 X 0, 124(0,13)
 ST 0, 116(0,13)
 LR 2, 112(0,13)
 SD 2, 120(0,13)
 LE 0, 144(0,13)
 DER 0, 2
 ST 0, 180(0,13)
 L 1E, 24(0,12)
 LA 1, 76(0,13)
 BALP 14,1E
 FC 0, 8(0,0)
 L¹² 1, E
 L 0, 166(0,13)
 L 5, 52(0,12)

A
 CCS
 SIM
 FF1FC
 CMFY#

1
 2
 NUM

NUM2
 NUM

4E000000P0000000C

4E000000P0000000C
 413243F6

WEWF
 FF1FC

NUM
 1E

00E752	18 03	IF	11, 3		
00E754	78 28 9 000	10 IE	2, 0(0, 5)		X
00E758	78 40 9 004	LE	4, 4(5, 5)		X
00E75C	33 44	LCEN	4, 4		
00E75E	18 26	LR	2, 6		
00E760	18 27	SF	2, 7		
00E762	85 20 0 C03	SLL	2, 3		
00E766	70 42 9 014	STE	4, 20(2, 5)		X
00E76A	70 22 5 C10	STE	2, 16(2, 5)		X
00E76E	1A E4	AF	6, 4		
00E770	67 7A C 0E4	PXLE	7, 10, 1E0(12)		10
00E774	5E P0 C 000	10CC02 L	11, 0(0, 12)		
00E778	58 P0 C 028	L	6, 40(0, 12)		P
00E77C	58 70 D 098	L	7, 136(0, 13)		P
00E780	59 00 D 010	L	0, 176(0, 13)		NUM2
00E784	57 00 D 07C	X	0, 124(0, 13)	4E00000000000000	
00E788	50 00 D C74	ST	0, 116(0, 13)		
00E78C	68 C0 D C7C	LD	0, 112(0, 13)		
00E790	6E 00 D 078	SD	0, 120(0, 13)	4E00000000000000	
00E794	70 00 C C24	STE	0, 36(0, 12)		.C02
00E798	1E 28	LR	2, 6		
00E79A	58 00 C 02C	L	0, 44(0, 12)		4
00E79E	50 00 C 030	ST	0, 48(0, 12)		.C02
00E7A2	58 30 C 02C	L	2, 44(0, 12)		4
00E7A6	1E 4E	LR	4, 6		
00E7A9	58 60 C 030	L	0, 48(0, 12)		.C02
00E7AC	18 7A	LR	7, 10		
00E7AE	18 82	LR	6, 2		
00E7B0	58 P0 D 0A8	L	11, 168(0, 13)		NUM
00E7B4	18 27	10000.3 LR	2, 7		
00E7B6	18 2A	SR	2, 10		
00E7B8	19 02	LR	0, 2		
00E7BA	57 00 D 07C	X	0, 124(0, 13)	4E00000000000000	
00E7BE	50 00 D C74	ST	0, 116(0, 13)		
00E7C2	68 20 D 070	LD	0, 112(0, 13)		
00E7C6	68 20 D C78	SD	0, 120(0, 13)	4E00000000000000	
00E7CA	7C 20 D 094	ME	2, 180(0, 13)		MEWF
00E7CE	70 20 D 0A4	STE	2, 164(0, 12)		ARC
00E7D2	5E F0 C C10	L	1E, 16(0, 12)		C05
00E7D6	41 10 D C60	LA	1, 56(0, 12)		
00E7DA	C5 FF	HALR	14, 1E		
00E7DC	47 00 0 C00	BC	0, 13(0, 0)		
00E7E0	70 00 C 039	STE	0, 56(0, 12)		.100
00E7E4	5E F0 C C14	L	1E, 20(0, 12)		SIF
00E7E8	41 10 D 060	LA	1, 56(0, 12)		
00E7EC	05 FF	HALR	14, 1E		
00E7EE	47 00 0 000	BC	0, 13(0, 0)		
00E7F2	70 00 C 03C	STE	0, 60(0, 12)		.101
00E7F6	78 00 C C38	LE	0, 56(0, 12)		.10C
00E7FA	70 00 C C40	STE	0, 64(0, 12)		.102
00E7FE	78 00 C 03C	LE	0, 60(0, 12)		.101
00E802	70 00 C C44	STE	0, 68(0, 12)		.102
00E806	78 08 9 000	LE	0, 6(0, 5)		X
00E80A	70 00 C 038	STE	0, 56(0, 12)		.10C
00E80E	78 08 9 C04	LE	0, 4(0, 9)		X
00E812	70 00 C 03C	STE	0, 60(0, 12)		.101
00E816	5E F0 C C1C	L	1E, 20(0, 12)		CMY4
00E81A	41 10 D C64	LA	1, 100(0, 13)		

00E81E	05 FF	DALR	14.1E	
00E820	47 00 0 000	EC	C, 13(0, 0)	
00E824	70 00 C C48	STE	0, 72(0,12)	.104
00E828	70 28 9 004	SIF	2, 4(8, 5)	X
00E82C	7E 00 C C48	LE	0, 72(0,12)	.104
00E830	70 08 9 000	STE	C, 0(0, 9)	X
00E834	70 00 9 004	20 LF	C, 4(8, 5)	X
00E838	7E 28 9 000	LF	2, 0(0, 9)	X
00E83C	7C 20 C 024	DE	2, 3(0,12)	.102
00E840	70 26 5 000	STE	2, 0(6, 5)	A
00E844	1A C3	AF	E, 3	
00E846	1A E4	AR	E, 4	
00E848	87 7A C 114	FXLE	7.10, 27(112)	100C03
00E84C	59 80 C 000	10CC04 L	11, 0(0,12)	
00E850	58 80 C 028	L	E, 4(0,12)	B
00E854	58 70 D CE0	L	7, 13(0,13)	Z
00E858	7E 20 5 004	LE	2, 4(0, 5)	A
00E85C	7D 20 D C8C	DE	2, 14(0,13)	4120000C
00E860	70 20 5 CC4	STE	2, 4(0, 5)	A
00E864	1B FF	SR	1E.1E	
00E866	5E F0 D 000	L	14, C(0,13)	
00E86A	07 FE	DCR	1E.14	
ADDRESS OF EPILOGUE				
00E86C	58 A0 D C04	L	10, 4(0,13)	
00E870	58 E0 A 00C	L	14, 12(0,10)	
00E874	58 80 A 018	L	11, 24(0,10)	
00E878	58 10 0 CC4	L	1, 4(0,11)	
00E87C	58 20 D 0A0	L	2, 16(0,12)	NUP
00E880	50 20 1 000	ST	2, 0(0, 1)	
00E884	10 EA	LR	1E.1C	
00E886	52 FF A C0C	MVI	12(10),2E5	
00E88A	58 2C A 01C	LM	2.12, 2E(10)	
00E88E	07 FE	BCR	1E.14	
ADDRESS OF PROLOGUE				
00E890	58 C0 D 040	L	12, 72(0,13)	
00E894	58 70 1 004	L	7, 4(0, 1)	
00E898	5E 20 7 C00	L	2, 0(0, 7)	
00E89C	50 20 D 0A8	ST	2, 16(0,12)	NUP
00E8A0	5E 20 1 C00	L	2, 0(0, 1)	
00E8A4	41 30 2 C00	LA	3, 0(0, 2)	
00E8A8	41 50 0 004	LA	E, 4	
00E8AC	1B 25	SR	2, E	
00E8AE	50 20 C 008	ST	2, E(0,12)	
00E8B2	50 30 C C0C	ST	2, 12(0,12)	A
00E8B6	50 30 D 04C	ST	3, 7(0,13)	
00E8BA	47 F0 C 050	BC	1E, BC(0,12)	
ADCON FOR PROLOGUE				
000020	0000E890	DC	XL4*0000E890*	
ADCON FOR SAVE AREA				
000024	C0C00028	DC	XL4*00000028*	
ADCON FOR EPILOGUE				
000028	00C0E86C	DC	XL4*0000E86C*	
ADCON FOR REG 12				
000070	00C0E8A0	DC	XL4*00C0E8A0*	
ADCONS FOR PARAMETER LISTS				
000078	00000000	DC	XL4*C0000000*	NUP
00007C	00C05228	DC	XL4*00005228*	X
000080	00C025E4	DC	XL4*000025E4*	1WK

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00008A	E0C000E0	DC	XL4'B00000E0*	WK
00008B	00C000CC	DC	XL4'ECC0C0CC*	ARC
00008C	00C0E6F0	DC	XL4'0000E6E0*	.102
000090	8000E6D0	DC	XL4'E000E6D0*	.100
TEMPORARIES AND GENERATED CONSTANTS				
00L6C0	00C00000	DC	XL4'00000000*	
00E6C4	00000000	DC	XL4'C00000CC*	
00E6C8	00C00000	DC	XL4'00000000*	
00E6CC	00C00004	DC	XL4'00000004*	
00E6D0	00C00000	DC	XL4'00CC00C0*	
00L6D4	00C00010	DC	XL4'00000010*	
00L6D8	00C00000	DC	XL4'C0C0C0C0*	
00E6DC	00C00000	DC	XL4'C0000000*	
00L6E0	00C00000	DC	XL4'00000000*	
00L6E4	00C00000	DC	XL4'C00000C0*	
00E6E8	00C00000	DC	XL4'00000000*	
00E6EC	0CC40006	DC	XL4'00C40006*	

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINECNT(40) SIZE(0750K) AUTODI(LCNF)

*OPTIONS IN EFFECT*SOURCE EBCCIC LIST NOCHECK (OBJECT MAP FORPAT GCSTMT XREF NCALC NOANSF TERM IUM FLAG(1)

STATISTICS SOURCE STATEMENTS = 16, PROGRAM SIZE = 55502, SUBPROGRAM NAME = FDCI

STATISTICS NO DIAGNOSTICS GENERATED

***** END OF COMPILATION *****

316K BYTES OF CORE NOT USED

REQUESTED OPTIONS: OPT=2,FCFMT,XREF,LIST,MPF,SIZE(750K)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LINECLNT(CC) SIZE(0750K) #LICEL(NONE)
SOURCE EECCLC LIST NCCCK OBJECT MAP FCFMT GUSTM DREF NCALC NOANSF TERM INM FLAG(1)

```

ISN 0002      SUBPDLTINE SELECT(A,E,FMAX),NUM,EFSLCN,FMEU          00001390
ISN 0003      DIMENSION A(1),E(1)                                00001400
ISN 0004      FMAX=0.0                                           00001410
ISN 0005      DO 10 I=1,NUM                                       00001420
ISN 0006      STD=STE(A(I)*A(I))                                  00001430
ISN 0007      10 IF (FMAX.LT. ABS(A(I))) FMAX=ABS(A(I))          00001440
ISN 0009      ICCUNT=0                                           00001450
ISN 0010      STD=STC/FLLCAT(NUM)                                 00001460
ISN 0011      STD=SORT(STD)                                       00001470
ISN 0012      EFSLCN=.05*STD                                      00001480
ISN 0013      DO 15 J=1,NUM                                       00001490
ISN 0014      IF (ABS(A(J)) .LE. EPSLON) GO TO 15                00001500
ISN 0016      ICCUNT=ICCUNT+1                                     00001510
ISN 0017      IF (A(J) .LT. -EFSLCN) E(ICCUNT)=A(J)+EFSLCN     00001520
ISN 0019      IF (A(J) .GT. EPSLON) E(ICCUNT)=A(J)-EFSLCN     00001530
ISN 0021      15 CONTINUE                                       00001540
ISN 0022      RATIC=FMAX/STD                                       00001550
ISN 0023      IF (RATIC .LT. 4.0) FMEL=0.0                       00001560
ISN 0025      IF (RATIC .LT. 15.0) .AND. FATIC .GE. 4.0) FMEU=10.0 00001570
ISN 0027      IF (RATIO .LT. 20.0) .AND. RATIC .GE. 15.0) FMEL=30.0 00001580
ISN 0029      IF (RATIC .LT. 25.0) .AND. RATIO .GE. 20.0) FMEL=50.0 00001590
ISN 0031      IF (RATIC .LT. 35.0) .AND. FATIC .GE. 25.0) FMEL=100.0 00001600
ISN 0033      IF (RATIO .GE. 25.0) FMEU=200.0                   00001610
ISN 0035      FMAX=FMAX-EFSLCN                                     00001620
ISN 0036      NUM=ICCLNT                                         00001630
ISN 0037      RETURN                                             00001640
ISN 0038      END                                               00001650
    
```

```

***** O R T R A N   C R O S S   R E F E R E N C E   L I S T I N G *****
SYMBOL  INTERNAL STATEMENT NUMBERS
A       0002 0003 0006 0007 0007 0007 0014 0017 0017 0015 0015
U       0002 0003 0017 0015
I       0015 0006 0006 0007 0007
J       0013 0014 0017 0017 0015 0019
ABS     0007 0007 0014
NUM     0002 0005 0010 0013 0036
STD     0006 0006 0010 0010 0011 0011 0012 0022
FMAX    0002 0004 0007 0007 0022 0025 0025
FMEU    0002 0023 0025 0027 0029 0031 0033
SORT    0011
FLLCAT 0010
RATIO   0022 0023 0025 0025 0027 0027 0025 0031 0031 0033
RATIO   0025
EPSLON  0002 0012 0014 0017 0017 0019 0019 0019 0022
ICOUNT  0005 0010 0016 0017 0015 0036
SELECT  0002
    
```

```

***** O R T R A N   C R O S S   R E F E R E N C E   L I S T I N G *****
LABEL  DEFINED  REFERENCES
10     0007  0005
15     0021  0013 0014
    
```

000000	47 F0 F 00C	SELF C1	BC	12,12(0,15)	
000004	07		DC	XL1'C7'	
000005	E2C'D3C5C3E340		CC	CL7'SELECT'	
00000C	50 EC D 00C		STM	14,12,12(13)	
000010	10 40		LR	4,12	
000012	58 CD F C20		LM	12,12,22(15)	
000016	50 40 D 004		ST	4,4(6,12)	
00001A	50 CD 4 C08		ST	12,8(0,4)	
00001E	C7 FC		BCR	12,12	
TEMPORARY FLR FIX/FLUAT					
000078	4E000000		DC	XL4'4E000000'	
00007C	C0C00000		DC	XL4'C0C00000'	
CONSTANTS					
000080	4E1C0000		DC	XL4'4E1C0000'	
000084	80C00000		DC	XL4'80C00000'	
0000E8	00C00000		CC	XL4'00C00000'	
0000EC	00C00000		DC	XL4'C0C00000'	
000090	00C00001		DC	XL4'C0C00001'	
000094	3FC00000		DC	XL4'3FC00000'	
000098	41400000		DC	XL4'41400000'	
00009C	41700000		CC	XL4'41700000'	
0000A0	41FC0000		CC	XL4'41FC0000'	
0000A4	42140000		DC	XL4'42140000'	
0000A8	42150000		CC	XL4'42150000'	
0000AC	421E0000		DC	XL4'421E0000'	
0000B0	42230000		DC	XL4'42230000'	
0000B4	42220000		CC	XL4'42220000'	
0000B8	42240000		DC	XL4'42240000'	
0000BC	42280000		CC	XL4'42280000'	
0000C0	00C00000		CC	XL4'00C00000'	
0000C4	00C00000		DC	XL4'00C00000'	
ADCONS FOR VARIABLES AND CONSTANTS					
ADCONS FOR EXTERNAL REFERENCES					
0000F8	00C00000		DC	XL4'00C00000'	A
000100	00C00000		DC	XL4'00C00000'	E
000108	00C00000		DC	XL4'00C00000'	SCRT
000138	58 40 D 000	100C01	L	4, 20E(0,12)	
00013C	58 60 D 008		L	6, 21E(0,12)	
000140	58 E0 D 0C0		L	11, 56E(0,12)	0
000144	58 A0 C 0EC		L	10, 23E(0,12)	4
000148	58 30 D C6E		L	3, 104E(0,12)	J
00014C	78 00 D 064		LE	6, 100E(0,12)	0
000150	70 00 D 080		STE	6, 17E(0,12)	FPA)
000154	58 20 D 048		L	2, 16E(0,12)	FUP
000158	89 20 0 002		SLL	2, 2	
00015C	50 20 D 104		ST	2, 260E(0,12)	CGE
000160	18 EA		LR	6,10	
000162	13 E2		LR	11, 2	
000164	7E C8 A 000	10CC72	LE	6, 01 8, 4)	A
00016E	38 46		LER	4, 6	
00016A	3C 46		NER	4, 6	
00016C	7A 40 D CAC		AF	4, 172E(0,12)	STD
000170	70 40 C 0AC		STE	4, 172E(0,12)	STD
000174	30 26	10	LFER	2, 6	
000176	70 20 D 0E4		STE	2, 22E(0,12)	CGI
00017A	79 20 D 080		CF	2, 17E(0,12)	FPA)

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00017E	47	CO	D	15E		EC	12. 250(0.13)	100004
000182	7C	20	D	CH0	100003	STE	2. 176(0.13)	FMAX
000186	87	8A	D	13C	100004	DXLE	6.10. 216(12)	100002
00018A	58	80	D	C50	100005	L	11. 56(0.13)	
00018E	18	2B				LR	2.11	
000190	50	80	D	0C4		ST	11. 196(0.13)	ICCLNT
000194	58	00	D	CAB		L	0. 166(0.13)	NUM
000198	57	00	D	05C		X	0. 52(0.13)	4E00000000000000
00019C	50	00	D	C5A		ST	0. 84(0.13)	
0001A0	68	20	D	050		LD	2. 80(0.13)	
0001A4	6E	20	D	C58		SD	2. 86(0.13)	4E00000000000000
0001A8	78	00	D	CAC		LE	0. 172(0.13)	STD
0001AC	20	02				DER	0. 2	
0001AE	70	00	D	0AC		STE	0. 172(0.13)	STC
0001B2	58	F0	D	0E0		L	15. 224(0.13)	SORT
0001B6	41	10	D	04C		LA	1. 70(0.13)	
0001BA	05	EF				FALR	14.15	
0001BC	47	00	D	00B		BC	0. 11(0. 01)	
0001C0	70	00	D	10B		STE	0. 264(0.13)	.10C
0001C4	70	00	D	CAC		STE	0. 172(0.13)	STD
0001C8	78	20	D	0CC		LE	2. 106(0.13)	3FCCCCCE
0001CC	7C	20	D	10B		ME	2. 264(0.13)	.100
0001D0	70	20	D	0C0		STE	2. 192(0.13)	E1SLCA
0001D4	33	22				LCER	2. 2	
0001C6	70	20	D	0E8		STP	2. 232(0.13)	.C02
0001DA	1E	E2				LP	2. 2	
0001CC	1E	7B				LR	1.11	
0001DE	1E	EB				LP	6.11	
0001E0	18	5A				LR	5.10	
0001E2	5E	80	D	104		L	11. 260(0.13)	.C0E
0001E6	78	69	A	000	100006	LE	0. 0(9. 4)	A
0001EA	30	46				LPER	4. 6	
0001EC	75	40	D	0C0		CF	4. 152(0.13)	EPSLCA
0001F0	47	C0	D	1F6		BC	12. 502(0.13)	1E
0001F4	1A	53			100007	AF	2. 2	
0001F6	1A	7A				AF	7.10	
0001F8	1A	8A				AR	8.10	
0001FA	79	60	D	0E8		CE	0. 232(0.13)	.C02
0001FE	47	A0	D	1EA		FC	10. 494(0.13)	100005
000202	39	26			100008	LER	2. 6	
000204	7A	20	D	0C0		AF	2. 152(0.13)	EPSLCA
000208	70	28	E	000		STE	2. 0(8. 6)	E
00020C	79	60	D	0C0	100009	CF	6. 152(0.13)	EPSLCA
000210	47	C0	D	1F6		FC	12. 502(0.13)	1E
000214	38	26			100010	LER	2. 6	
000216	78	20	D	0C0		SE	2. 152(0.13)	EPSLCA
00021A	70	27	E	000		STE	2. 0(7. 6)	P
00021E	87	5A	D	18E	15	DXLE	5.10. 446(12)	100006
000222	50	50	D	0C4		ST	2. 152(0.13)	ICCLNT
000226	58	80	D	0C0	100011	L	11. 56(0.13)	0
00022A	78	20	D	080		LF	2. 176(0.13)	FMAX
00022E	70	20	D	0AC		CE	2. 172(0.13)	STD
000232	70	20	D	00B		STF	2. 184(0.13)	FATIC
000236	79	20	D	C70		CE	2. 112(0.13)	41A0000C
00023A	47	A0	D	21E		HC	10. 542(0.13)	100012
00023E	78	00	D	0C4	100012	LF	0. 100(0.13)	0
000242	70	00	D	084		STF	0. 180(0.13)	EMUL
000246	78	00	D	080	100012	LF	0. 104(0.13)	FATIC

00024A	79 00 D 07E		CE	0, 120(0.13)	41F00000
00024E	47 A0 C 23E		FC	10, 274(0.13)	10CC1E
000252	78 00 D 089	20CC01	LF	0, 184(0.13)	FATIC
000256	79 00 D 070		CE	0, 112(0.13)	41A00000
00025A	47 50 C 23E		EC	0, 274(0.13)	10CC1E
00025E	78 00 D 07A	10CC14	LE	0, 112(0.13)	41A00000
000262	70 00 D 084		STE	0, 180(0.13)	FMEU
000266	78 00 D 088	100013	LE	0, 184(0.13)	FATIC
00026A	79 00 D 07C		CE	0, 124(0.13)	42140000
00026E	47 A0 D 25E		BC	10, 200(0.13)	10CC17
000272	78 00 D 08E	20CC02	LE	0, 184(0.13)	FATIC
000276	79 00 D 07E		CE	0, 120(0.13)	41F00000
00027A	47 50 D 25E		PC	0, 200(0.13)	10CC17
00027E	78 00 D 0E4	10CC16	LE	0, 132(0.13)	421E0000
000282	70 00 D 0E4		STE	0, 180(0.13)	FMEU
000286	78 00 D 088	10CC17	LE	0, 184(0.13)	FATIC
00028A	79 00 D 0E0		CE	0, 128(0.13)	42190000
00028E	47 A0 D 27E		EC	10, 238(0.13)	10CC15
000292	78 00 C 08C	20CC03	LE	0, 188(0.13)	FATIC
000296	79 00 D 07C		CE	0, 124(0.13)	42140000
00029A	47 50 D 27E		EC	0, 238(0.13)	10CC15
00029E	78 00 D 0E0	10001E	LE	0, 140(0.13)	42320000
0002A2	70 00 D 084		STE	0, 180(0.13)	FMEU
0002A6	78 00 D 088	100019	LE	0, 184(0.13)	FATIC
0002AA	79 00 D 0E0		CE	0, 136(0.13)	42230000
0002AE	47 A0 D 29E		PC	10, 270(0.13)	100021
0002B2	78 00 D 08E	20CC01	LE	0, 184(0.13)	FATIC
0002B6	79 00 D 0E0		CE	0, 128(0.13)	42190000
0002BA	47 50 D 25E		PC	0, 270(0.13)	10CC21
0002BF	78 00 D 090	10CC20	LE	0, 144(0.13)	42640000
0002C2	70 00 D 084		STE	0, 180(0.13)	FMEU
0002C6	78 00 D 088	100021	LE	0, 184(0.13)	FATIC
0002CA	79 00 D 088		CE	0, 136(0.13)	42230000
0002CE	47 50 D 282		EC	0, 250(0.13)	10CC22
0002D2	78 00 D 094	10CC22	LF	0, 148(0.13)	42CECCCC
0002D6	70 00 D 084		STE	0, 180(0.13)	FMEU
0002DA	78 00 D 0E0	100023	LE	0, 176(0.13)	FMAX
0002DE	78 00 D 0C0		SE	0, 192(0.13)	FPSLGN
0002E2	70 00 C 080		STE	0, 176(0.13)	FMAX
0002E6	58 00 D 0CA		L	0, 196(0.13)	ICCLAT
0002EA	50 00 D 0AE		ST	0, 166(0.13)	NUM
0002EE	18 FF		SF	10, 10	
0002F0	58 00 D 0CC		L	14, 0(0.13)	
0002F4	07 FE		DCR	10, 14	
*ADDRESS OF EPILOGUE					
0002F6	58 A0 D 004		L	10, 4(0.13)	
0002FA	58 E0 A 00C		L	14, 12(0.10)	
0002FE	58 B0 A 018		L	11, 24(0.10)	
000302	5E 10 B 0CE		L	1, 8(0.11)	
000306	78 20 D 080		LF	0, 176(0.13)	FMAX
00030A	70 20 1 C00		STE	0, 0(0. 1)	
00030E	58 20 B 00C		L	0, 12(0.11)	
000312	58 30 D 0A8		L	0, 166(0.13)	NUM
000316	50 30 2 C00		ST	0, 0(0. 2)	
00031A	58 40 E 010		L	4, 16(0.11)	
00031E	78 40 D 0C0		LF	4, 192(0.13)	FPSLGN
000322	70 40 4 C00		STE	0, 0(0. 4)	
000326	5E 50 E 014		L	0, 20(0.11)	

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00032A	78 60 D 004	LE	0, 180(0,13)	FREU
00032E	70 60 S 000	STE	0, 0(0, 0)	
000332	18 DA	LR	13,1C	
000334	52 FF A 00C	MVI	12(10),255	
000338	5E 2C A 01C	LM	2,12, 28(1C)	
00033C	07 FE	RCR	12,14	
ADDRESS OF PROLOGUE				
00033E	5E 7A 1 008	LM	7,10, 9(1)	
000342	78 20 7 000	LE	2, 0(0, 7)	
000346	70 20 D 080	STE	2, 176(0,13)	FRA>
00034A	5E 2C 8 000	L	2, 0(0, 8)	
00034E	50 20 D 0A8	ST	2, 16E(0,13)	FRA>
000352	78 20 5 000	LE	2, 0(0, 5)	
000356	70 20 D 0C0	STE	2, 152(0,13)	EFSLEN
00035A	78 20 A 000	LE	2, 0(0,10)	
00035E	70 20 D 084	STE	2, 180(0,13)	FREL
000362	5E 20 1 000	L	2, 0(0, 1)	
000366	41 30 2 000	LA	2, 0(0, 2)	
00036A	41 50 0 0C4	LA	5, 4	
00036E	1E 25	SR	2, 5	
000370	50 20 D 0D0	ST	2, 20E(0,13)	
000374	5C 30 D 0D4	ST	2, 212(0,13)	A
000378	58 20 1 004	L	2, 4(0, 1)	
00037C	41 30 2 000	LA	2, 0(0, 2)	
000380	41 50 0 004	LA	5, 4	
000384	1E 25	SR	2, 5	
000386	50 20 D 0D8	ST	2, 216(0,13)	
00038A	50 30 D 0DC	ST	2, 220(0,13)	E
00038E	47 F0 D 110	FC	1E, 272(0,13)	
ADCLN FOR PROLOGUE				
000020	0000033E	CC	XL4*0000033E*	
ADCLN FOR SAVL AREA				
000024	00000028	DC	XL4*00000028*	
ADCLN FOR EPILOGUE				
000028	000002FE	DC	XL4*000002FE*	
ADCLNS FOR PARAMETER LISTS				
000074	00000004	CC	XL4*00000004*	STC
TEMPORARIES AND GENERATED CONSTANTS				
00010C	00000000	CC	XL4*00000000*	
000110	00000000	CC	XL4*00000000*	
000114	00000004	CC	XL4*00000004*	
000118	00000000	DC	XL4*00000000*	
00011C	00000000	CC	XL4*00000000*	
000120	00000000	DC	XL4*00000000*	
000124	00000000	CC	XL4*00000000*	
000128	00000000	DC	XL4*00000000*	
00012C	00000000	CC	XL4*00000000*	
000130	00000000	CC	XL4*00000000*	
000134	00000000	DC	XL4*00000000*	

*OPTIONS IN EFFECT*NAME(PAIN) OPTIMIZE(2) LINECOUNT(60) SIZE(0750K) ALTCBL(NINF)

*OPTIONS IN EFFECT*SOURCE ERCCIC LIST NDCCK OBJECT WAF FCFMAY GLESTX *REF NCALC PDANSF TRM DIM FLAG(1)

STATISTICS SOURCE STATEMENTS = 37, PROGRAM SIZE = 514, SCRIPTEGFAF NAME =SELECT

REQUESTED OPTIONS: OPT=2,FCRPAI,XREF,LIST,MAP,SIZE(750K)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LINECOUNT(60) SIZE(0750K) AUTOBLINDNE
SOURCE ERCCIC LIST KCODEK OBJECT MAP FCRPAI GGSTMT XREF NOALC NUANSF TERN IBM FLAG(1)

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15H 0002      SUBROUTINE CCMENU(RDATA,NUM,EPSLON,XMAX,FMEL)      00000910
              MEU-LAW CONVERTER                                00000920
              C                                               00000930
              C                                               00000940
              C      RDATA : INPUT ARRAY OF COEFFICIENTS OF SEISMIC DATA TRACE 00000950
              C      RMEL  : MEU-VALUE SET AS 500              00000960
              C      XMAX  : ABSOLUTE MAXIMUM VALUE OF RDATA ARRAY 00000970
              C      NUM   : NO. OF DATA POINTS PER SWEEPS     00000980
              C      EPSLON : THRESHOLD VALUE OF INSIGNIFICANT DCI COEFFICIENTS 00000990
              C                                               00010000
15H 0003      DIMENSION RDATA(1)                                00010010
              C      APPLY MEU-LAW                              00010020
15H 0004      RMEU=FMEL+1.0                                     00010030
15H 0005      DEN=ALG(RMEL)                                    00010040
15H 0006      EPSLON=EPSLON/XMAX                               00010050
15H 0007      EPSLON=1.0+FMEL*EPSLON                           00010060
15H 0008      EPSLON=XMAX*ALG(EPSLON)                          00010070
15H 0009      EPSLON=(EPSLON/DEN)                              00010080
15H 0010      DO 10 I=1,NUM                                     00010090
15H 0011      RDATA(I)=0.0                                     00010100
15H 0012      SIGN=1.0                                         00010110
15H 0013      ABSR=ABS(RDATA(I))                               00010120
15H 0014      IF (RDATA(I).LT.0.0) SIGN=-1.0                  00010130
15H 0015      ABSR=ABSR/XMAX                                    00010140
15H 0016      ABSR=1.0+RMEL*ABSR                               00010150
15H 0017      ABSR=XMAX*ALG(ABSR)                              00010160
15H 0018      ABSR=(ABSR/DEN)*SIGN                             00010170
15H 0019      RDATA(I)=ABSR                                    00010180
15H 0020      10 CONTINUE                                      00010190
15H 0021      20 RETLN                                         00010200
15H 0022      END                                             00012000
    
```

***** FORTRAN CROSS REFERENCE LISTING*****

SYMBOL	INTERNAL	STATEMENT	NUMBERS	CROSS	REFERENCE	LISTING
I	0010	0011	0013	0014	0020	
ABS	0013					
DEN	0005	0007	0015			
NUM	0002	0010				
ABSR	0012	0016	0017	0017	0018	0015
ALOG	0005	0008	0018			
RMEU	0002	0004	0007	0017		
SIGN	0012	0014	0015			
XMAX	0002	0006	0008	0016	0018	
RDATA	0002	0003	0011	0013	0014	0020
FMEL	0004	0005				
CONMLU	0002					
EPSLON	0002	0006	0006	0007	0008	0005

***** FORTRAN CROSS REFERENCE LISTING*****

LABEL	DEFINED	REFERENCES
10	0021	0010
20	0022	

000174	7E 29 B 000	LE	2, 0(5, 8)	FLATA
000178	30 02	LPER	0, 2	
00017A	70 00 D C7E	STE	0, 120(0,13)	ABSF
00017E	32 22	LIER	2, 2	
000180	41 A0 0 1A4	FC	10, 25(0,12)	10CCCC
000184	70 00 D 000	LE	0, 17(0,13)	.C02
0001E8	70 00 D 080	STE	0, 128(0,12)	SIGN
00018C	78 20 D C78	LE	2, 120(0,13)	ABSF
000190	70 20 D 0E4	DE	2, 122(0,13)	XMAX
000194	7C 20 D C7C	ME	2, 124(0,13)	PAEL
000198	7A 20 D C60	AE	2, 9(0,13)	41100000
00019C	70 20 D 078	STE	2, 120(0,13)	AFSF
0001A0	5E F0 D 0A0	L	15, 160(0,13)	ALOG
0001A4	41 10 D 054	LA	1, 84(0,13)	
0001A8	05 EF	EALR	14,15	
0001AA	47 00 0 012	FC	0, 18(0, 0)	
0001AE	78 20 D C84	LE	2, 122(0,12)	XMAX
0001B2	3C 20	MER	2, 0	
0001B4	70 20 D C70	DE	2, 112(0,13)	DEF
0001B8	7C 20 D 080	ME	2, 128(0,12)	SIGN
0001BC	70 29 B 000	STE	2, 0(5, 8)	EDATA
0001C0	87 9A 0 13C	10 EXLE	5,10, 31(12)	100002
0001C4	10 FF	20 SP	15,16	
0001C6	58 E0 D 000	L	14, 0(0,12)	
0001CA	07 FE	ECR	15,14	
ADDRESS OF EPILOGUE				
0001CC	5E A0 D C04	L	10, 4(0,13)	
0001D0	58 E0 A 00C	L	14, 12(0,10)	
0001D4	5E B0 A 018	L	11, 24(0,10)	
0001D8	5E 10 0 C0E	L	1, 8(0,11)	
0001DC	78 20 D 090	LE	2, 144(0,13)	EFSLEN
0001E0	70 20 1 000	STE	2, 0(0, 1)	
0001E4	18 0A	LR	13,10	
0001E6	92 FF A 00C	MVI	12(10),255	
0001EA	5E 2C A C1C	LM	2,12, 28(10)	
0001EE	07 FE	RCR	15,14	
ADDRESS OF PROLOGUE				
0001F0	58 7A 1 C04	LM	7,10, 4(1)	
0001F4	58 20 7 000	L	2, 0(0, 7)	
0001F8	50 20 D C74	ST	2, 116(0,13)	NUM
0001FC	78 2C 8 000	LE	2, 0(0, 8)	
000200	70 20 D 070	STE	2, 144(0,13)	EFSLEN
000204	7E 20 9 C00	LE	2, 0(0, 9)	
000208	70 20 D 0E4	STE	2, 132(0,13)	XMAX
00020C	78 20 A 000	LE	2, 0(0,10)	
000210	7C 20 D C7C	STE	2, 124(0,13)	FMEL
000214	58 20 1 000	L	2, 0(0, 1)	
000218	41 30 2 000	LA	2, 0(0, 2)	
00021C	41 50 0 C04	LA	5, 4	
000220	1E 25	SR	2, 2	
000222	50 20 D CA4	ST	2, 164(0,12)	
000226	50 30 D 0A8	ST	2, 176(0,13)	FLATA
00022A	47 F0 D 0C4	FC	15, 196(0,13)	
ADCLN FOR PROLOGUE				
000020	000001F0	DC	XLA*000001F0'	
ADCLN FOR SAVE AREA				
000024	00000020	DC	XI*00000020'	
ADCLN FOR EPILOGUE				

ADCON, FOR PARAMETER LISTS	00002B 00C001C	DC	X14'000001CC'	
	000074 00C00000	DC	X14'00000000'	FMEU
	000078 00C0000B	DC	X14'00C0000B'	E'5LEN
TEMPORARIES AND GENERATED CONSTANTS	00007C 00C000A0	DC	X14'00C000A0'	ABSF
	0000D4 00C00000	DC	X14'0000C000'	
	0000D8 00C00000	DC	X14'00CC0C00'	
	0000DC 00C00004	CC	X14'0000C004'	
	0000E0 00C00000	DC	X14'00CC0000'	
	0000E4 00000000	CC	X14'00000000'	
	0000E8 00C00000	DC	X14'00000000'	

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINECOUNT(20) SIZE(0750K) ALICDRL(ACNE)

*OPTIONS IN EFFECT*SOURCE EBCCIC LIST ACDECK (OBJECT MAP FC*AT CC*PT) *FEF NCALC NOANSF TERM IBM FLAG(1)

STATISTICS SOURCE STATEMENTS = 22, PROGRAM SIZE = 556, SUBPROGRAM NAME =CCNMEU

STATISTICS NO DIAGNOSTICS GENERATED

***** END OF COMPILATION *****

316K BYTES OF CORE NOT USED

REQUESTED OPTIONS: OPT=2,FORMAT,XREF,LIST,MAP,SIZE(750K)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LINKCOUNT(0) SIZE(0750K) AUTODBL(NONE) SOURCE FBCCIC LIST NOCHECK OBJECT MAP FORMAT GCSTMT XREF NOALC NOANSF TERM IBM FLAG(1)

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1SN 0002      SUBROUTINE CCDE(RDATA,EPSLCN,XMAX,ICCODE,NLM)          00001210
C                                                     00001220
C      RDATA(I)      :      INPUT ARRAY          00001230
C      EPSLCN       :      EFFECTIVE LEVEL      00001240
C      XMAX         :      MAXIMUM INFLU VALUE   00001250
C      ICCODE       :      MULTIPLY ARRAY      00001260
C      NLM         :      NUMBER OF DATA IN THE INPLT ARRAY 00001270
C      IULVL       :      UPPER LEVEL LIMIT     00001280
C      ILLVL       :      LOWER LEVEL LIMIT     00001290
C                                                     00001300
1SN 0003      DIMENSION RDATA(I), ICCODE(I)          00001310
1SN 0004      DELTA=XMAX/14.0                          00001320
1SN 0005      IULVL=03                                  00001330
1SN 0006      ILLVL=-64                                00001340
C                                                     00001350
1SN 0007      I=0                                       00001360
1SN 0008      DO 10 L=1,NUM                             00001370
1SN 0009      I=111                                     00001380
1SN 0010      IF(RDATA(J).LT.0.0) GO TO 15              00001390
1SN 0011      CODE=RDATA(I)/DELTA                      00001400
1SN 0012      ICCODE(I)=INT(CODE)                      00001410
1SN 0013      GO TO 20                                  00001420
1SN 0014      15 CODE=RDATA(I)/DELTA-1.0               00001430
1SN 0015      ICCODE(I)=INT(CODE)                      00001440
1SN 0016      20 IF(ICCDE(I).GT. IULVL) ICCODE(I)=IULVL 00001450
1SN 0017      IF(ICCODE(I).LT. ILLVL) ICCODE(I)=ILLVL 00001460
1SN 0018      10 CONTINUE                              00001470
1SN 0019      RETURN                                   00001480
1SN 0020      ENR                                       00001490

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*****ORTRAN CROSS REFERENCE LISTING*****

SYMBOL	INTERNAL STATEMENT NUMBERS
I	0007 0009 0009 0012 0012 0015 0016 0017 0017 0015 0015
J	0008 0010
INT	0013 0016
NUM	0002 0008
CODE	0012 0013 0015 0016
XMAX	0002 0004
CCDE	0002
DELTA	0004 0012 0015
ICCODE	0002 0003 0013 0016 0017 0017 0015 0015
ILLVL	0006 0019 0019
IULVL	0005 0017 0017
RDATA	0002 0003 0010 0012 0015
EPSLCN	0002

*****ORTRAN CROSS REFERENCE LISTING*****

LABEL	DEFINED	REFERENCES
10	0021	0008
15	0015	0010
20	0017	0014

000000	47 F0 F C0C	CODEF	EC	15,12(10,15)	
000004	C7		DC	XL1*07*	
000005	C3E6C4C5C94040		CC	CL7*CODER *	
00000C	90 EC D 00C		STM	14,12,12(13)	
000010	1B 4D		LR	4,13	
000012	9B CD F 020		LM	12,12,32(15)	
000016	50 40 D C04		ST	4,4(10,13)	
00001A	50 00 4 00B		ST	12,E(10,4)	
00001E	C7 FC		ECR	15,12	
TEMPERARY FOR FIX/FLCAT					
00007B	00000000		CC	XL4*C0C00000*	
00007C	C0C00000		DC	XL4*00000000*	
CONSTANTS					
000080	4FC00000		DC	XL4*4F0R0000*	
000084	00C00000		DC	XL4*C0C00000*	
000088	00C00000		DC	XL4*00C0C0C0*	
0000EC	00C00000		DC	XL4*00000000*	
000090	00000001		DC	XL4*C0C0C001*	
000094	00C0003F		DC	XL4*0000003F*	
000098	00C00040		DC	XL4*0000C040*	
00009C	41100000		DC	XL4*411000C0*	
0000A0	42400000		DC	XL4*42400000*	
0000A4	00C00000		DC	XL4*00000000*	
0000A8	00C00000		DC	XL4*00000000*	
ADCONS FOR VARIABLES AND CONSTANTS					
ADCONS FOR EXTERNAL REFERENCES					
0000E0	00000000		DC	XL4*C0C0C0C0*	ICODE
0000E8	00C00000		DC	XL4*00000000*	IFDATA
000108	58 E0 D 0C0	100001	L	C, 152(0,13)	
00010C	58 90 D 0B8		L	C, 184(0,13)	
000110	78 E0 D C94		LF	C, 14E(0,13)	XMA>
000114	58 E0 D 060		L	11, 9C(0,13)	0
000118	5F A0 D 0CC		L	10, 204(0,13)	4
00011C	3E 46		LEP	4, 6	
00011E	7D 40 D 078		DE	4, 12C(0,13)	42400000
000122	5E 00 D C0C		L	0, 10E(0,13)	E3
000126	50 00 D 0A4		ST	C, 164(0,13)	IUNVL
00012A	58 00 D 070		L	C, 112(0,13)	64
00012E	13 00		ECR	C, 0	
000130	50 00 C 0A0		ST	C, 16C(0,13)	ILLVL
000134	E8 20 D C8C		L	2, 140(0,13)	NUM
000138	E9 20 0 C02		SLL	2, 2	
00013C	18 4B		LR	4,11	
00013E	1E 5D		LR	5,11	
000140	1E 7A		LP	7,1C	
000142	18 8B		LF	E,11	
000144	1E 02		LF	11, 2	
000146	1A 8A	100002	AR	E,1C	
000148	1A 4A		AR	4,10	
00014A	1A 5A		AP	5,10	
00014C	78 07 6 000		LE	C, C(7, C)	IFDATA
000150	32 C0		LICR	0, 0	
000152	47 50 D 14C		DC	C, 332(0,13)	IF
000156	78 65 6 000	100003	LF	C, C(5, C)	IFDATA
00015A	20 E4		ECR	C, 4	
00015C	28 00		SCR	C, 0	

00015E	38 06	LE	C, E		
000160	6A 00 0 050	AC	0, 8E(0,13)	4F09000000000000	
000164	60 C0 0 050	STD	0, 80(0,13)		
000168	58 20 0 054	L	2, 84(0,13)		
00016C	50 25 5 000	ST	2, 0(5, 5)		ICODE
000170	47 F0 0 1EC	HC	1E, 264(0,13)		20
000174	78 24 0 000	15 LE	2, C(4, C)		ADATA
000178	30 24	DER	2, 4		
00017A	38 C2	LEP	C, 2		
00017C	78 60 0 074	SE	C, 116(0,13)	41100000	
000180	28 00	SCR	C, C		
000182	38 06	LER	C, E		
000184	6A 00 0 050	AC	C, 8E(0,12)	4F08000000000000	
000188	60 00 0 050	STD	0, 80(0,13)		
00018C	58 20 0 054	L	2, 84(0,13)		
000190	5C 24 5 000	ST	2, 0(4, 5)		ICODE
000194	58 08 9 000	20 L	C, C(8, 5)		ICODE
000198	59 00 0 0A4	C	C, 164(0,12)		LULVL
00019C	47 C0 0 180	HC	12, 204(0,13)		100000
0001A0	58 00 0 0A4	100004 L	0, 164(0,13)		LULVL
0001A4	50 08 9 000	ST	C, C(8, 5)		ICODE
0001A8	5E 09 5 000	100005 L	C, 0(0, 0)		ICODE
0001AC	59 00 0 0A0	C	C, 160(0,13)		LULVL
0001B0	47 A0 0 154	EC	1C, 404(0,13)		IC
0001B4	5E 00 0 0A0	100006 L	C, 160(0,13)		LULVL
0001B8	50 08 9 000	ST	C, C(8, 5)		ICODE
0001BC	E7 7A 0 11E	10 FXLE	7,10, 226(13)		100007
0001C0	78 C0 0 094	100007 LE	C, 14E(0,13)		XMAX
0001C4	58 80 0 060	L	13, 56(0,12)		C
0001C8	18 FF	SR	1E,1E		
0001CA	58 E0 0 000	L	14, C(0,13)		
0001CE	07 FE	PCR	1E,14		
ADDRESS OF EPILOGUE					
0001D0	58 00 0 004	L	12, 4(0,12)		
0001D4	5E E0 0 00C	L	14, 12(0,13)		
0001D8	52 FF 0 00C	MVI	12(12),255		
0001DC	98 2C 0 01C	LM	2,12, 22(12)		
0001E0	C7 FF	ECR	1E,14		
ADDRESS OF PROLOGUE					
0001E2	58 7A 1 004	LM	7,1C, 4(1)		
0001E6	78 20 7 000	LF	2, 0(0, 7)		
0001EA	70 20 0 0AE	STE	2, 16E(0,13)		EFSLCF
0001EE	7E 20 0 000	LF	2, C(0, 2)		
0001F2	70 20 0 054	STE	2, 14E(0,13)		XMAX
0001F6	5E 20 A 000	L	2, 0(0,10)		
0001FA	50 20 0 0EC	ST	2, 140(0,12)		NUM
0001FE	5E 20 1 000	L	2, 0(0, 1)		
000202	41 30 2 000	LA	3, 0(0, 2)		
000206	41 50 0 004	LA	2, 4		
00020A	18 25	SR	2, E		
00020C	50 20 0 000	ST	2, 152(0,12)		
000210	50 30 0 0C4	ST	2, 19E(0,13)		ADATA
000214	58 20 1 00C	L	2, 12(0, 1)		
000218	41 30 2 000	LA	3, 0(0, 2)		
00021C	41 50 0 0C4	LA	E, 4		
000220	1E 25	SR	2, E		
000222	50 20 0 01E	ST	2, 184(0,12)		
000226	50 30 0 01C	ST	2, 104(0,13)		ICODE

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ADCLN FOR PROLOGUE	00022A 47 F0 D 0E0	DC	1E, 224(0,13)
	000020 00C001F2	DC	X1A*000001F2*
ADCLN FOR SAVE AREA	000024 00C0002B	DC	XL4*0C00002E*
ADCLN FOR EPILOGUE	00002B 00000100	DC	XL4*00000100*
TEMPORARIES AND GENERATED CONSTANTS	0000F0 00C00000	DC	XL4*00000000*
	0000F4 00C00004	DC	XL4*00000004*
	0000F8 00C00008	DC	XL4*00C00000*
	0000FC 00C0000C	DC	XL4*00000000*
	000100 00C00000	DC	X1A*0C000000*
	000104 00C00000	DC	XL4*00000000*

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINECOUNT(0) SIZE(0750K) AUTOCOL(NCF)

*OPTIONS IN EFFECT*SOURCE EBCCIC LIST NODFCK (OBJECT MAP FORMAT GDSIMP XREF NCALC NOANSF TERM IBM FLAG(1))

STATISTICS SOURCE STATEMENTS = 22, PROGRAM SIZE = 15E, SUBPROGRAM NAME = CDDER

STATISTICS NO DIAGNOSTICS GENERATED

***** END OF COMPILATION *****

216K BYTES OF CORE NOT USED

STATISTICS NO DIAGNOSTICS THIS STEP

REQUESTED OPTIONS: OPT=2,FCFMT,XREF,LIST,MAP,SIZE(750K)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LIFECLNT(60) SIZE(0750K) AUTOCCL(NONE)
SOURCE EBCCIC LIST NODECK OBJECT MAP FCFMT GDSMT PREF NCALC NUANSF TERM IOP FLAG(1)

```

ISN 0002      SURFLINE DECODE(ICODE,MAPPER,DELTA,RDATA,NUM)          00000290
              C  VARIABLE DESCRIPTION                                00000290
              C  ICODE : ENCODED CCT TRANSFORM COEF. OF SEISMIC DATA 00000300
              C  MAPPER : COMPRESSION FLAG                          00000310
              C  DELTA : STEP SIZE                                  00000320
              C  RDATA : DECODED DATA ARRAY                       00000330
              C                                                    00000340
ISN 0003      DIMENSION ICODE(1),MAPPER(1),RDATA(1)                 00000350
ISN 0004      DO 10 I=1,NUM                                          00000360
ISN 0005      RDATA(I)=0.0                                          0000027C
ISN 0006      IF (MAPPER(I) .EQ. 0) GO TO 10                        0000038C
ISN 0008      CODE=ICODE(I)                                         00000390
ISN 0009      RDATA(I)=DELTA*CODE                                   00000400
ISN 0010      10 CONTINUE                                          00000410
ISN 0011      RETURN                                               00000420
ISN 0012      END                                                  00000430
    
```

*****CFORAN CROSS REFERENCE LISTING*****

```

SYMBOL  INTERNAL STATEMENT NUMBERS
I       0004 0003 0006 000E 0005
NUM     0002 0004
CODE    0008 0009
DELTA   0002 0009
ICODE   0002 0003 0003
RDATA   0002 0003 0005 0005
DECODE  0002
MAPPER  0002 0003 0006
    
```

*****FORTRAN CROSS REFERENCE LISTING*****

```

LABEL  DEFINED  REFERENCES
10     0010    0004 0006

          000000 47 F0 F 00C          ECODE  DC  15,12(0,15)
          000004 07                    DC  XL1*07*
          000005 C4(5C3D6(4C540        DC  CL7*(CODE *
          00000C 50 EC D 00C          STM   14,12,12(13)
          000010 18 4D                LR    4,13
          000012 58 CD F 020          LW    12,13,32(15)
          000016 50 40 D 004          ST    4,4(0,13)
          00001A 50 D0 4 00B          ST    12,8(0,4)
          00001E 07 FC                ECR   12,12

TEMPORARY FOR FIX/FLOAT
00007B 4E000000          CC     XL4*4E000000*
00007C 00000000          CC     XL4*00000000*

CONSTANTS
000080 4E000000          CC     XL4*4E000000*
00008A 80000000          CC     XL4*80000000*
00008B 00000000          DC     XL4*00000000*
00008C 00000000          DC     XL4*00000000*
    
```

000070	00000001	DC	X14'CCCC0001'		
000054	00000000	EC	X14'00000000'		
000090	00000000	DC	X14'CCCC0000'		
ADJGNS FOR VARIABLES AND CONSTANTS					
ADJGNS FOR EXTERNAL REFERENCES					
000009	00000000	DC	X14'00000000'		ICCODE
0000C0	00000000	DC	X14'00000000'		BCDATA
0000C8	00000000	DC	X14'00000000'		MAFFER
0000DC	58 60 C 0A0	100001	L	6, 160(0,13)	
0000E0	58 70 D C90		L	7, 144(0,13)	
0000E4	58 80 D C90		L	8, 152(0,13)	
0000E8	78 40 D 000		LE	4, 128(0,13)	DELTA
0000EC	58 A0 D 0AC		L	10, 172(0,13)	4
0000F0	58 50 D 078		L	5, 120(0,13)	NUM
0000F4	18 25		LR	11, 2	
0000F6	65 80 D CC2		SLL	11, 2	
0000FA	18 9A		LR	5,10	
0000FC	78 00 D 0E4	100002	LE	6, 100(0,13)	0
000100	70 C5 E C00		STE	6, 0(9, 8)	BCDATA
000104	58 09 6 000		L	6, 0(9, 8)	MAFFER
000108	12 00		LTR	0, 0	
00010A	47 60 D 102		BC	8, 256(0,13)	10
00010E	58 09 7 000	100003	L	6, 0(9, 7)	ICCODE
000112	57 00 D C5C		X	0, 52(0,13)	4E000000E0000000
000116	50 00 D 054		ST	6, 84(0,13)	
00011A	68 60 D 050		LG	6, 80(0,13)	
00011E	68 60 D 05E		SC	6, 80(0,13)	4E000000E0000000
000122	3E 26		IER	2, 6	
000124	1C 24		MER	2, 4	
000126	70 25 B C00		STE	2, 0(9, 8)	BCDATA
00012A	67 5A D 0D4	10	EXLE	5,10, 212(13)	100002
00012E	18 FF	100004	SR	12,12	
000130	58 60 C 000		L	14, 0(0,13)	
000134	C7 FE		FCR	12,14	
ADDRESS OF EPILOGUE					
000136	58 00 D 004		L	12, 4(0,13)	
00013A	58 80 D C0C		L	14, 12(0,13)	
00013E	52 FF C 00C		MVI	12(13),255	
000142	5E 2C D C1C		LP	2,12, 20(13)	
000146	C7 FE		BCR	12,14	
ADDRESS OF PROLOGUE					
000148	5E 79 1 C0E		LP	7, 9, 8(13)	
00014C	7E 20 7 CCC		LE	2, 0(0, 7)	
000150	70 20 D 090		STE	2, 128(0,13)	DELTA
000154	5E 20 5 C00		L	2, 0(0, 9)	
000158	50 20 D 078		ST	2, 120(0,13)	NUM
00015C	5E 20 1 000		L	2, 0(0, 1)	
000160	41 20 2 C00		LA	2, 0(0, 2)	
000164	41 50 0 004		LA	5, 4	
000168	18 25		SR	2, 5	
00016A	50 20 D 650		ST	2, 144(0,13)	
00016E	50 30 D 094		ST	2, 148(0,13)	ICCODE
000172	5E 20 1 C04		L	2, 4(0, 1)	
000176	41 20 2 CCC		LA	2, 0(0, 2)	
00017A	41 50 0 004		LA	5, 4	
00017E	18 25		SR	2, 5	
000180	50 20 D 0A0		ST	2, 160(0,13)	
000184	50 30 C 0A4		ST	2, 164(0,13)	MAFFER

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000138	58 20 1 C0C	L	2, 12(0, 1)	
00016C	41 20 2 C00	LA	2, 0(0, 2)	
000190	41 20 0 C04	LA	2, 4	
000194	1E 25	SR	2, 5	
000196	20 20 0 C58	ST	2, 152(0, 13)	
00019A	50 20 0 05C	ST	2, 152(0, 13)	SEATA
00019E	47 F0 0 0R4	EC	15, 1PC(0, 12)	
ADLN FOR PROLOGUE				
000020	00000140	DC	XL4'00000140'	
ADLN FOR SAVE AREA				
000024	00000020	DC	XL4'00000020'	
ADLN FOR EPILOGUE				
000028	00000136	CC	XL4'00000136'	
TEMPORARIES AND GENERATED CONSTANTS				
000000	00000000	EC	XL4'00000000'	
000004	00000004	EC	XL4'00000004'	
000008	00000000	DC	XL4'00000000'	

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINECOUNT(C) SIZE(0750K) AUTODBL(NGNE)

*OPTIONS IN EFFECT*SOURCE EBCCIC LIST NDECK(OBJECT) MAP FORMAT GCSTMT XFF NCALC NOANSF TERM IBM FLAG(1)

STATISTICS SOURCE STATEMENTS = 11, PROGRAM SIZE = 416, SUBPROGRAM NAME =DECODE

STATISTICS NO DIAGNOSTICS GENERATED

***** END OF COMPILATION *****

3168 BYTES OF CORE NOT USED

REQUIRED OPTIONS: OPT=2,FORMAT,XREF,LIST,MAP,SIZE(750K)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LINECOUNT(CC) SIZE(0750K) PUTCEBL(NONE)
SOURCE EECIC LIST NOCHECK OBJECT MAP FERMAT GO(1M) XREF NCALC NOANSF TERM IBM FLAG(1)

```

ISN 0002      SUBROUTINE INVMEL(RDATA,NUM,XMAX,RMEU,EPSLN)      0000044C
              C      MEL-LAW INVERTER                        00000450
              C      RDATA      : INPUT ARRAY OF MEL-LAW CONVERTED TRACE DATA 00000460
              C      RMEU       : MEL-VALUE SET AS 50C        00000470
              C      XMAX       : ABSOLUTE MAXIMUM VALUE OF INPUT ARRAY      00000480
              C      NUM        : NUMBER OF DATA POINTS TO BE INVERTED      00000490
              C      EPSLN      : THRESHOLD VALUE FOR THE INSIGNIFICANT CCT COEF. 00000500
              C
ISN 0003      DIMENSION RDATA(1)                            00000510
ISN 0004      COEF=ALGG(RMEU+1)/XMAX                        00000520
ISN 0005      COEF1=XMAX/RMEU                              00000530
ISN 0006      DO 10 I=1,NUM                                00000540
              C      10 INCREASE COMPRESSED VALUE                00000560
              C      IF(RDATA(I) .GE. 0.0) CG 10 10          00000570
ISN 0007      RVAL=RDATA(I)*COEF                          00000580
ISN 0009      SIGN=1.0                                     00000590
ISN 0010      IF(RVAL .LT. 0.0) SIGN=-1.0                 00000600
ISN 0011      RVAL=AES(RVAL)                               00000610
ISN 0013      RVAL=EXP(RVAL)                              00000620
ISN 0014      RVAL=RVAL - 1.0                             00000630
ISN 0015      RVAL=RVAL*COEF1                             00000640
ISN 0016      RDATA(I)=RVAL*SIGN                          00000650
ISN 0017      IF(RDATA(I) .LT. 0.0) RDATA(I)=RDATA(I)+EPSLN 00000660
ISN 0018      IF(RDATA(I) .GE. 0.0) RDATA(I)=RDATA(I)+EPSLN 00000670
ISN 0020      CONTINUE                                    00000680
ISN 0022      10 CONTINUE                                 00000690
ISN 0023      RETURN                                      00000700
ISN 0024      END
    
```

*****FORTRAN CROSS REFERENCE LISTING*****

```

SYMBOL  INTERNAL STATEMENT NUMBERS
I       0006 0007 0005 0017 0018 0018 0020 0020 0020
AHS     0013
EXF     0014
NUM     0002 0006
ALGG    0004
COEF    0004 0009
RMEU    0002 0004 0005
RVAL    0005 0011 0013 0013 0014 0014 0015 0015 0016 0016 0017
SIGN    0010 0011 0017
XMAX    0002 0004 0005
COEF1   0005 0016
EPSLN   0002
RDATA   0002 0003 0007 0005 0017 0018 0018 0018 0020 0020 0020
EPSLNH  0016 0020
INVMEL  0002
    
```

*****FORTRAN CROSS REFERENCE LISTING*****

```

LABEL  DEFINED  REFERENCES
10     0022    0006 0007
    
```

000000	47 F0 F 00C	INVMU	BC	15,12(0,15)
000004	07		DC	XL1*(7)
000005	C5E5E504C5E440		CC	CL7*INVMU*
00000C	90 EC D 00C		STM	14,12,12(13)
000010	18 40		LR	4,12
000012	5E CD F 020		LR	12,13,22(15)
000016	50 40 D 004		ST	4,4(0,13)
00001A	50 00 4 003		ST	12,2(0,4)
00001E	C7 FC		ECR	15,12

CONSTANTS

000080	00C00000		CC	XL4*00000000*
000084	CCCC0001		DC	XL4*CCCC0001*
000088	41100000		DC	XL4*41100000*
00008C	0CC00000		CC	XL4*00000000*
000090	00C00000		DC	XL4*CCCC0000*

ADDONS FOR VARIABLES AND CONSTANTS

ADDONS FOR EXTERNAL REFERENCES

000008	00C00000		DC	XL4*CCCC0000*	EXP
00000C	00C00000		CC	XL4*00000000*	ALOG
000000	0CC00000		DC	XL4*00000000*	REATA
0000F8	58 80 D 0A8	100001	L	8, 16(0,13)	
0000FC	5E A0 D 0C0		L	10, 192(0,13)	A
000100	58 70 D 070		L	7, 112(0,13)	NUM
000104	78 00 D 060		LE	C, 5(0,13)	41100000
000108	7A 00 D C78		AF	C, 120(0,13)	FMEL
00010C	70 00 D 0C8		STE	C, 100(0,13)	100
000110	58 F0 D CA4		L	15, 164(0,13)	ALG
000114	41 10 D C4C		LA	1, 7(0,13)	
000118	05 EF		BALR	14,15	
00011A	47 00 D C04		EC	0, 4(0, 0)	
00011E	70 00 D 0E4		DE	C, 132(0,13)	XMAX
000122	70 00 D 074		STE	C, 11(0,13)	CCFF
000126	7E 00 D 0E4		LE	0, 132(0,13)	XMAX
00012A	70 00 D 07E		DE	C, 120(0,13)	FMEL
00012E	70 00 D 038		STE	0, 136(0,13)	COFF1
000132	78 00 D C60		LE	C, 5(0,13)	41100000
000136	33 00		LCER	C, 0	
000138	70 00 D 08C		STE	0, 188(0,13)	100
00013C	18 B7		LR	11, 7	
00013E	85 80 0 002		SLL	11, 2	
000142	1E 5A		LR	5,10	
000144	7E 29 8 000	100002	LE	2, C(5, 8)	FLA1A
000148	70 20 D 090		STE	2, 17(0,13)	1001
00014C	22 22		LTER	2, 2	
00014E	47 80 D 19A		DC	8, 410(0,13)	10
000152	7C 20 D 074	100003	MF	2, 11(0,13)	CCFF
000156	70 20 D C7C		STE	2, 124(0,13)	FVAL
00015A	78 00 C 060		LE	C, 5(0,13)	41100000
00015E	70 00 D C80		STE	C, 12(0,13)	SIGN
000162	22 22		LTER	2, 2	
000164	47 A0 E 148		EC	10, 128(0,13)	100005
000168	78 00 D 08C	100004	LE	0, 188(0,13)	1000
00016C	70 00 E C80		STE	C, 12(0,13)	SIGN
000170	7E CD D C7C	100005	LE	0, 124(0,13)	FVAL
000174	20 00		LTER	C, C	
000176	70 00 E 07C		STE	C, 124(0,13)	FVAL
00017A	5E F0 D 0A0		L	15, 140(0,13)	EXP

00017E	41 10 D 0E0	LA	1, 80(0,12)	
000182	05 FF	BAIR	14,15	
000184	47 00 0 00E	RC	0, 14(0, 0)	
000188	38 20	LER	2, C	
00018A	7B 20 D 0E0	SE	2, 5(C,12)	41100C0C
00018E	7C 20 D C99	ME	2, 13(0,13)	COFF1
000192	7C 20 D 0E0	MF	2, 12(0,13)	SIG1
000196	70 29 8 000	STE	2, 0(9, 2)	RDATA
00019A	22 22	LTER	2, 2	
00019C	47 A0 0 1E4	RC	10, 20(0,13)	100007
0001A0	7E 25 8 C00	LE	2, 0(9, 2)	RDATA
0001A4	7B 20 D C50	SE	2, 144(0,13)	EPSEL1
0001A8	70 29 8 000	STE	2, 0(9, 2)	PCATA
0001AC	7E 25 8 C00	LE	2, 0(9, 2)	PCATA
0001B0	70 20 D 018	STE	2, 184(0,13)	CO2
0001B4	32 22	LTER	2, 2	
0001B6	47 50 D 19A	RC	5, 410(0,13)	IC
0001BA	7A 20 D 050	AE	2, 144(0,13)	EPSEL1
0001BE	70 29 8 C00	STE	2, 0(9, 2)	RDATA
0001C2	87 9A D 11C	EXLE	5,10, 204(13)	100002
0001C6	1E FF	SR	15,15	
0001C8	5E E0 0 C00	L	14, 0(0,13)	
0001CC	07 FE	RCR	15,14	
ADDRESS OF EPILOGUE				
0001CE	5E A0 D C04	L	10, 4(0,13)	
0001D2	5B E0 A 00C	L	14, 12(0,10)	
0001D6	5B 80 A 018	L	11, 24(0,10)	
0001DA	5B 10 B C0C	L	1, 12(0,11)	
0001DE	7B 20 D C78	LE	2, 120(0,13)	EP1
0001E2	70 20 1 C00	STE	2, 0(0, 1)	
0001E6	1B 0A	LP	12,10	
0001E8	92 FF A C0C	MVI	12(10),255	
0001EC	5B 2C A C1C	LM	2,12, 28(10)	
0001F0	07 FE	RCR	15,14	
ADDRESS OF PROLOGUE				
0001F2	5B 7A 1 C04	LM	7,10, 4(1)	
0001F6	5B 20 7 000	L	2, 0(0, 7)	
0001FA	50 20 D C70	ST	2, 112(0,13)	NUM
0001FE	7E 20 8 000	LE	2, 0(0, 8)	
000202	70 20 D 0F4	STE	2, 132(C,12)	KNAX
000206	7E 20 9 000	LE	2, 0(0, 9)	
00020A	70 20 D 078	STE	2, 120(0,13)	EP1
00020E	7B 20 A 000	LE	2, C(0,10)	
000212	70 20 D C8C	STE	2, 140(0,13)	EP50K
000216	5B 20 1 000	L	2, 0(0, 1)	
00021A	41 20 2 C00	LA	2, 0(0, 2)	
00021E	41 50 0 C04	LA	5, 4	
000222	1E 25	SR	2, 2	
000224	50 20 0 0A8	ST	2, 16(0,13)	
000228	50 30 0 CAC	ST	3, 172(0,13)	FLATA
00022C	47 F0 D 0C0	FC	15, 20(0,13)	
ADDN FOR PROLOGUE				
000020	00001F2	DC	X14*00001F2*	
ADDN FOR SAVE AREA				
000024	000002B	DC	X14*000002B*	
ADDN FOR EPILOGUE				
000028	00001CE	DC	X14*00001CE*	
ADDNS FOR PARAMETER LISTS				

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000074	00C000F0	CC	XL4*E000C0F0*	.100
000078	00C000A4	CC	XL4*E00000A4*	EVAL
TEMPORARIES AND GENERATED CONSTANTS				
000008	00C00000	DC	XL4*00000000*	
00000C	00C00000	CC	XL4*00000000*	
0000E0	00C00000	DC	XL4*00000000*	
0000E4	00C00000	DC	XL4*CCCC0000*	
0000E8	00C00004	DC	XL4*00000004*	
0000EC	00C00000	DC	XL4*CCCC0000*	
0000F0	00C00000	CC	XL4*00000000*	
0000F4	00C00000	DC	XL4*00000000*	

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINECOUNT(60) SIZE(0750K) ALIGCD(LINE)

*OPTIONS IN EFFECT*SOURCE EBCDIC LIST NOCHECK OBJECT MAP FORMAT (0)SYMT XREF NCALC NOANSI TERM IBM FLAG(1)

STATISTICS SOURCE STATEMENTS = 23, PROGRAM SIZE = 860, SUBPROGRAM NAME = INVMED

STATISTICS NO DIAGNOSTICS GENERATED

***** END OF COMPILATION *****

316K BYTES OF CORE NOT USED

REQUESTED OPTIONS: OPT=2,FORMAT,XREF,LIST,MAP,SIZE(756K)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LINECOUNT(20) SIZE(0750K) ALIASEDL(NONE)
SOURCE EBCLIC LIST NOCHECK OBJECT MAP FORMAT GCSTMT XREF NCALC NOANSF TERM IBM FLAG(1)

```

C
ISN 0002      C      SUBROUTINE FIDCT(A,NUM)                                00000710
C              INVERSE DCT USING FFTCC (IBM MATH PACKAGE)                00000720
C              NUM : NUMBER OF DATA POINTS                             00000730
C              A   : DATA TO BE TRANSFORMED                             00000740
C              NUM : NUMBER OF DATA POINTS                             00000750
C              WK  : WORKING AREA OF FFTCC                                00000760
C              WK  : WORKING AREA OF FFTCC                                00000770
C              X   : INPUT AND OUTPUT ARRAY OF FFTCC                     00000780
C              X   : INPUT AND OUTPUT ARRAY OF FFTCC                     00000790
C
ISN 0003      C      DIMENSION WK(2600),WK(2600),A(1)                    00000800
ISN 0004      C      COMPLEX X(4750)                                     00000810
ISN 0005      C      WWP=3.141592/FLCAT(2*NUM)                            00000820
ISN 0006      C      DO 10 I=1,NUM                                         00000830
ISN 0007      C      ARG=WRW*FLCAT(I-1)                                     00000840
ISN 0008      C      X(I)=CMPLX(A(I),0.0)*CMPLX(COS(ARG),SIN(ARG))        00000850
ISN 0009      C      10 CONTINUE                                           00000860
ISN 0010      C      CALL FFTCC(X,NUM,WK,WK)                               00000870
ISN 0011      C      A(I)=REAL(X(I))                                       00000880
ISN 0012      C      20 CONTINUE                                           00000890
ISN 0013      C      RETURN                                               00000900
ISN 0014      C      END                                                  00000910
    
```

```

***** FORTRAN CROSS REFERENCE LISTING *****
SYMBOL  INTERNAL STATEMENT NUMBERS
A       0002 0003 0006 0011
I       0006 0007 0008 000E 0011 0011
X       0004 0008 0010 0011
WK      0003 0010
ARG     0007 0008 0008
COS     000E
WK      0003 0010
NUM     0002 0005 0006 0010
SIN     0008
REAL    0011
WRW     0005 0007
CMPLX   000E 0008
FFTCC   0010
FIDCT   0002
FLAG    000E 0007
    
```

```

***** FORTRAN CROSS REFERENCE LISTING *****
LABEL  DEFINED  REFERENCES
10     0009  0006
20     0012

000000  47 F O I  C O C          F I D C T      DC   15,12(0,15)
000004  07                                DC   X(1),C7*
000005  C6C5C4C3E340A0          DC   CL7*FIDCT *
00000C  50 E C D  C O C          STM  14,12,12(13)
000010  13 4D                    LR   4,12
    
```

000012 90 00 F 020
 000016 50 00 D 004
 00001A 50 00 4 008
 00001E 07 FC

LP 12.12.22(1E)
 ST 4.4(0.12)
 ST 12.2(0.4)
 BCR 11.12

TEMPORARY FOR FIX/FLOAT

000090 4E000000
 000054 00100000

DC XL4*4E000000*
 FC XL4*00000000*

CONSTANTS

000098 4E000000
 00009C 80000000
 0000A0 00000000
 0000A4 00000001
 0000A8 00000002
 0000AC 413243F6
 0000B0 00000000
 0000B4 00000000

DC XL4*4E000000*
 DC XL4*80000000*
 DC XL4*00000000*
 DC XL4*00000001*
 DC XL4*00000002*
 DC XL4*413243F6*
 DC XL4*00000000*
 DC XL4*00000000*

ADDNS FOR VARIABLES AND CONSTANTS

00E680 0002568
 00E684 00005208

DC XL4*00002568*
 DC XL4*00005208*

ADDNS FOR EXTERNAL REFERENCES

00E688 00000000
 00E690 00000000
 00E694 00000000
 00E698 00000000
 00E69C 00000000

DC XL4*00000000*
 DC XL4*00000000*
 DC XL4*00000000*
 DC XL4*00000000*
 DC XL4*00000000*

A
 CCS
 SIN
 FITCC
 CMFY4

00E604 58 50 C 008
 00E608 58 60 C 004
 00E60C 58 80 C 000
 00E6E0 58 A0 D 07C
 00E6E4 58 40 C 024
 00E6E8 58 30 C 02C
 00E6EC 58 20 D 058
 00E6F0 85 20 0 001
 00E6F4 18 02
 00E6F6 57 00 D 074
 00E6FA 50 00 D 06C
 00E6FE 6E 20 D 068
 00E702 6B 20 D 070
 00E706 73 00 D 084
 00E70A 30 02
 00E70C 70 00 D 09C
 00E710 19 2A
 00E712 50 A0 D 090
 00E716 78 00 D 078
 00E71A 70 00 C 020
 00E71E 50 40 C 028
 00E722 50 30 C 030
 00E726 5E 70 C 030
 00E72A 58 80 C 028
 00E72E 18 92
 00E730 5E 80 D 058
 00E734 18 29
 00E736 18 2A
 00E738 18 02
 00E73A 57 00 D 074
 00E73E 50 00 D 06C
 00E742 68 20 D 068

100001

L 5, 8(0.12)
 L 6, 4(0.12)
 L 11, 0(0.12)
 L 10, 124(0.13)
 L 4, 30(0.12)
 L 2, 44(0.12)
 L 2, 152(0.13)
 SLL 2, 1
 LR 0, 2
 X C, 116(0.13)
 ST C, 10E(0.12)
 LD 2, 104(0.13)
 SD 2, 112(0.13)
 LE C, 132(0.13)
 DER C, 2
 STE C, 150(0.12)
 LR 2, 10
 ST 1C, 144(0.13)
 LE C, 120(0.12)
 STE 0, 32(0.12)
 ST 4, 40(0.12)
 ST 2, 4E(0.12)
 L 7, 48(0.12)
 L 6, 40(0.12)
 LR 5, 2
 L 11, 152(0.13)
 LR 2, 5
 SR 2, 10
 LR C, 2
 X C, 110(0.12)
 ST C, 10E(0.12)
 LD 2, 104(0.12)

4E00000080000000

4E00000080000000

413243F6

WFWF

1

0

+C01

+C02

+C03

+C04

+C02

NUM

4E00000080000000

00E746	68 20 D 070	SD	2, 112(0,12)	4E0000C0E0000000
00E74A	7C 20 D C5C	ME	2, 156(0,12)	NRWF
00E74E	70 20 D C94	STE	2, 148(0,13)	ARC
00E752	78 27 5 000	LE	2, 0(7, 8)	A
00L756	70 20 C C3E	STE	2, 56(0,12)	.T01
00E75A	5E F0 C 010	L	1E, 16(0,12)	CCS
00E75E	41 10 D C4C	LA	1, 76(0,13)	
00E762	05 EF	BALR	14,1E	
00E764	47 00 0 008	BC	0, 8(0, 0)	
00E768	70 00 C 034	STE	0, 52(0,12)	.T00
00L76C	58 F0 C 014	L	1E, 20(0,12)	FIN
00E770	41 10 D 04C	LA	1, 76(0,13)	
00E774	C5 EF	BALR	14,1E	
00E776	47 00 0 CC8	BC	0, 8(0, 0)	
00E77A	70 00 C 040	STE	0, 64(0,12)	.T02
00E77E	7E 00 C C34	LE	0, 52(0,12)	.T0C
00L782	70 00 C C48	STE	0, 72(0,12)	.T0E
00E786	78 00 C 040	LE	0, 64(0,12)	.T02
00E78A	70 00 C 04C	STE	0, 76(0,12)	.T0C
00L78E	78 00 C 020	LE	0, 32(0,12)	.C01
00E792	70 00 C 02C	STE	0, 60(0,12)	.T0E
00L796	58 F0 C 01C	L	1E, 28(0,12)	CPFY#
00E79A	41 10 D C50	LA	1, 80(0,13)	
00E79E	05 EF	BALR	14,1E	
00E7A0	47 00 0 008	BC	0, 8(0, 0)	
00L7A4	70 00 C 034	STE	0, 52(0,12)	.T00
00E7A8	70 28 6 004	STE	2, 4(8, 6)	X
00L7AC	78 00 C C34	LE	0, 52(0,12)	.T0C
00E7B0	70 C8 6 C00	STE	0, 0(8, 6)	X
00E7E4	1A 73	10 AR	7, 3	
00E7B6	1A EA	AR	E, 4	
00L7B8	87 9A C 014	DXLE	5,10, 1E0(12)	100002
00E7BC	50 98 D 090	ST	5, 144(0,12)	I
00L7C0	58 80 C 000	100003 L	11, 0(0,12)	
00L7C4	58 F0 C 018	L	1E, 24(0,12)	FFTC
00L7C8	41 10 D C5E	LA	1, 88(0,13)	
00E7CC	05 EF	BALR	14,1E	
00E7CE	47 00 0 00A	BC	0, 10(0, 0)	
00E7D2	5E 20 D C50	L	2, 144(0,13)	I
00E7D6	89 20 0 003	SLL	2, 3	
00L7DA	78 22 6 C00	LE	2, 0(2, 6)	X
00E7DE	78 02 6 C04	LF	0, 4(2, 6)	X
00L7E2	58 20 D 090	L	2, 144(0,13)	I
00L7E6	89 20 0 C02	SLL	2, 2	
00L7EA	70 22 5 C00	STE	2, 0(2, 8)	/
00L7EE	1E FF	20 SR	1E,1E	
00L7F0	58 E0 D C00	L	1A, 0(0,12)	
00E7F4	07 FE	OCR	1E,14	
ADDRESS OF EPILOGUE				
00E7F6	58 A0 D CC4	L	10, 4(0,13)	
00E7FA	58 E0 A 00C	L	14, 12(0,1C)	
00E7FE	5E B0 A C18	L	11, 24(0,1C)	
00L802	58 10 0 C04	L	1, 4(0,11)	
00L806	58 20 D C9A	L	2, 152(0,13)	NIA
00E80A	50 20 I 000	ST	2, 0(0, 1)	
00L80E	18 DA	LC	1E,1C	
00L810	92 FF A 00C	MVI	12(1C),2EE	
00L814	5E 2C A 01C	LM	2,12, 20(10)	

ADDRESS OF PROLOGUE	OPERATOR	OPERAND	OPERATION	OPERAND	OPERATION	OPERAND	OPERATION	OPERAND
00E810	C7 FE		PCP	15,14				
00E81A	22 C0 D C48		L	12, 72(0,13)				
00E81E	52 70 1 C04		L	7, 4(0, 11)				
00E822	5E 20 7 000		L	2, C(0, 7)				
00E826	50 20 0 058		ST	2, 152(0,13)				NUM
00E82A	5E 20 1 000		L	2, C(0, 11)				
00E82E	41 30 2 000		LA	3, C(0, 2)				
00E832	41 50 0 004		LA	1, 4				
00E836	1B 25		SR	2, 5				
00E83B	20 20 C CCE		SI	2, 2(0,12)				
00E83C	50 30 C 00C		SI	3, 12(0,12)				A
00E840	47 F0 C C54		PC	15, 84(0,12)				
ADCN FOR PROLOGUE	000020	00C0E01A	DC	XL4*C0C0E01A*				
ADCN FOR SAVE AREA	000024	00C0002B	DC	XL4*C000C02B*				
ADCN FOR EPILOGUE	000028	00C0E7F6	DC	XL4*C000E7F6*				
ADCN FOR REG 12	000070	00C0E6B0	DC	XL4*0000E6B0*				
ADCN FOR PARAMETER LISTS	000074	00C000DC	DC	XL4*E0C000DC*				ARG
	000078	00C0E6B8	CC	XL4*C000E6B8*				.101
	00007C	0000E6C8	DC	XL4*E0C0E6C8*				.105
	000080	00005210	DC	XL4*00005210*				X
	000084	00C000C0	CC	XL4*000000C0*				NUM
	000088	00C0296C	DC	XL4*0000296C*				INX
	0000EC	00C000CC	CC	XL4*000000CC*				BY
TEMPORARIES AND GENERATED CONSTANTS	00E6A0	00C00000	DC	XL4*CCC00000*				
	00E6A4	C0C0000E	CC	XL4*0000000E*				
	00E6A8	00C00000	DC	XL4*C0C00000*				
	00E6AC	00000004	CC	XL4*00C0C004*				
	00E6B0	00C00000	CC	XL4*C0000000*				
	00E6B4	00C00000	DC	XL4*CCC000C0*				
	00E6B8	00C00000	CC	XL4*00000000*				
	00E6BC	00C00000	DC	XL4*C0000000*				
	00E6C0	00000000	DC	XL4*C000C0C0*				
	00E6C4	00C00005	DC	XL4*00E00005*				
	00E6C8	00580005	DC	XL4*C05E0005*				
	00E6CC	00C00000	DC	XL4*C0000000*				
	00E6D0	00C00006	DC	XL4*C0C00006*				

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINECOUNT(20) SIZE(0750K) AUTOCCEL(ON)

*OPTIONS IN EFFECT*SOURCE EECIC LIST NOCHECK OBJECT MAP FORMAT GC5INT PDEF NCALC NOANSF TERM (0) FLAC(1)

STATISTICS SOURCE STATEMENTS = 13, PROGRAM SIZE = 554CC, SUBPROGRAM NAME = FIDCT

STATISTICS NO DIAGNOSTICS GENERATED

***** END OF COMPILATION *****

212K BYTES OF CORE NOT USED

STATISTICS NO DIAGNOSTICS THIS STEP

REQUESTED OPTIONS: OPT=2,FORMAT,XREF,LIST,MAP,SIZE(750K)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LINECOUNT(C) SIZE(0750K) AUTOCORREL(C)
 SOURCE EXECIC LIST NCECK CEJECT MAP FORMAT GOEINT XREF NCALC NOANSF TRM IHM FLAG(1)

```

C
C
C
ISN 0002 C      SUBROUTINE ALICE(N,MM,R,A,ALPHA,RC)
C
C      THIS PROGRAM CALCULATES THE FILTER COEFFICIENTS FOR
C      A LINEAR PREDICTOR OF ORDER 'MM'. THE REFLECTION COEFFICIENTS
C      ARE SAVED FOR A LATTICE IMPLEMENTATION OF THE INVERSE FILTER.
C      THIS ROUTINE IS GIVEN BY MARKEL AND GRAY (LINEAR PREDICTION OF
C      SPEECH).
C
C      N - NO. OF POINTS
C      X - VECTOR OF SAMPLED SPEECH PTS
C      M - ORDER OF INVERSE FILTER
C      A - VECTOR OF INVERSE FILTER COEFFICIENTS
C      ALPHA - VECTOR OF CROSS CORRELATION COEFFICIENTS
C      RC - VECTOR OF REFLECTION COEFFICIENTS
C      R - AUTOCORRELATION COEFFICIENTS
C      ARRAY - VECTOR TO HOUSE ALL COEFFICIENTS CALCULATED
C
C
ISN 0003 C      INTEGER NDCWF,ALICE,LF
ISN 0004 C      DATA LP/E/,NDCWF/E/,ALICE/S/
ISN 0005 C      DIMENSION F(20),I(256),X(256),A(20),ALPHA(20),RC(20),
C      $      ARRAY(44)
ISN 0006 C      M=MM
ISN 0007 C      MF = M + 1
ISN 0008 C      DO 15 K = 1,MC
C
C      R(K) = 0.0
C      L = N - K + 1
C
C      DO 10 NF = 1,L
C      NFK = NF+K-1
C      R(K) = F(K)+X(NF)*X(NFK)
C      CONTINUE
C
ISN 0015 C      15 CONTINUE
C
C      A(1) = 1.
C      ALPHA(1) = 0(1)
C      IF(M.EQ.0) GO TO 60
C      RC(1) = -R(2)/R(1)
C      A(2) = R(1)
C      ALPHA(2) = R(1) + F(2)*R(1)
C      IF(M.EQ.1) GO TO 60
C
C
ISN 0025 C      DO 40 M(KC) = 2,M
C
ISN 0026 C      S = 0.0
    
```

```

C
ISN 0027          DC 20 IP = 1,MINC          00002250
ISN 0028          MNC=MINC-IP+2          00002260
ISN 0029          S = S + MNC)*A(IP)      00002270
ISN 0030          CONTINUE                00002280
C
ISN 0031          PC(MINC) = -S/ALPHA(MINC) 00002290
ISN 0032          MH = MINC/2 + 1         00002300
C
ISN 0033          DO 30 IP = 2,MH         00002310
ISN 0034          IP = MINC - IP + 2     00002320
ISN 0035          AT = A(IP) + FC(MINC)*A(IP) 00002330
ISN 0036          A(IP) = A(IP) + FC(MINC)*A(IP) 00002340
ISN 0037          A(IP) = A(IP) + FC(MINC)*A(IP) 00002350
ISN 0038          CONTINUE                00002360
C
ISN 0039          A(MINC+1) = FC(MINC)     00002370
ISN 0040          ALPHA(MINC+1) = ALPHA(MINC)-ALPHA(MINC)*PC(MINC)*FC(MINC) 00002380
C
ISN 0041          IF (ALPHA(MINC)) 60,60,40 00002390
C
ISN 0042          40 CONTINUE              00002400
C
ISN 0043          60 CONTINUE              00002410
C
ISN 0044          RETURN                   00002420
ISN 0045          END                     00002430
    
```

***** F O R T R A N C R O S S R E F E R E N C E L I S T I N G *****

SYMBOL	INTERNAL STATEMENT NUMBERS															
A	0002	0003	0010	0011	0025	0025	0035	0036	0036	0037	0039					
K	0008	0009	0010	0012	0013	0013										
L	0010	0011														
M	0006	0007	0018	0022	0025											
N	0002	0010														
R	0002	0003	0009	0012	0013	0017	0020	0020	0022	0022	0025					
S	0020	0029	0029	0031												
A	0002	0003	0013	0013												
AT	0035	0037														
IP	0034	0035	0036	0036												
IP	0027	0028	0029	0032	0034	0035	0036	0037								
IX	0005															
LP	0003	0004														
MH	0032	0033														
MM	0002	0006														
NP	0007	0008														
NP	0011	0012	0013													
PC	0002	0003	0020	0021	0022	0031	0035	0036	0035	0040	0040					
MNC	0028	0025														
MPK	0012	0013														
ALPH	0002															
MINC	0025	0027	0028	0031	0031	0032	0034	0035	0036	0035	0035	0040	0040	0040	0040	0041
ALPHA	0002	0003	0017	0022	0031	0040	0040	0040	0041							
ALPHA	0005															
AUTC	0003	0004														
END	0003	0004														

***** C O N T A I N E R C O D E S R E F E R E N C E L I S T I N G *****

LABEL	DEFINED	REFERENCES
10	0014	0011
11	0007	
15	0015	0008
20	0030	0027
30	0030	0033
40	0042	0025 0041
50	0043	0018 0023 0041 0041

000000	47 F0 F 00C	ALPH	PC	15,12(0,15)
000004	C7		CC	XL1*07*
000005	C1E4E7D(404040)		CC	CL1*ALIC
00000C	90 EC D 00C		STM	14,12,12(12)
000010	1E 4D		LP	4,13
000012	58 CD F 020		LM	12,12,32(15)
000016	50 40 D C04		ST	4,4(0,13)
00001A	50 DC 4 (CB		ST	12,8(0,4)
00001E	07 FC		HCR	15,12

CONSTANTS

000078	00C00000	DC	XL4*C0C0C0C0*
00007C	00C00000	DC	XL4*C0000000*
000080	00C00001	DC	XL4*C0C0C001*
000084	00C00002	DC	XL4*C0C0C0C2*
000088	411C0000	CC	XL4*4110C000*
00008C	00C00000	DC	XL4*C0C0C000*
000090	00C00000	CC	XL4*00C00000*

ALCONS FOR VARIABLES AND CONSTANTS
ALCONS FOR EXTERNAL REFERENCES

0000E0	00C00000	CC	XL4*00000000*	A
0000E8	00C00000	DC	XL4*00000000*	F
0000F0	00000000	DC	XL4*C0C0C0C0*	B
0000F8	00C00000	CC	XL4*00000000*	HC
000100	00C00000	DC	XL4*C0C0C0C0*	ALFFZ
000120	58 70 D 000	100071	L	7, 21(0, 13)
000124	58 EC D 000		L	8, 20(0, 13)
000128	58 50 D 00E		L	5, 18(0, 13)
00012C	58 E0 D 0C0		L	11, 19(0, 13)
000130	5E A0 D CEB		L	10, 8(0, 13)
000134	58 E0 D 0E4		L	6, 22(0, 13)
000138	5E 20 D C5C		L	2, 14(0, 13)
00013C	50 20 E 074		ST	2, 11(0, 13)
000140	1A 2A	11	AF	2,10
000142	5C 20 D 054		ST	2, 14(0, 13)
000146	50 60 D 0F4		ST	6, 24(0, 13)
00014A	1E 3A		LF	2,10
00014C	5E 70 D 078		L	7, 12(0, 13)
000150	18 E2		LR	8, 2
000152	18 5E		LF	5, 6
000154	78 E0 D 054		LF	6, 8(0, 13)
00015E	1E 5H		LR	5,11
00016A	58 E0 D C14	100072	L	14, 24(0, 13)
00016E	78 E5 000		STM	6, C(14, 5)
000162	18 27		LF	2, 7
000164	1H 23		SR	2, 3

000166	18 EA		LR	11.10	
000169	1A P2		FR	11. 2	
00016A	58 40 D CFA		L	4. 244(0.13)	.C04
00016E	18 89		LR	6. 5	
000170	1E 5A		LR	5.10	
000172	58 70 D 0C8		L	7. 200(0.13)	
000176	18 29	100003	LR	2. 5	
000178	1A 23		FR	2. 3	
00017A	1E 22		LR	6. 2	
00017C	1E 5A		SR	6.10	
00017E	78 28 7 C0C		LE	2. 0(0. 7)	
000182	18 26		LR	2. 6	
000184	89 20 0 002		SUI	2. 2	
00018P	7C 22 7 C00		WE	2. 0(2. 7)	
00018C	7A 24 5 000		AE	2. 0(4. 5)	F
000190	70 24 5 000		STE	2. 0(4. 5)	F
000194	5A E0 D 0C4	10	A	6. 220(0.13)	4
000198	87 9A D 14E		RMLE	5.10. 324(13)	100002
00019C	5E 50 D 0E4	15	L	5. 220(0.13)	4
0001A0	58 80 D 054		L	6. 140(0.13)	MF
0001A4	58 70 D 078		L	7. 120(0.13)	N
0001A8	58 00 D CFA		L	6. 244(0.13)	.C04
0001AC	1A 09		AR	6. 5	
0001AE	50 00 D 0FA		ST	6. 244(0.13)	.C04
0001B2	1A 2A		AR	3.10	
0001B4	15 38		CR	2. 6	
0001B6	47 C0 D 132		TC	12. 200(0.13)	100002
0001BA	5E 70 D 008	100005	L	7. 210(0.13)	
0001BE	58 80 D 000		L	6. 200(0.13)	
0001C2	5E 50 D 08E		L	5. 184(0.13)	
0001C6	58 E0 D 0C0		L	11. 192(0.13)	
0001CA	58 60 D 0EA		L	6. 220(0.13)	4
0001CE	78 C0 D C6C		LE	6. 50(0.13)	4110000
0001D2	70 00 9 004		STE	6. 4(6. 5)	A
0001D6	7E 20 0 C04		LE	2. 4(0.11)	F
0001DA	70 20 7 C04		STE	2. 4(0. 7)	ALPHA
0001DE	58 00 D C74		L	6. 110(0.13)	N
0001E2	12 00		LTR	6. 0	
0001E4	47 80 D 22E		UC	6. 500(0.13)	60
0001E8	78 20 0 C08	100006	LE	2. 6(0.11)	F
0001EC	78 40 0 C04		LE	4. 4(0.11)	F
0001F0	38 62		LEP	6. 2	
0001F2	20 64		REP	6. 4	
0001FA	33 06		LCER	6. 6	
0001F6	70 C0 8 C04		STF	6. 4(0. 6)	60
0001FA	7E C0 8 C04		LF	6. 4(0. 8)	60
0001FE	70 60 9 C08		STE	6. 4(6. 5)	7
000202	3C 26		WER	2. 6	
000204	2A 24		AER	2. 4	
000206	70 20 7 008		STE	2. 6(0. 7)	ALPHA
00020A	55 40 D C74		C	10. 110(0.13)	N
00020E	47 80 D 22E		UC	6. 500(0.13)	60
000212	58 00 D C5C	100007	L	6. 92(0.13)	2
00021E	50 C0 D CAP		ST	6. 160(0.13)	MINC
00021A	58 00 D 0FD		L	6. 244(0.13)	0
00021E	50 00 D 0FC		ST	6. 230(0.13)	.C02
000222	58 20 0 C5C		L	2. 62(0.13)	2
000226	19 26		LE	5. 7	

00022B	10 P7	LR	11. 7	
00022A	58 80 0 018	L	6. 184(0.13)	
00022F	58 70 0 000	L	7. 208(0.13)	
000232	78 20 0 C54	100009 LE	2. 84(0.13)	C
000236	70 20 0 07C	SE	2. 124(0.13)	S
00023A	18 29	LR	2. 5	
00023C	18 49	LR	4. 5	
00023E	19 72	LR	7. 2	
000240	1E 8A	LR	E.10	
000242	5E 80 0 CA8	L	11. 168(0.13)	MINC
000246	38 62	LR	6. 2	
00024E	5E 60 0 C18	L	6. 184(0.13)	
00024C	58 50 0 C00	L	5. 152(0.13)	
000250	18 28	100009 LR	2.11	
000252	10 28	SR	2. 8	
000254	18 52	LR	5. 2	
000256	1A 93	AR	5. 3	
00025E	1E 25	LR	2. 5	
00025A	65 20 0 C02	SL	2. 2	
00025E	7E 22 5 000	LE	2. 0(2. 5)	F
000262	7C 27 6 C00	ME	2. 0(7. 6)	Z
000266	3A 62	AF	6. 2	
000268	1A 74	20 AF	7. 4	
00026A	87 8A 0 229	FXLE	E.10. 5E2(13)	100005
00026E	7C 60 0 C7C	STE	6. 124(0.13)	S
000272	58 70 0 000	100010 L	7. 208(0.13)	
000276	5E 80 0 018	L	6. 184(0.13)	
00027A	5E 80 0 C08	L	11. 216(0.13)	
00027E	5E 90 0 0F4	L	5. 228(0.13)	4
000282	5E E0 0 CEC	L	14. 228(0.13)	002
000286	78 20 0 07C	LE	2. 124(0.13)	S
00028A	70 2E 0 C00	CE	2. 0(14.11)	ALPHA
00028E	33 02	LCER	0. 2	
000290	70 0E 7 000	STE	6. 0(14. 7)	RC
000294	5E 40 0 DAB	L	4. 168(0.13)	MINC
000298	6E 40 0 C2D	SDDA	4. 32	
00029C	1E 43	ER	4. 2	
00029E	1E 2A	LR	2.10	
0002A0	1A 25	AR	2. 5	
0002A2	50 30 0 018	ST	2. 136(0.13)	IF
0002A6	78 0E 7 C00	LE	0. 0(14. 7)	RC
0002AA	70 00 0 0E0	SL	6. 224(0.13)	001
0002AF	18 49	LR	4. 5	
0002B0	5E 50 0 CA8	L	5. 168(0.13)	MINC
0002B4	58 70 0 018	L	7. 240(0.13)	P
0002B8	5E 80 0 CEB	L	6. 136(0.13)	IF
0002BC	1E 82	LR	11. 2	
0002BE	58 90 0 018	L	5. 184(0.13)	
0002C2	18 25	100011 LR	2. 5	
0002C4	10 28	SR	2. 8	
0002C6	18 62	LR	6. 2	
0002CB	1A 63	AT	6. 3	
0002CA	18 26	LR	2. 6	
0002CC	65 20 0 C02	SL	2. 2	
0002C0	78 62 5 000	LE	6. 0(2. 5)	Z
0002D4	3E 26	LR	2. 6	
0002D6	7C 20 0 C1C	ME	2. 224(0.13)	001
0002DA	7E 47 5 000	L	4. 0(7. 6)	Z

0002DE	3A 24	ALR	2, 4	
0002E0	7C 40 D 0E0	ME	4, 224(0,13)	.CO1
0002E4	3A 4E	AFR	4, 6	
0002E6	70 42 9 000	STE	2, 0(2, 5)	A
0002EA	70 27 9 000	STE	2, 0(7, 9)	A
0002EE	1A 7A	20 AR	1, 4	
0002F0	87 8A D 29A	EXLE	F,10, 0(0,13)	100C11
0002F4	5E 70 D 010	100C12 L	7, 20(0,13)	
0002F8	58 80 C 010	L	8, 184(0,13)	
0002FC	5E 80 D 008	L	11, 21(0,13)	
000300	58 50 D 0FA	L	5, 22(0,13)	A
000304	5E 80 D 01C	L	14, 23(0,13)	.CO2
00030E	78 00 D 0E0	LE	0, 224(0,13)	.CO1
00030C	70 0E A 004	STE	0, 4(14, 8)	A
000310	78 2E B 000	LE	2, 0(14,11)	ALFFA
000314	38 42	LER	4, 2	
000316	7C 40 D 0E0	ME	4, 224(0,13)	.CO1
00031A	7C 40 D 0E0	ME	4, 224(0,13)	.CO1
00031E	3D 24	SEP	2, 4	
000320	70 2E B 004	STE	2, 4(14,11)	ALFFA
000324	78 0E B 000	LE	0, 0(14,11)	ALFFA
000328	2E 00	LTER	0, 0	
00032A	47 30 D 30A	RC	3, 77(0,13)	4C
00032E	47 F0 D 32C	FC	1E, 80(0,13)	CC
000332	58 00 D 0E0	40 L	0, 22(0,13)	.CO2
000336	1A 05	AR	0, 5	
000338	50 C0 D 0EC	ST	0, 23(0,13)	.CO2
00033C	58 20 D 0A8	L	2, 16(0,13)	MINC
000340	1A 2A	AF	2,10	
000342	50 20 D 0A8	ST	2, 16(0,13)	MINC
000346	59 20 D 074	C	2, 11(0,13)	A
00034A	47 C0 D 20A	FC	12, 22(0,13)	100C0E
00034E	58 70 D 008	60 L	7, 21(0,13)	
000352	58 80 D 000	L	6, 20(0,13)	
000356	5E 50 D 01E	L	5, 184(0,13)	
00035A	58 E0 C 0C0	L	11, 152(0,13)	
00035E	58 C0 D 0E4	L	6, 22(0,13)	A
000362	1B FF	SF	1E,1E	
000364	58 E0 D 000	L	14, 0(0,13)	
000368	07 FF	FCR	1E,14	
ADDRESS OF EPILOGUE				
00036A	58 D0 D 004	L	12, 4(0,13)	
00036E	58 F0 D 00C	L	14, 12(0,13)	
000372	52 FF C 00C	MVI	12(13),2E5	
000376	5E 2C D 01C	LM	2,12, 2(13)	
00037A	C7 FE	HCR	1E,14	
ADDRESS OF PROLOGUE				
00037C	5E 70 1 000	LM	7,11, 0(1)	
000380	58 20 7 000	L	2, 0(0, 7)	
000384	50 20 D 07E	ST	2, 120(0,13)	A
000388	5E 20 5 000	L	2, 0(0, 9)	
00038C	50 20 C 090	ST	2, 144(0,13)	AA
000390	58 20 1 004	L	2, 4(0, 1)	
000394	41 30 2 000	LA	2, 0(0, 2)	
000398	41 50 0 004	LA	5, 4	
00039C	1B 25	SR	2, 5	
0003A0	E0 20 C 00E	ST	2, 20(0,13)	
0003A2	50 30 C 00C	ST	2, 204(0,13)	A

```

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINESCOUNT(C) SIZE(0750K) ALTCODE(LNCF)
*OPTIONS IN EFFECT*SOURCE EBCCIC LIST CHECK EJECT P/D FORMAT GDSMT XREF NCALC NMANSE IFRM IRM FLACED
*STATISTICS* SOURCE STATEMENTS = 44. PROGRAM SIZE = 1024. SUBPROGRAM NAME = AUTC
*STATISTICS* J DIAGNOSTIC GENERATOR. HIGHEST SEVERITY CODE IS: 4
***** END LI COMPILATION *****
      3004 BYTES IN CORE NOT USED

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***** I N T R A N C R O S S R E F E R E N C E L I S T I N G *****
 SYMBOL INTERNAL STATEMENT NUMBERS
 MM 0005 0006 0017
 MC 0002 0003 0012 0015
 MD 0002 0004 0017
 INVER5 0002

***** D E F I N E D C R O S S R E F E R E N C E L I S T I N G *****
 LABEL DEFINED REFERENCES
 10 0008 0006
 20 0010 0005
 30 0013 0011
 40 0015 0014

000000 47 F0 F C0C INVER5 DC 15,12(0,15)
 000004 07 EC XL1*(7)
 000005 C9E5E5C9E5E240 EC (L7*INVER5)
 00000C 50 EC D C0C SIM 14,12,12(13)
 000010 18 4D LF 4,13
 000012 58 CD F C20 IF 12,13,22(15)
 000016 50 40 D 004 ST 4,4(0,12)
 00001A 50 00 4 C0H ST 12,E(0,4)
 00001E 07 FC PCR 12,12

CONSTANTS

000078 C0C0C000 DC XL4*0000C0C0*
 00007C 00C00001 DC XL4*C0C00C01*
 000080 00C0C000 DC XL4*000000C0*
 0000E4 C0C0C000 DC XL4*C0C0C000*

ADCLN FOR VARIABLES AND CONSTANTS
 ADCLR FOR EXTERNAL REFERENCES

0000F0 00C00000 EC XL4*C0C000C0* A
 0000F8 00C00000 EC XL4*00C000C0* F
 000100 C0C0C000 EC XL4*000000C0* J
 000108 00C00000 DC XL4*C000C0C0* RC
 000110 00C00000 EC XL4*00000000* RC
 000138 58 70 D C0C 100001 L 7, 224(0,13)
 00013C 58 80 C 0C8 L 8, 200(0,13)
 000140 78 60 D C50 LF 6, 80(0,13) C
 000144 58 00 D C68 L 11, 104(0,13) A
 000148 58 A0 D 0F0 L 10, 24(0,13) A
 00014C 58 20 D C54 L 2, 84(0,13) I
 000150 1A 2B AR 2,11
 000152 50 20 D 070 ST 2, 112(0,13) MA
 00015C 89 20 C C02 SLL 2, 2
 00015A 50 20 D 0FB ST 2, 24(0,13) *C02
 00015E 18 5A LF 5,10
 000160 1E B2 LF 11, 2
 000162 70 69 C 074 100002 STE 6, 110(5,13) F
 000166 87 5A D 12A 10 EXLE 5,10, 214(12) 10C02
 00016A 58 00 D C68 100003 L 11, 104(0,13) M
 00017E 18 2B CR 2,11
 000170 89 20 C C02 SLL 2, 2
 000174 50 20 D 114 ST 2, 20(0,13) *C05
 000178 5C 20 C 0FC ST 2, 252(0,13) *C02

00017C	58 00 0 0FB	L	0, 240(0,13)	.C02
000180	5C 00 0 10B	SI	0, 264(0,13)	.C06
000184	58 50 0 0CC	L	0, 100(0,13)	A
00018E	05 50 0 002	SLL	0, 2	
00019C	18 2A	LR	0, 10	
0001EE	18 4A	LR	4, 10	
000190	1E 02	LR	0, 2	
000192	58 30 0 00B	L	0, 216(0,13)	
000196	7E 26 3 000	100004 LE	0, 0(0, 2)	X
00019A	70 2C 8 004	STE	0, 4(0, 0)	A
00019E	1E 54	LR	0, 4	
0001A0	18 4A	LR	10, 4	
0001A2	58 00 0 104	L	11, 260(0,13)	.C05
0001A6	78 29 0 074	100005 LE	0, 116(0,13)	E
0001AA	7C 25 7 0CC	ME	0, 0(0, 7)	RC
0001AE	7A 29 8 000	AE	0, 0(0, 0)	A
0001B2	70 25 8 004	STE	0, 4(0, 0)	A
0001B6	07 5A 0 17E	30 BXLE	0, 10, 102(13)	10000E
0001EA	1E 54	10000C LR	0, 4	
0001EC	1E 54	LR	10, 4	
0001EE	58 00 0 0FC	L	11, 252(0,13)	.C03
0001F2	78 29 8 000	40 LE	0, 0(0, 0)	A
0001F6	7C 25 7 0CC	ME	0, 0(0, 7)	RC
0001FA	7A 29 0 074	AF	0, 116(0,13)	E
0001FE	70 25 0 07B	STE	0, 120(0,13)	E
000102	07 5A 0 17A	BXLE	0, 10, 410(13)	40
000106	78 20 0 004	100007 LE	0, 4(0, 0)	A
00010A	70 20 0 07B	STE	0, 120(0,13)	E
00010E	5E 00 0 10B	L	14, 264(0,13)	.C0C
000112	78 2E 8 000	LE	0, 0(14, 0)	A
000116	58 00 0 00B	L	10, 232(0,13)	
00011A	70 26 0 000	STF	0, 0(0, 0)	FC
00011E	07 64 0 16E	20 EXLE	0, 4, 300(13)	100004
0001F2	58 00 0 00B	100009 L	11, 104(0,13)	M
0001F6	1E FF	SR	10, 10	
0001FA	58 00 0 000	L	14, 0(0, 13)	
0001FC	07 FE	RCP	10, 14	
ADDRESS OF EPILOGUE				
0001FE	58 00 0 004	L	12, 4(0, 13)	
000202	58 00 0 00C	L	14, 12(0, 13)	
000206	92 FF 0 00C	MVI	12(13), 255	
00020A	5E 2C 0 01C	LM	0, 12, 20(13)	
00020E	07 FE	RCP	10, 14	
ADDRESS OF PROLOGUE				
000210	58 7B 1 000	LM	7, 11, 0(0)	
000214	58 20 7 000	L	0, 0(0, 7)	
000218	50 20 0 06C	ST	0, 100(0, 13)	A
00021C	5E 20 5 0CC	L	0, 0(0, 5)	
000220	50 20 0 06B	ST	0, 104(0, 13)	M
000224	58 20 1 004	L	0, 4(0, 1)	
000228	41 30 0 000	LA	0, 0(0, 2)	
00022C	41 50 0 004	LA	0, 4	
000230	1B 25	SF	0, 5	
000232	50 20 0 00B	ST	0, 216(0, 13)	
000236	50 30 0 00C	SI	0, 220(0, 13)	X
00023A	5E 20 1 00C	I	0, 12(0, 1)	
00023E	41 30 0 000	LA	0, 0(0, 2)	
000242	41 50 0 004	LA	0, 4	

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000246	1B 25	SR 2, 5		
000248	50 20 D 000	ST 2, 200(0,13)		
00024C	50 30 D 004	ST 3, 212(0,13)		
000250	58 20 1 01C	L 2, 16(0, 1)		R
000254	41 30 2 000	LA 3, 0(0, 2)		
00025E	41 50 0 004	LA 5, 4		
00025C	1E 25	SR 2, 5		
00025E	50 20 C 009	ST 4, 200(0,13)		
000262	50 30 D 00C	ST 3, 204(0,13)		A
000266	58 20 1 014	L 2, 20(0, 1)		
00026A	41 30 2 000	LA 3, 0(0, 2)		
00026E	41 50 0 004	LA 5, 4		
000272	1E 25	SR 2, 5		
000274	50 20 D 0E0	ST 2, 224(0,13)		
00027E	50 30 D 0E4	ST 3, 228(0,13)		PC
00027C	58 20 1 01E	L 2, 24(0, 1)		
000280	41 30 2 000	LA 3, 0(0, 2)		
0002E4	41 50 0 004	LA 5, 4		
000288	1B 25	SR 2, 5		
0002EA	50 20 D 0E8	ST 2, 232(0,13)		
00028E	50 30 D 0E0	ST 3, 236(0,13)		PC
000292	47 F0 D 110	EC 15, 272(0,13)		
ADCON FOR PROLOGUE	000020 0000210	DC XL4*0000210*		
ADCON FOR SAVE AREA	000024 000002E	DC XL4*000002E*		
ADCON FOR EPILOGUE	000028 000001E	DC XL4*000001E*		
TEMPORARIES AND GENERATED CONSTANTS				
000118	00000004	DC XL4*00000004*		
00011C	00000000	DC XL4*00000000*		
000120	00000000	DC XL4*00000000*		
000124	00000000	DC XL4*00000000*		
000128	00000000	DC XL4*00000000*		
00012C	00000000	DC XL4*00000000*		
000130	00000000	DC XL4*00000000*		
000134	00000000	DC XL4*00000000*		

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINECOUNT(0) SIZE(0750K) ALTCODE(NONE)

*OPTIONS IN EFFECT*SOURCE EPCCIC LIST NOCHECK OBJECT MAP FCNAT COSTIME NONE NALC NOANSI TERM IMP FLAGED

STATISTICS SOURCE STATEMENTS = 15, PROGRAM SIZE = 672, SUFFIXES NAME = INVS

STATISTICS NO DIAGNOSTICS GENERATED

***** END OF COMPILATION *****

2108 BYTES OF CORE NOT USED

REQUESTED OPTIONS: OPT=2,FORMAT,XREF,LIST,PAG,SIZE(750K)

OPTIONS IN EFFECT: NAME(MAIN) OPTIMIZE(2) LINECOUNT(20) SIZE(0750K) ALI(0)BL(0)NONE

SOURCE EXECIO LIST NOEXEC OBJECT XREF FORMAT GUSMT XREF NCALC NOANSI TERM IBM FLAC(1)

```

C
C
ISN 0002      C      SUBROUTINE SYNTHZIN(M,RC,YC,Y)
C
C      THIS ROUTINE IMPLEMENTS THE SYNTHESIS LATTICE FILTER.
C      ITS OUTPUT IS THE ORIGINAL INPUT RECONSTRUCTED FROM THE
C      PREDICTION RESIDUAL.
C
C      RC - REFLECTION COEFFICIENTS
C      RC - PREDICTION RESIDUAL
C      M - ORDER OF THE FILTER
C      N - NUMBER OF SAMPLES OF THE PRED. RESIDUAL
C      Y - RECONSTRUCTED SIGNAL
C
C
ISN 0003      C      DIMENSION A(20),E(20),FC(20),YC(256),Y(256)
C
ISN 0004      C      MM=MM1
ISN 0005      C      DO 100 J=1,MM
ISN 0006      C      B(J)=C.
ISN 0007      100  CCONTINUE
C
ISN 0008      C      DO 200 J=1,N
ISN 0009      C      A(MM)=FC(J)
ISN 0010      C      DO 300 I=1,M
ISN 0011      C      IP1=I-1
ISN 0012      C      A(IP1)=A(IP1+1)-E(IP1)*RC(IP1)
ISN 0013      300  CCONTINUE
ISN 0014      C      DO 400 I=1,M
ISN 0015      400  B(I+1)=E(I)/(1+RC(I))
ISN 0016      C      E(I)=A(I)
ISN 0017      C      Y(J)=A(I)
ISN 0018      200  CCONTINUE
ISN 0019      C      RETURN
ISN 0020      C      END
    
```

```

*****FORTRAN CROSS REFERENCE LISTING*****
SYMBOL  INTERNAL STATEMENT NUMBERS
A      0003 0009 0012 0012 0015 0016 0017
B      0003 0006 0012 0015 0015 0016
C      0010 0011 0014 0015 0015 0015 0015
J      0005 0006 0006 0005 0017
M      0002 0004 0010 0011 0014
N      0002 0003
Y      0002 0003 0017
MM     0004 0005 0005
RC     0002 0003 0012 0015
RU     0002 0003 0005
IP1    0011 0012 0012 0012 0012
    
```

*****FORTRAN CROSS REFERENCE LISTING*****
SYMBOL INTERNAL STATEMENT NUMBERS
SYNTHZ 0002

*****FORTRAN CROSS REFERENCE LISTING*****

LABEL	DEFINED	REFERENCES
100	CCC7	0005
200	0018	0008
300	0013	0010
400	0015	0014

000000	47 F0 F 00C	SYNTHZ	FC	15,12(0,15)
000004	07		DC	XL1'07'
000005	E2E05E10E940		EC	CL7'SYNTHZ'
00000C	90 EC D 00C		STM	14,12,12(13)
000010	18 40		LR	4,12
000012	5E CD F 020		LN	12,13,22(15)
000016	50 40 D 0C4		ST	4,4(0,13)
00001A	50 00 4 008		SI	13,8(0,4)
00001E	C7 FC		FCR	15,12

CONSTANTS

000078	0000000	DC	XL4'00000000'
00007C	0000001	DC	XL4'00000001'
000080	00000000	DC	XL4'00000000'
000084	00000000	EC	XL4'00000000'

ADDONS FOR VARIABLES AND CONSTANTS
ADDONS FOR EXTERNAL REFERENCES

000148	0000000	DC	XL4'00000000'	Y	
000150	00000000	DC	XL4'00000000'	FC	
000158	00000000	DC	XL4'00000000'	RE	
00017C	5E 70 D 12E	100001	L	7, 29(0,13)	
000180	78 60 D 050		LF	6, EC(0,13)	0
0001E4	5E 80 D 06E		L	11, 104(0,13)	N
000188	5E 40 D 128		L	4, 212(0,13)	4
00018C	5E 30 D 054		L	3, 84(0,13)	1
000150	18 23		LF	2, 2	
000192	1A 20		AR	1,11	
000194	50 20 D 070		ST	2, 112(0,13)	NA
000158	89 20 0 002		SLL	2, 2	
00015C	50 20 D 150		ST	2, 32(0,13)	COE
0001A0	1E 54		LR	5, 4	
0001A2	1E A4		LR	10, 4	
0001A4	18 12		LR	11, 2	
0001A6	70 65 D 00E	100002	STF	6, 200(0,13)	U
0001AA	87 9A E 17E	100	PXLE	5,10, 382(13)	100002
0001AE	5E 60 D 00E	100001	L	11, 104(0,13)	N
0001B2	5E 00 D 150		L	6, 32(0,13)	COE
0001B6	50 00 D 148		SI	6, 32(0,13)	COA
0001EA	1E 00		LF	6,11	
0001BC	65 00 0 002		SLL	6, 2	
0001C0	50 00 D 140		SI	6, 32(0,13)	CO2
0001C4	5E 20 D 00C		L	5, 108(0,13)	N
0001C8	65 20 0 002		SLL	5, 2	
0001CC	18 24		LF	2, 4	

0001C0	18 C2		LF	6, 2	
0001D0	5E F0 D 120	10CC04	L	15, 204(0,13)	
0001D4	78 26 F 000		LE	2, 0(0,15)	FC
0001DE	5E F0 D 11F		L	14, 228(0,13)	.C04
0001DC	70 2E D 07E		STE	2, 120(14,13)	A
0001E0	18 53		LF	5, 2	
0001E2	18 A3		LF	10, 2	
0001E4	58 E0 D 0E8		L	11, 104(0,13)	A
0001E8	13 2F	100005	LR	2, 11	
0001EA	18 29		SR	2, 5	
0001EC	18 E2		LR	6, 2	
0001EE	1A EA		AR	8, 10	
0001F0	1E 28		LR	2, 8	
0001F2	89 20 0 002		SLL	2, 2	
0001F6	78 22 D CCE		LE	2, 200(2,13)	B
0001FA	7C 22 7 000		ME	2, 0(2, 7)	FC
0001FE	33 22		LCER	2, 2	
000200	7A 22 D C7C		AE	2, 124(2,13)	A
000204	70 22 C 070		STE	2, 120(2,13)	A
000208	E7 5A D 1C0	200	EXLE	5, 10, 454(13)	10CCCE
00020C	18 94	10CC0E	LF	5, 4	
00020E	18 A4		LR	10, 4	
000210	5E B0 D 140		L	11, 320(0,13)	.C02
000214	18 E7		LR	8, 7	
000216	78 29 D C7B	400	LE	2, 120(5,13)	A
00021A	7C 25 B C00		ME	2, 0(5, 8)	RC
00021E	7A 29 D 0CE		AE	2, 200(5,13)	E
000222	70 29 D 0CC		STE	2, 204(5,13)	E
000226	E7 5A D 1EE		EXLE	5, 10, 454(13)	400
00022A	78 20 D 07C	10CC07	LE	2, 124(0,13)	A
00022E	70 20 D 0CC		STE	2, 204(0,13)	E
000232	58 F0 D 120		L	12, 288(0,13)	
000236	70 26 F 000		STE	2, 0(0,15)	Y
00023A	E7 64 D 1AE	200	EXLE	6, 4, 424(13)	10CCCA
00023E	58 B0 D C6B	10CC0E	L	11, 104(0,13)	A
000242	18 FF		SR	12, 15	
000244	58 E0 D C00		L	14, 0(0,13)	
000248	07 FE		BCR	12, 14	
ADDRESS OF EPILOGUE					
00024A	58 D0 D 004		L	12, 4(0,13)	
00024E	58 E0 C 00C		L	14, 12(0,13)	
000252	52 FF D 00C		MVI	12(13), 2E5	
000256	58 2C D 01C		LM	2, 12, 2E(13)	
00025A	07 FE		BCR	12, 14	
ADDRESS OF PROLOGUE					
00025C	5E 78 1 000		LM	2, 11, 0(1)	
000260	5E 20 7 C00		L	2, 0(0, 7)	
000264	50 20 D C6C		ST	2, 108(0,13)	A
000268	5E 20 B 000		L	2, 0(0, 8)	
00026C	50 20 D C6E		ST	2, 104(0,13)	A
000270	5E 20 1 C0E		L	2, 0(0, 1)	
000274	41 30 2 000		LA	2, 0(0, 2)	
000278	41 50 0 C01		LA	5, 4	
00027C	18 2E		SR	2, 5	
00027E	50 20 D 120		ST	2, 290(0,13)	
000282	50 30 D 12C		ST	2, 200(0,13)	BC
000286	5E 20 1 00C		L	2, 12(0, 1)	
00028A	41 30 2 C00		LA	2, 0(0, 2)	

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	0002EE	41 50 0 004	LA	5, 4	
	0002E2	1B 25	SR	2, 5	
	000294	50 20 D 130	ST	2, 304(0,13)	
	000258	50 30 D 124	ST	3, 208(0,13)	PC
	00029C	58 20 1 010	L	2, 16(0, 1)	
	0002A0	41 30 2 000	LA	3, 0(0, 2)	
	0002A4	41 50 0 004	LA	5, 4	
	0002A8	1B 25	SR	2, 5	
	0002AA	50 20 D 120	ST	2, 288(0,13)	
	0002AE	50 30 D 124	ST	3, 292(0,13)	Y
	0002B2	47 40 D 154	DC	15, 340(0,13)	
ADCN FOR PROLOGUE	000020	C0C025C	DC	XL4*C0C0025C*	
ADCN FOR SAVE ARLA	000024	00C0028	DC	XL4*C000C028*	
ADCN FOR EPILOGUE	000028	C0C024A	DC	XL4*C0C0C24A*	
TEMPORARIES AND GENERATED CONSTANTS	0001C0	00C00004	DC	XL4*C0C00004*	
	000164	00C00000	DC	XL4*00000000*	
	000168	00C00000	DC	XL4*00C000C0*	
	00016C	00C00000	DC	XL4*00C000C0*	
	000170	00C00000	DC	XL4*00000000*	
	000174	00C00000	DC	XL4*00000000*	
	000178	00C00000	DC	XL4*00000000*	

*OPTIONS IN EFFECT*NAME(MAIN) OPTIMIZE(2) LINECNT(0) SIZE(07500) AUTOCPL(NONE)

*OPTIONS IN EFFECT*SOURCE FBCCIC LIST ALUECK OBJECT MAP FORMAT GUEINT PREF NCALC NOANSI TERM IBM FLAG(1)

STATISTICS SOURCE STATEMENTS = 15, PROGRAM SIZE = 694, SUPPFC(FAN NAME =SYNTH2

STATISTICS NO DIAGNOSTIC GENERATE

***** END OF COMPILATION *****

316K BYTES OF CORE NOT USED

APPENDIX D

SATELLITE COMMUNICATION TERMINOLOGIES

Transponder. The equipment which receives a signal, amplifies it, changes its frequency, and retransmits it. This process is necessary in order to avoid interfering the weak incoming signal with the powerful transmitted signal. Most satellites have more than one transponder, and the bandwidth of most transponders is 36 MHz, which can carry one televised signal.

Earth Station. A large dish-shaped antenna pointed toward a satellite. The size of the dish usually varies with the size of the beam angle. Most earth stations simply transmit and receive the telecommunication signals with a fixed antenna. The most inexpensive channels operate with a UHF channel.

Channel Capacity (C). The maximum rate at which nearby errorless data transmission is theoretically possible. For certain types of communication channels it has been shown that [14]

$$C = B \log_2 (1 + S/N)$$

where B is the bandwidth, S is the signal power, and N is the noise power. For example, for a 3000 Hz of bandwidth and $S/N = 10^3$, the capacity of channel needs to be approximately 30 k bits/second. At the present time, the actual data rate on such channels ranges from 150 to 9000 bits/second [23].

Multiplexing. Any technique which permits more than one independent signal to share one physical facility. This may include space, frequency, and/or time division multiplexing. In a satellite, a high level of multiplexing is needed so that many signals can share the large bandwidth.

Access Methods of Satellite Links [34, 35]. A satellite interconnects large numbers of earth stations scattered over thousands of miles.

An efficient solution to allocate subchannels to many users (multiple-access) is required. The simplest way to subdivide satellite capacity by frequency is to give different users different transponders, simply referred to as the multiple-transponder technique. This approach is less satisfactory for the users who need much smaller capacity and/or variable channel assignment. For this reason, techniques which allow for sharing the same transponder among many earth stations have been developed. They are frequency division multiple access (FDMA) and time division multiple access (TDMA).

FDMA. With FDMA the transponder bandwidth is divided into smaller bandwidths. An earth station transmits on one or more of these divisions. The signal is used to modulate a carrier and multiple carriers are used so that signals can be spaced from each other. Figure 40 illustrates one example of FDMA, where 800 channel carriers are spaced 45 KHz apart. The first channel is used to control the allocation of voice carriers to earth stations. A carrier at the center of each slot is modulated with the voice channel. The more carriers a transponder shares with this scheme the lower the overall capacity. This is due to the effect of guard bands between the carrier's bands and intermodulation.

TDMA. One of the objectives of TDMA is to employ a single carrier for the transmission via one transponder. With TDMA, each earth station is allowed to transmit a high-speed burst of bits for a brief period of time. The times of bursts are carefully controlled so that no two bursts overlap. For the period of its bursts the earth station has the entire transponder bandwidth available to it. The set of bursts is illustrated in Figure 41. The first burst in a frame contains no

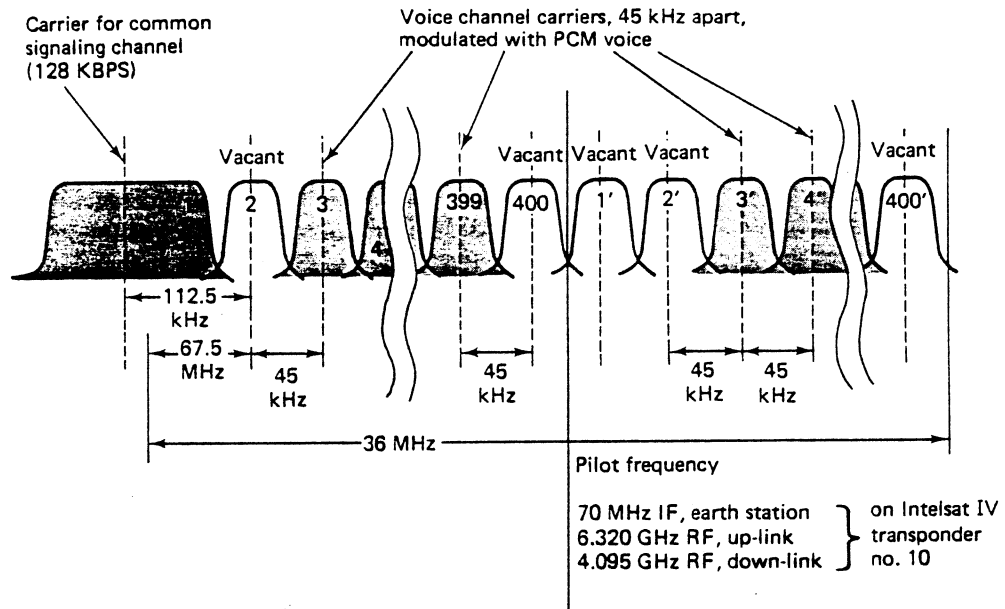


Figure 40. Frequency Multiple Access Method Example (After Martin [35])

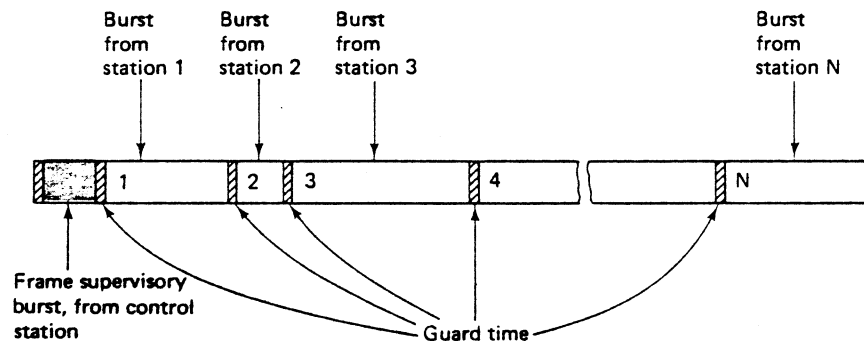


Figure 41. Time Division Multiple Access Method Example (After Martin [35])

traffic but serves to synchronize and identify a frame. This method is superior to FDMA in such aspects as flexibility in channel capacity, no intermodulation interference, and no transponder saturation. However, TDMA has other types of problems associated with the synchronization and control of the high-speed digital bit stream.

2
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